## The AtariBASIC



A complete explanation of the inside workings of Atari BASIC, along with the original source code. For intermediate and advanced programmers.

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From COMPUTE! Books and Optimized Systems Software, Inc.

## The Atari BASIC

 SOURCE

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# Publisher's Foreword 

It's easy to take a computer language like Atari BASIC for granted. But every PEEK and POKE, every FOR-NEXT loop and IF-THEN branch, is really a miniprogram in itself. Taken together, they become a powerful tool kit. And, as Atari owners know, there are few home-computer languages as powerful and versatile - from editing to execution - as Atari BASIC.

With this book, the Atari BASIC tool kit is unlocked. The creators of Atari BASIC and COMPUTE! Publications now offer you, for the first time, a detailed, inside look at exactly how a major computer manufacturer's primary language works.

For intermediate programmers, the thorough and careful explanations in Parts 1 and 2 will help you understand exactly what is happening in your Atari computer as you edit and run your programs.

For advanced programmers, Part 3 provides a complete listing of the source code for Atari BASIC, so that your machine language programs can make use of the powerful routines built into that 8 K cartridge.

And for programmers at all levels, by the time you're through studying this book you'll feel that you've seen a whole computer language at work.

Special thanks are due to Bill Wilkinson, the creative force behind Atari BASIC and many other excellent programs for Atari and other computers, for his willingness to share copyrighted materials with computer users. Readers of COMPUTE! Magazine already know him as a regular columnist, and in this book he continues his tradition of clear explanations and understandable writing.

## Acknowledgments

As far as we know, this is the first time that the actual source listing of a major manufacturer's primary computer language has been made available to the general public.

As with our previous COMPUTE! Publications book Inside Atari DOS, this book contains much more than simply a source listing. All major routines are examined and explained. We hope that when you finish reading this book you will have a better understanding of and appreciation for the design and work which go into as sophisticated a program as Atari BASIC.

This book is the result of the efforts of many people. The initial credit must go to Richard Mansfield of COMPUTE! Publications for serving as our goad and go-between. Without his (and COMPUTE!'s) insistence, this book might never have been written. Without his patience and guidance, the contents of this book might not have been nearly as interesting.

To Kathleen O'Brien and Paul Laughton must go the lion's share of the authoring credits. Between them, they have done what I believe is a very creditable job of explaining a very difficult subject, the internal workings of Atari BASIC. In fact, Part I of this book is entirely their work. Of course, their ability to explain the listing may not be so surprising. After all, between them they wrote almost all of the original code for Atari BASIC. So, even though Paul and Kathleen are not associated with Optimized Systems Software, we were pleased to have their invaluable help in writing this book and hope that they receive some of the credit which has long been due them.

Mike Peters was responsible for taking our old, almost unreadable copies of the source code diskettes for Atari BASIC and converting them to another machine, using another assembler, and formatting the whole thing into an acceptable form for this book. This isn't surprising either, since Mike keypunched the original (yes, on cards).

And I am Bill Wilkinson, the one responsible for the rest of this book. In particular, I hope you will find that a good amount of the material in Part II will aid you in understanding how to make the best use of this book.

The listing of Atari BASIC is reproduced here courtesy of OSS, Inc., which now owns its copyright and most other associated rights.

## Preface

In 1978, Atari, Inc., purchased a copy of Microsoft BASIC for the 6502 microprocessor (similar to the version from which Applesoft is derived). After laboring for quite some time, the people of Atari still couldn't make it do everything they wanted it to in the ROM space they had available. And there was a deadline fast approaching: the January 1979 Las Vegas Consumer Electronics Show (CES).

At that time, Kathleen, Paul, Mike and I all worked for Shepardson Microsystems, Inc. (SMI). Though little known by the public, SMI was reasonably successful in producing some very popular microcomputer software, including the original Apple DOS, Cromemco's 16K and 32K BASICs, and more. So it wasn't too surprising that Atari had heard of us.

And they asked us: Did we want to try to fix Microsoft BASIC for them? Well, not really. Did we think we could write an all-new BASIC in a reasonable length of time? Yes. And would we bet a thousand dollars a week on our ability to do so?

While Bob Shepardson negotiated with Atari and I wrote the preliminary specifications for the language (yes, I'm the culprit), time was passing all too rapidly. Finally, on 6 October 1978, Atari's Engineering Department gave us the okay to proceed.

The schedule? Produce both a BASIC and a Disk File Manager (which became Atari DOS) in only six months. And, to make sure the pressure was intense, they gave us a $\$ 1000-\mathrm{a}-$ week incentive (if we were early) or penalty (if we were late).

But Paul Laughton and Kathleen O'Brien plunged into it. And, although the two of them did by far the bulk of the work, there was a little help from Paul Krasno (who implemented the transcendental routines), Mike Peters (who did a lot of keypunching and operating), and me (who designed the floating point scheme and stood around in the way a lot). Even Bob Shepardson got into the act, modifying his venerable IMP-16 assembler to accept the special syntax table mnemonics that Paul invented (and which we paraphrase in the current listing via macros).

Atari delivered the final signed copy of the purchase order on 28 December 1978, two and a half months into the project. But it didn't really matter: Paul and Kathy were on vacation, having delivered the working product more than a week before!

So Atari took Atari BASIC to CES, and Shepardson Microsystems faded out of the picture. As for the bonus for early delivery - there was a limit on how much the incentive could be. Darn.

The only really unfortunate part of all this was that Atari got the BASIC so early that they moved up their ROM production schedule and committed to a final product before we had a chance to do a second round of bug fixing.

And now? Mike and I are running Optimized Systems Software, Inc. And even though Paul and Kathleen went their own way, we have kept in touch enough to make this book possible.

## Part One $\square$

## How Atari BASIC Works



# Atari BASIC: A High-Level Language Translator 

The programming language which has become the de facto standard for the Atari Home Computer is the Atari 8K BASIC Cartridge, known simply as Atari BASIC. It was designed to serve the programming needs of both the computer novice and the experienced programmer who is interested in developing sophisticated applications programs. In order to meet such a wide range of programming needs, Atari BASIC was designed with some unique features.

In this chapter we will introduce the concepts of high level language translators and examine the design features of Atari BASIC that allow it to satisfy such a wide variety of needs.

## Language Translators

Atari BASIC is what is known as a high level language translator. A language, as we ordinarily think of it, is a system for communication. Most languages are constructed around a set of symbols and a set of rules for combining those symbols.

The English language is a good example. The symbols are the words you see on this page. The rules that dictate how to combine these words are the patterns of English grammar. Without these patterns, communication would be very difficult, if not impossible: Out sentence this believe, of make don't this trying if sense you to! If we don't use the proper symbols, the results are also disastrous: @twu2 yeggopt gjsiem, keorw?

In order to use a computer, we must somehow communicate with it. The only language that our machine really understands is that strange but logical sequence of ones and zeros known as machine language. In the case of the Atari, this is known as 6502 machine language.

When the 6502 central processing unit (CPU) "sees" the sequence 01001000 in just the right place according to its rules of syntax, it knows that it should push the current contents of

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the accumulator onto the CPU stack. (If you don't know what an "accumulator" or a "CPU stack" is, don't worry about it. For the discussion which follows, it is sufficient that you be aware of their existence.)

Language translators are created to make it simpler for humans to communicate with computers. There are very few 6502 programmers, even among the most expert of them, who would recognize 01001000 as the push-the-accumulator instruction. There are more 6502 programmers, but still not very many, who would recognize the hexadecimal form of $01001000, \$ 48$, as the push-the-accumulator instruction. However, most, if not all, 6502 programmers will recognize the symbol PHA as the instruction which will cause the 6502 to push the accumulator.

PHA, \$48, and even 01001000, to some extent, are translations from the machine's language into a language that humans can understand more easily. We would like to be able to communicate to the computer in symbols like PHA; but if the machine is to understand us, we need a language translator to translate these symbols into machine language.

The Debug Mode of Atari's Editor/Assembler cartridge, for example, can be used to translate the symbols $\$ 48$ and PHA to the ones and zeros that the machine understands. The debugger can also translate the machine's ones and zeros to $\$ 48$ and PHA. The assembler part of the Editor/Assembler cartridge can be used to translate entire groups of symbols like PHA to machine code.

## Assemblers

An assembler - for example, the one contained in the Assembler/Editor cartridge - is a program which is used to translate symbols that a human can easily understand into the ones and zeros that the machine can understand. In order for the assembler to know what we want it to do, we must communicate with it by using a set of symbols arranged according to a set of rules. The assembler is a translator, and the language it understands is 6502 assembly language.

The purpose of 6502 assembly language is to aid program authors in writing machine language code. The designers of the 6502 assembly language created a set of symbols and rules that matches 6502 machine language as closely as possible.

This means that the assembler retains some of the

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disadvantages of machine language. For instance, the process of adding two large numbers takes dozens of instructions in 6502 machine language. If human programmers had to code those dozens of instructions in the ones and zeros of machine language, there would be very few human programmers.

But the process of adding two large numbers in 6502 assembly language also takes dozens of instructions. The assembly language instructions are easier for a programmer to read and remember, but they still have a one-to-one correspondence with the dozens of machine language instructions. The programming is easier, but the process remains the same.

## High Level Languages

High level languages, like Atari BASIC, Atari PILOT, and Atari Pascal, are simpler for people to use because they more closely approximate human speech and thought patterns. However, the computer still understands only machine language. So the high level languages, while seeming simple to their users, are really much more complex in their internal operations than assembly language.

Each high level language is designed to meet the specific need of some group of people. Atari Pascal is designed to implement the concept of structured programming. Atari PILOT is designed as a teaching tool. Atari BASIC is designed to serve both the needs of the novice who is just learning to program a computer and the needs of the expert programmer who is writing a sophisticated application program, but wants the program to be accessible to a large number of users.

Each of these languages uses a different set of symbols and symbol-combining rules. But all these language translators were themselves written in assembly language.

## Language Translation Methods

There are two different methods of performing language translation - compilation and interpretation. Languages which translate via interpretation are called interpreters. Languages which translate via compilation are called compilers.

Interpreters examine the program source text and simulate the operations desired. Compilers translate the program source text into machine language for direct machine execution.

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The compilation method tends to produce faster, more efficient programs than does the interpretation method. However, the interpretation method can make programming easier.

## Problems with the Compiler Method

The compiler user first creates a program source file on a disk, using a text editing program. Then the compiler carefully examines the source program text and generates the machine language as required. Finally, the machine language code is loaded and executed. While this three-step process sounds fairly simple, it has several serious " gotchas."

Language translators are very particular about their symbols and symbol-combining rules. If a symbol is misspelled, if the wrong symbol is used, or if the symbol is not in exactly the right place, the language translator will reject it. Since a compiler examines the entire program in one gulp, one misplaced symbol can prevent the compiler from understanding any of the rest of the program - even though the rest of the program does not violate any rules! The result is that the user often has to make several trips between the text editor and the compiler before the compiler successfully generates a machine language program.

But this does not guarantee that the program will work. If the programmer is very good or very lucky, the program will execute perfectly the very first time. Usually, however, the user must debug the program.

This nearly always involves changing the source program, usually many times. Each change in the source program sends the user back to step one: after the text editor changes the program, the compiler still has to agree that the changes are valid, and then the machine code version must be tested again. This process can be repeated dozens of times if the program is very complex.

## Faster Programming or Faster Programs?

The interpretation method of language translation avoids many of these problems. Instead of translating the source code into machine language during a separate compiling step, the interpreter does all the translation while the program is running. This means that whenever you want to test the program you're writing, you merely have to tell the interpreter to run it. If things don't work right, stop the program, make a few changes, and run the program again at once.

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You must pay a few penalties for the convenience of using the interpreter's interactive process, but you can generally develop a complex program much more quickly than the compiler user can.

However, an interpreter is similar to a compiler in that the source code fed to the interpreter must conform to the rules of the language. The difference between a compiler and an interpreter is that a compiler has to verify the symbols and symbol-combining rules only once - when the program is compiled. No evaluation goes on when the program is running. The interpreter, however, must verify the symbols and symbol-combining rules every time it attempts to run the program. If two identical programs are written, one for a compiler and one for an interpreter, the compiled program will generally execute at least ten to twenty times faster than the interpreted program.

## Pre-compiling Interpreter

Atari BASIC has been incorrectly called an interpreter. It does have many of the advantages and features of an interpretive language translator, but it also has some of the useful features of a compiler. A more accurate term for Atari's BASIC Language Translator is pre-compiling interpreter.

Atari BASIC, like an interpreter, has a text editor built into it. When the user enters a source line, though, the line is not stored in text form, but is translated into an intermediate code, a set of symbols called tokens. The program is stored by the editor in token form as each program line is entered. Syntax and symbol errors are weeded out at that time.

Then, when you run the program, these tokens are examined and their functions simulated; but because much of the evaluation has already been done, the execution of an Atari BASIC program is faster than that of a pure interpreter. Yet Atari BASIC's program-building process is much simpler than that of a compiler.

Atari BASIC has advantages over compilers and interpreters alike. With Atari BASIC, every time you enter a line it is verified for language correctness. You don't have to wait until compilation; you don't even have to wait until a test run. When you type RUN you already know there are no syntax errors in your program.

# Internal Design Overview 


#### Abstract

Atari BASIC is divided into two major functional areas: the Program Constructor and the Program Executor. The Program Constructor is used when you enter and edit a BASIC program. The source line pre-compiler, also part of the Program Constructor, translates your BASIC program source text lines into tokenized lines. The Program Executor is used to execute the tokenized program - when you type RUN, the Program Executor takes over.

Both the Program Constructor and the Program Executor are designed to use data tables. Some of these tables are already contained in BASIC's ROM (read-only memory). Others are constructed by BASIC in the user RAM (randomaccess memory). Understanding these various tables is an important key to understanding the design of Atari BASIC.


## Tokens

In Atari BASIC, tokens are the intermediate code into which the source text is translated. They represent source-language symbols that come in various lengths - some as long as 100 characters (a long variable name) and others as short as one character (" $+{ }^{\prime \prime}$ or " " - "). Every token, however, is exactly one eight-bit byte in length.

Since most BASIC Language Symbols are more than one character long, the representation of a multi-character BASIC Language Symbol with a single-byte token can mean a considerable saving of program storage space.

A single-byte token symbol is also easier for the Program Executor to recognize than a multi-character symbol, since it can be evaluated by machine language routines much more quickly. The SEARCH routine - 76 bytes long - located at $\$$ A462 is a good example of how much assembly language it takes to recognize a multi-character symbol. On the other hand, the two instructions located at $\$ \mathrm{AB} 42$ are enough to

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determine if a one-byte token is a variable. Because routines to recognize Atari BASIC's one-byte tokens take so much less machine language, they execute relatively quickly.

The 256 possible tokens are divided into logical numerical groups that also make them simpler to deal with in assembly language. For example, any token whose value is 128 ( $\$ 80$ ) or greater represents a variable name. The logical grouping of the token values also means faster execution speeds, since, in effect, the computer only has to check bit 7 to recognize a variable.

The numerical grouping of the tokens is shown below: Token Value (Hex) Description

| 00-0D | Unused |
| :--- | :--- |
| 0 E | Floating Point Numeric Constant. |
|  | The next six bytes will hold its value. |


| OF | String Constant. <br> The next byte is the string length. <br> A string of that length follows. |
| :--- | :--- |
| 10-3C | Operators. <br> See table starting at \$A7E3 for specific <br> operators and values. |

3D-54 Functions.
See table starting at \$A820 for specific functions and values.

55-7F Unused.
80-FF Variables.
In addition to the tokens listed above, there is another set of single-byte tokens, the Statement Name Tokens. Every statement in BASIC starts with a unique statement name, such as LET, PRINT, and POKE. (An assignment statement such as " $\mathrm{A}=\mathrm{B}+\mathrm{C}$," without the word LET, is considered to begin with an implied LET.) Each of these unique statement names is represented by a unique Statement Name Token.

The Program Executor does not confuse Statement Name Tokens with the other tokens because the Statement Name Tokens are always located in the same place in every statement - at the beginning. The Statement Name Token value is derived from its entry number, starting with zero, in the Statement Name Table at \$A4AF.

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## Tables

A table is a systematic arrangement of data or information. Tables in Atari BASIC fall into two distinct types: tables that are part of the Atari BASIC ROM and tables that Atari BASIC builds in the user RAM area.

## ROM Tables

The following is a brief description of the various tables in the Atari BASIC ROM. The detailed use of these tables will be explained in subsequent chapters.
Statement Name Table (\$A4AF). The first two bytes in each entry point to the information in the Statement Syntax Table for this statement. The rest of the entry is the name of the statement name in ATASCII. Since name lengths vary, the last character of the statement name has the most significant bit turned on to indicate the end of the entry. The value of the Statement Name Token is derived from the relative (from zero) entry number of the statement name in this table.
Statement Execution Table (\$AA00). Each entry in this table is the two-byte address of the 6502 machine language code which will simulate the execution of the statement. This table is organized with the statements in the same order as the statements in the Statement Name Table. Therefore, the Statement Name Token can be used as an index to this table. Operator Name Table (\$A7E3). Each entry comprises the ATASCII text of an Operator Symbol. The last character of each entry has the most significant bit turned on to indicate the end of the entry. The relative (from zero) entry number, plus 16 (\$10), is the value of the token for that entry. Each of the entries is also given a label whose value is the value of the token for that symbol. For example, the ";" symbol at \$A7E8 is the fifth (from zero) entry in the table. The label for the ";" token is CSC, and the value of CSC is $\$ 15$, or 21 decimal ( $1 \downarrow 16+5$ ).
Operator Execution Table (\$AA70). Each two-byte entry points to the address, minus one, of the routine which simulates the execution of an operator. The token value, minus 16 , is used to access the entries in this table during execution time: The entries in this table are in the same order as in the Operator Name Table.
Operator Precedence Table (\$AC3F). Each entry represents the relative execution precedence of an individual operator. The table entries are accessed by the operator tokens,

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minus 16. Entries correspond with the entries in the Operator Name Table. (See Chapter 7.)
Statement Syntax Table (\$A60D). Entries in this table are used in the process of translating the source program to tokens. The address pointer in the first part of each entry in the Statement Name Table is used to access the specific syntax information for that statement in this table. (See Chapter 5.)

## RAM Tables

The tables that BASIC builds in the user RAM area will be explained in detail in Chapter 3. The following is a brief description of these tables:
Variable Name Table. Each entry contains the source ATASCII text for the corresponding user variable symbol in the program. The relative (from zero) entry number of each entry in this table, plus 128, becomes the value of the token representing the variable.
Variable Value Table. Each entry either contains or points to the current value of a variable. The entries are accessed by the token value, minus 128.
Statement Table. Each entry is one tokenized BASIC program line. The tokenized lines are kept in this table in ascending numerical order by line number.
Array/String Table. This table contains the current values for all strings and numerical arrays. The location of the specific values for each string and/or array variable is accessed from information in the Variable Value Table.
Runtime Stack. This is the LIFO Runtime Stack, used to control the execution of GOSUB/RETURN and similar statements.

## Pre-compiler

Atari BASIC translates the BASIC source lines from text to tokens as soon as they are entered. To do this, Atari BASIC must recognize the symbols of the BASIC Language. BASIC also requires that its symbols be combined in certain specific patterns. If the symbols don't follow the required patterns, then Atari BASIC cannot translate the line. The process of checking a source line for the required symbol patterns is called syntax checking.

BASIC performs syntax checking as part of the tokenizing process. When the Program Editor receives a completed line of

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input, the editor hands the line to the syntax routine, which examines the first word of the line for a statement name. If a valid statement name is not found, then the line is assumed to be an implied LET statement.

The grammatical rules for each statement are contained in the Statement Syntax Table. A special section of code examines the symbols in the source line, under the direction of the grammatical rules set forth in the Statement Syntax Table. If the source line does not conform to the rules, then it is reported back as an error. Otherwise, the line is translated to tokens. The result of this process is returned to the Program Editor for further processing.

## Program Editor

When Atari BASIC is not executing statements, it is in the edit mode. When the user enters a source line and hits return, the editor accepts the line into a line buffer, where it is examined by the pre-compiler. The pre-compiler returns either tokens or an error text line.

If the line started with a line number, the editor inserts the tokenized line into the Statement Table. If the Statement Table already contains a line with the same line number, then the old line is removed from the Statement Table. The new line is then inserted just after the statement with the next lower line number and just before the statement with the next higher line number.

If the line has no line number, the editor inserts the line at the end of the Statement Table. It then passes control to the Program Executor, which will carry out the statement(s) in the line at the end of the Statement Table.

## Program Executor

The Program Executor has a pointer to the statement that it is to execute. When control is passed to the executor, the pointer points to the direct (command) line at the end of the statement table. If that statement causes some other line to be executed (RUN, GOTO, GOSUB, etc.), the pointer is changed to the new line. Lines continue to be executed as long as nothing stops that execution (END, STOP, error, etc.). When the program execution is stopped, the Program Executor returns control to the editor.

When a statement is to be executed, the Statement Name Token (the first code in the statement) directs the interpreter to the specific code that executes that statement. For instance, if that token represents the PRINT statement, the PRINT execution code is called. The execution code for each statement then examines the other tokens and simulates their operations.

## Execute Expression

Arithmetic and logical expressions ( $\mathrm{A}+\mathrm{B}, \mathrm{C} / \mathrm{D}+\mathrm{E}, \mathrm{F}<\mathrm{G}$, etc.) are simulated with the Execute Expression code. Expression operators (,,$+- *$, etc.) have execution precedence - some operators must be executed before some others. The expression $1+3^{*} 4$ has a value of 13 rather than 16 because * had a higher precedence than + . To properly simulate expressions, BASIC rearranges the expression with higher precedence first.

BASIC uses two temporary storage areas to hold parts of the rearranged expression. One temporary storage area, the Argument Stack, holds arguments - values consisting of constants, variables, and temporary values resulting from previous operator simulations. The other temporary storage area, the Operator Stack, holds operators. Both temporary storage areas are managed as Last-In/First-Out (LIFO) stacks.

## LIFO Stacks

A LIFO (Last In/First Out) stack operates on the principle that the last object placed in the stack storage area will be the first object removed from it. If the letters A, B, C, and D, in that order, were placed in a LIFO stack, then D would be the first letter removed, followed by C, B, and A. The operations required to rearrange the expression using these stacks will be explained in Chapter 7.

BASIC also uses another LIFO stack, the Runtime Stack, in the simulation of statements such as GOSUB and FOR. GOSUB requires that BASIC remember where in the statement table the GOSUB was located so it will return to the right spot when RETURN is executed. If more than one GOSUB is executed before a RETURN, BASIC returns to the statement after the most recent GOSUB.

## Memory Usage

Many of BASIC's functions are controlled by a set of tables built in RAM not already occupied by BASIC or the Operating System (OS). Figure 3.1 is a diagram of memory use by both programs. Every time a BASIC programmer enters a statement, memory requirements for the RAM tables change. Memory use by the OS also varies. Different graphics modes, for example, require different amounts of memory.

These changing memory requirements are monitored, and this series of pointers keeps BASIC and the OS from overlaying each other in memory:

- High memory address (HMADR) at location \$02E5
- Application high memory (APHM) at location $\$ 000 \mathrm{E}$
- Low memory address (LMADR) at location \$02E7

When a graphics mode requires larger screen space, the OS checks the application high memory address (APHM) that has been set by BASIC. If there is enough room for the new screen, the OS uses the upper portion of space and sets the pointer HMADR to the bottom of the screen to tell the application how much space the OS is now using.

BASIC builds its table toward high memory from low memory. The pointer to the lowest memory available to an application, called LMADR in the BASIC listing, is set by the OS to tell BASIC the lowest memory address that BASIC can use. When BASIC needs more room for one of its tables, BASIC checks HMADR. If there is enough room, BASIC uses the space and puts the highest address it has used into APHM for OS.

BASIC's operation consists primarily of building, reading, and modifying tables. Pointers to the RAM tables are kept in consecutive locations in zero page starting at $\$ 80$. These tables are, in order,

- Multipurpose Buffer
- Variable Name Table
- Variable Value Table
- String/Array Table


## Chapter Three

- Statement Table
- Runtime Stack

BASIC reserves space for a buffer at LMADR. It then builds the tables contiguously (without gaps), starting at the top of the buffer and extending as far as necessary towards APHM. When a new entry needs to be added to a table, all data in the tables above is moved upward the exact amount needed to fit the new entry into the right place.

Figure 3-1. Memory Usage


Free RAM


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## Variable Name Table

The Variable Name Table (VNT) is built during the pre-compile process. It is read, but not modified, during execution - but only by the LIST statement. The VNT contains the names of the variables used in the program in the order in which they were entered.

The length of entries in the Variable Name Table depends on the length of the variable name. The high order bit of the last character of the name is on. For example, the ATASCII code for the variable name ABC is 414243 (expressed in hexadecimal). In the Variable Name Table it looks like this:

## 4142 C3

The $\$$ character of a string name and the ( character of an array element name are stored as part of the variable name. The table entries for variables C, AA\$, and X(3) would look like this:

| C | C3 |  |  |
| :--- | :--- | :--- | :--- |
| AA\$ | 41 | 41 | A4 |
| X(3) | 58 | A8 |  |

It takes only two bytes to store $X(3)$ because this table stores only X(.

A variable is represented in BASIC by a token. The value of this token is the position (relative to zero) of the variable name in the Variable Name Table, plus $\$ 80$. BASIC references an entry in the table by using the token, minus $\$ 80$, as an index. The Variable Name Table is not changed during execution time.

The zero page pointer to the Variable Name Table is called VNTP in the BASIC listing.

## Variable Value Table

The Variable Value Table (VVT) is also built during the precompile process. It is both read and modified during execution.
There is a one-to-one correspondence in the order of entries between the Variable Name Table and the Variable Value Table. If XXX is the fifth variable in the Variable Name Table, then XXX's value is the fifth entry in the Variable Value Table. BASIC references a table entry by using the variable token, minus $\$ 80$, as an index.

Each entry in the Variable Value Table consists of eight bytes. The first two bytes have the following meaning:

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type = one byte, which indicates the type of variable $\$ 00$ for floating point variable
\$40 for array variable
$\$ 80$ for string variable
vnum $=$ one byte, which indicates the relative position of the variable in the tables

The meaning of the next six bytes varies, depending on the type of variable (floating point, string, or array). In all three cases, these bytes are initialized to zero during syntaxing and during the execution of the RUN or CLR.

When the variable is a floating point number, the six bytes represent its value.

When the variable is an array, the remaining six bytes have the following format:

disp $=$ the two-byte displacement into string/array space of this array variable
$\operatorname{dim} 1=$ two bytes indicating the first dimension value
$\operatorname{dim} 2=$ two bytes indicating the second dimension value
All three of these values are set appropriately when the array is DIMensioned during execution.

When the variable is a string, the remaining six bytes have the following meaning:


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disp = the two-byte displacement into string/array space of this string variable. This value is set when the string is DIMensioned during execution.
curl $=$ the two-byte current length of the string. This value changes as the length of the string changes during execution.
maxl = the two-byte maximum possible length of this string.
This value is set to the DIM value during execution.
When either a string or an array is DIMensioned during execution, the low-order bit in the type byte is turned on, so that the array type is set to $\$ 41$ and the string type to $\$ 81$.

The zero page pointer to the Variable Value Table is called VVTP in the BASIC listing.

## Statement Table

The Statement Table, built as each statement is entered during editing, contains tokenized forms of the statements that were entered. This table determines what happens during execution.

The format of a Statement Table entry is shown in Figure 3-2. There can be several tokens per statement and several statements per line.

Figure 3-2. Format of a Statement Table Entry

$\operatorname{lnum}=$ the two-byte line number (low-order, high-order)
llen $=$ the one-byte line length (the displacement to the next line in the table)
slen $=$ the one-byte statement length (the displacement to the next statement in the line)
snt $=$ the one-byte Statement Name Token
toks $=$ the other tokens that make up the statement (this is variable in length)
eos = the one-byte end of statement token
eol = the one-byte end of line token
The zero page pointer to the Statement Table is called STMTAB in the BASIC listing.

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## String/Array Table

The String/Array Table (also called String/Array Space) is created and modified during execution. Strings and arrays can be intermixed in the table, but they have different formats. Each array or string is pointed to by an entry in the Variable Value Table. The entry in the String/Array Table is created when the string or array is DIMensioned during execution. The data in the entry changes during execution as the value of the string or an element of the array changes.

An entry in the String/Array Table is not initialized to any particular value when it is created. The elements of arrays and the characters in a string cannot be counted upon to have any particular value. They can be zero, but they can also be garbage - data previously stored at those locations.

## Array Entry

For an array, the String/Array Table contains one six-byte entry for each array element. Each element is a floating point number, stored in raveled order. For example, the entry in the String/Array Table for an array that was dimensioned as A(1,2) contains six elements, in this order:

$$
\mathrm{A}(0,0) \quad \mathrm{A}(0,1) \quad \mathrm{A}(0,2) \quad \mathrm{A}(1,0) \quad \mathrm{A}(1,1) \quad \mathrm{A}(1,2)
$$

## String Entry

A string entry in the String/Array Table is created during execution, when the string is DIMensioned. The size of the entry is determined by the DIM value. The "value" of the string to BASIC at any time is determined by the data in the String/Array Table and the current length of the string as set in the Variable Value Table.

The zero page pointer to the String/Array Table is called STARP in the BASIC listing.

The Runtime Stack is created during execution. BASIC uses this LIFO stack to control processing of FOR/NEXT loops and GOSUBs. When either a FOR or a GOSUB statement is encountered during execution, an entry is put on the Runtime Stack. When a NEXT, RETURN, or a POP statement is encountered, entries are pulled off the stack.

Both the FOR entry and the GOSUB entry have a four-byte header:

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type $=$ one byte indicating the type of element
GOSUB type $=0$
FOR type $=$ non-zero
lnum $=$ the two-byte number of the line which contains the statement (low-order, high-order)
disp = one byte indicating the displacement into the line in the Statement Table of the token which caused this stack entry.

The FOR-type byte is actually the token representing the loop control variable from the FOR statement. (In the statement FOR $I=1$ to 10 , I is the loop control variable.) So the FOR-type byte will have a value of $\$ 80$ through $\$ F F$ - the possible values of a variable token.

The FOR entry contains 12 additional bytes, formatted like this:

sval = the six-byte (floating point) limit value at which to stop the loop
step $=$ the six-byte (floating point) STEP value to increment by

The GOSUB entry consists entirely of the four-byte header. The LIST and READ statements also put a GOSUB type entry on the Runtime Stack, so that the line containing the LIST or READ can be found again when the statement has finished executing.

The zero page pointer to the Runtime Stack is called RUNSTK in the BASIC listing.

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## Zero Page Table Pointers

The starting addresses of the tables change dynamically during both program construction and program execution. BASIC keeps the current start addresses of the tables and other pointers required to manage memory space in contiguous zeropage cells. Each pointer is a two-byte address, low byte first.

Since these zero page cell addresses remain constant, BASIC is always able to find the tables. Here are the zero page pointers used in memory management, their names in the BASIC listing, and their addresses:

Multipurpose Buffer
Variable Name Table
VNTP
\$80, \$81
VNT dummy end
Variable Value Table
Statement Table
Current Statement Pointer
String/Array Table
Runtime Stack
Top of used memory
VNTD $\$ 84, \$ 85$

VVTP $\$ 86, \$ 87$
STMTAB $\$ 88, \$ 89$
STMCUR \$8A, \$8B
STARP
RUNSTK
\$8C, \$8D
\$8E, \$8F
MEMTOP \$90,\$91

## Memory Management Routines

Memory Management routines allocate space to the BASIC tables as needed. There are two routines: expand, to add space, and contract, to delete space. Each routine has one entry point for cases in which the number of bytes to be added or deleted is less than 256, and another when it is greater than or equal to 256.

The EXPAND and CONTRACT routines often move many thousands of bytes each time they are called. The 6502 microprocessor is designed to move fewer than 256 bytes of data very quickly. When larger blocks of data are moved, the additional 6502 instructions required can make the process very slow. The EXPAND and CONTRACT routines circumvent this by using the less-than-256-byte fast-move capabilities in the movement of thousands of bytes. The end result is a set of very fast and very complex data movement routines.

All of this complexity does have a drawback. The infamous Atari BASIC lock-up problem lives in these two routines. If an EXPAND or CONTRACT requires that an exact multiple of 256 bytes be moved, then the routines move things from the wrong

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place in memory to the wrong place in memory, whereupon the computer locks up and won't respond. The only way to avoid losing hours of work this way is to SAVE to disk or cassette frequently.

## EXPAND (\$A881)

Parameters at entry:

| register the zero page address containing the pointer to <br> $X$ the location after which space is to be added |  |
| ---: | :--- |
| $\mathrm{Y}=$ | the low-order part of the number of bytes to |
|  | expand |
| $\mathrm{A}=$ | the high-order part of the number of bytes to |
|  | expand |

The routine creates a hole in the table memory, starting at a requested location and continuing the requested number of bytes.

The routine first checks to see that there is enough free memory space to satisfy the request.

It adds the requested expand size to each of the zero-page table pointers between the one pointed to by the $X$ register and MEMTOP. Then each pointer will point to the correct address when EXPAND is done.

EXPAND then creates space at the address indicated by the $X$ register. The number of bytes required is contained in the $Y$ and A registers. (Y contains the least significant byte, while A contains the most significant.) All data from the requested address to the address pointed to by MEMTOP is moved toward high memory by the requested number of bytes. This creates a hole of the proper size.

The routine then sets Application High Memory (APHM) to the value in MEMTOP. This tells the OS the highest memory address that BASIC is currently using.

## EXPLOW (\$A87F)

Parameters at entry:

```
X register = zero page address containing the pointer to the
        location after which space is to be added
    Y = number of bytes to expand (low-order byte only)
```


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This is an additional entry point for the EXPAND routine. It is used when the number of bytes to be added to the table is less than 256.

This routine first loads the 6502 accumulator with zero to indicate the most significant byte of the expand length. It then functions exactly like EXPAND.

## CONTRACT (\$A8FD)

Parameters at entry:

$$
\begin{aligned}
& \text { register } \\
& \text { X = zero page address containing the pointer to the } \\
& \text { starting location where space is to be removed } \\
& \mathrm{Y}=\text { the low-order part of the number of bytes to } \\
& \text { contract } \\
& \mathrm{A}=\text { the high-order part of the number of bytes to } \\
& \text { contract }
\end{aligned}
$$

This routine removes a requested number of bytes at a requested location by moving all the data from higher in the tables downward the exact amount needed to replace the unwanted bytes.

It subtracts the requested contract size from each of the zero page table pointers between the one pointed to by the $X$ register and MEMTOP. Then each pointer will point to the correct address when CONTRACT is done.

The routine sets application high memory (APHM) to the value in MEMTOP to indicate to the OS the highest memory address that BASIC is currently using.

The block of data to be moved downward is defined by starting at the address pointed to by the zero-page address pointed to in X, plus the offset number stored in Y and A , and then continuing to the address specified at MEMTOP. Each byte of data in that block is moved downward in memory by the number of bytes specified in Y and A , effectively erasing all the data between the specified address and that address plus the requested offset.
CONTLOW (\$A8FB)
Parameters at entry:

[^0]\[

$$
\begin{aligned}
\mathrm{Y} & =\begin{array}{l}
\text { the number of bytes to contract (low-order byte } \\
\text { only) }
\end{array}
\end{aligned}
$$
\]

This routine is used to remove fewer than 256 bytes from the tables at a requested location by moving all the data from higher in the tables downward the exact amount needed to replace the unwanted bytes.

This routine first loads the 6502 accumulator with zero to serve as the most significant byte of the contract length. It then functions exactly like CONTRACT.

## Miscellaneous Memory Allocations

Besides the tables, which change dynamically, BASIC also uses buffers and stacks at fixed locations.

The Argument/Operator Stack is allocated at BASIC's low memory address and occupies 256 bytes. During pre-compiling it is used as the output buffer for the tokens. During execution, it is used while evaluating an expression. This buffer/stack is referenced by a pointer at location $\$ 80$. This pointer has several names in the BASIC listing: LOMEM, ARGOPS, ARGSTK, and OUTBUFF.

The Syntax Stack is used during the process of syntaxing a statement. It is referenced directly - that is, not through a pointer. It is located at $\$ 480$ and is 256 bytes long.

The Line Buffer is the storage area where the statement is placed when it is ENTERed. It is the input buffer for the edit and pre-compile processes. It is 128 bytes long and is referenced directly as LBUFF. Often the address of LBUFF is also put into INBUFF so that the buffer can be referenced through a pointer, though INBUFF can point to other locations during various phases of BASIC's execution.

## Program Editor

The Atari keyboard is the master control panel for Atari BASIC. Everything BASIC does has its origins at this control panel. The Program Editor's job is to service the control panel and respond to the commands that come from it.

The editor gets a line from the user at the keyboard; does some preliminary processing on the line; passes the line to the pre-compiler for further processing; inserts, deletes, or replaces the line in the Statement Table; calls the Program Executor when necessary; and then waits to receive the user's next line input.

## Line Processing

The Program Editor, which starts at \$A060, begins its process by resetting the 6502 CPU stack. Resetting the CPU stack is a drastic operation that can only occur at the beginning of a logical process. Each time Atari BASIC prepares to get a new line from the user, it restarts its entire logical process.

## Getting a Line

The Program Editor gets a user's line by calling CIO. The origin of the line is transparent to the Program Editor. The line may have been typed in at the keyboard or entered from some external device like the disk (if the ENTER command was given). The Program Editor simply calls CIO and asks it to put a line of not more than 255 bytes into the buffer pointed to by INBUFF (\$F3). INBUFF points to the 128 -byte area defined at LBUFF (\$580).

The OS's screen editor, which is involved in getting a line from the keyboard, will not pass BASIC a line that is longer than 120 bytes. Normally, then, the 128-byte buffer at LBUFF is big enough to contain the user's line.

Sometimes, however, if a line was originally entered from the keyboard with few blanks and many abbreviations, then LISTed to and re-ENTERed from the disk, an input line may be longer than 128 bytes. When this happens, data in the $\$ 600$ page is overlaid. A LINE TOO LONG error will not necessarily

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occur at this point. A LINE TOO LONG error occurs only if the Pre-compiler exceeds its stack while processing the line or if the tokenized line OUTBUFF exceeds 256 bytes. These overflows depend on the complexity of the line rather than on its actual length.

When CIO has put a line into the line buffer (LBUFF) and the Program Editor has regained control, it checks to see if the user has changed his mind and hit the break key. If the user did indeed hit break, the Program Editor starts over and asks CIO for another line.

## Flags and Indices

In order to help control its processing, the Program Editor uses flags and indices. These must be given initial values.
CIX and COX. The index CIX (\$F2) is used to access the user's input line in the line buffer (LBUFF), while COX (\$94) is used to access the tokenized statement in the output buffer (OUTBUFF). These buffers and their indices are also used by the pre-compiler. The indices are initialized to zero to indicate the beginning of the buffers.
DIRFLG. This flag byte (\$A6) is used by the editor to remember whether a line did or did not have a line number, and also to remember if the pre-compiler found an error in that line. DIRFLG is initialized to zero to indicate that the line has a line number and that the pre-compiler has not found an error.
MAXCIX. This byte ( $\$ 9 \mathrm{~F}$ ) is maintained in case the line contains a syntax error. It indicates the displacement into LBUFF of the error. The character at this location will then be displayed in inverse video. The Program Editor gives this byte the same initial value as CIX, which is zero.
SVVNTP. The pointer to the current top of the Variable Name Table (VNTD) is saved as SVVNTP (\$AD) so that if there is a syntax error in this line, any variables that were added can be removed. If a user entered an erroneous line, such as 100 $\mathrm{A}=$ XAND B, the variable XAND would already have been added to the variable tables before the syntax error was discovered. The user probably meant to enter $100 \mathrm{~A}=\mathrm{X}$ AND B, and, since there can only be 128 variables in BASIC, he probably does not want the variable XAND using up a place in the variable tables. The Program Editor uses SVVNTP to find the entry in the Variable Name Table so it can be removed.

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SVVVTE. The process used to indicate which variable entries to remove from the Variable Value Table in case of error is different. The number of new variables in the line (SVVVTE,\$B1) is initialized to zero. The Program Pre-compiler increments the value every time it adds a variable to the Variable Value Table. If a syntax error is detected, this number is multiplied by eight (the number of bytes in each entry on the Variable Value Table) to get the number of bytes to remove, counting backward from the most recent value entered.

## Handling Blanks

In many places in the BASIC language, blanks are not significant. For example,

$$
100 \text { IFX = 6THENGOTO500 }
$$

has the same meaning as
100 IF X $=6$ THEN GOTO 500.
The Program Editor, using the SKIPBLANK routine (\$DBA1), skips over unnecessary blanks.

## Processing the Line Number

Once the editor has skipped over any leading blanks, it begins to examine the input line, starting with the line number. The floating point package is called to determine if a line number is present, and, if so, to convert the ATASCII line number to a floating point number. The floating point number is converted to an integer, saved in TSLNUM for later use, and stored in the tokenized line in the output buffer (OUTBUFF).

The routine used to store data into OUTBUFF is called :SETCODE (\$A2C8). When :SETCODE stores a byte into OUTBUFF, it also increments COX, that buffer's index.

BASIC could convert the ATASCII line number directly to an integer, but the routine to do this would not be used any other time. Routines to convert ATASCII to floating point and floating point to integer already exist in BASIC for other purposes. Using these existing routines conserves ROM space.

An interesting result of this sequence is that it is valid to enter a floating point number as a line number. For example, $100.1,10.9$, or 2.05 E 2 are valid line numbers. They would be converted to 100, 11, and 205 respectively.

If the input line does not start with a line number, the line is considered to be a direct statement. DIRFLG is set to $\$ 80$ so
that the editor can remember this fact. The line number is set to 32768 (\$8000). This is one larger than the largest line number a user is allowed to enter. BASIC later makes use of this fact in processing the direct statement.
Line length. The byte after the line number in the tokenized line in OUTBUFF is reserved so that the line length (actually the displacement to the next line) can be inserted later. (See Chapter 2.) The routine :SETCODE is called to reserve the byte by incrementing (COX) to indicate the next byte.
Saving erroneous lines. In the byte labeled STMSTRT, the Program Editor saves the index into the line buffer (LBUFF) of the first non-blank character after the line number. This index is used only if there is a syntax error, so that all the characters in the erroneous line can be moved into the tokenized line buffer and from there into the Statement Table.

There are advantages to saving an erroneous line in the Statement Table, because you can LIST the error line later. The advantage is greatest, not when entering a program at the keyboard, but when entering a program originally written in a different BASIC on another machine (via a modem, perhaps). Then, when a line that is not correct in Atari BASIC is entered, the line is flagged and stored - not discarded. The user can later list the program, find the error lines, and re-enter them with the correct syntax for Atari BASIC.
Deleting lines. If the input line consists solely of a line number, the Program Editor deletes the line in the Statement Table which has that line number. The deletion is done by pointing to the line in the Statement Table, getting its length, and calling CONTRACT. (See Chapter 3.)

## Statement Processing

The user's input line may consist of one or more statements. The Program Editor repeats a specific set of functions for each statement in the line.

## Initializing

The current index (COX) into the output buffer (OUTBUFF) is saved in a byte called STMLBD. A byte is reserved in OUTBUFF by the routine :SETCODE. Later, the value in

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STMLBD will be used to access this byte, and the statement length (the displacement to the next statement) will be stored here.

## Recognizing the Statement Name

After the editor calls SKBLANK to skip blanks, it processes the statement name, now pointed to by the input index (CIX). The editor calls the routine SEARCH (\$A462) to look for this statement name in the Statement Name Table. SEARCH saves the table entry number of this statement name into location STENUM.

The entry number is also the Statement Name Token value, and it is stored into the tokenized output buffer (OUTBUFF) as such by :SETCODE. The SEARCH routine also saves the address of the entry in SRCADR for use by the pre-compiler.

If the first word in the statement was not found in the Statement Name Table, the editor assumes that the statement is an implied LET, and the appropriate token is stored. It is left to the pre-compiler to determine if the statement has the correct syntax for LET.

The editor now gives control to the pre-compiler, which places the appropriate tokens in OUTBUFF, increments the indices CIX and COX to show current locations, and indicates whether a syntax error was detected by setting the 6502 carry flag on if there was an error and clearing the carry flag if there was not. (See Chapter 5.)

## If a Syntax Error Is Detected

If the 6502 carry flag is set when the editor regains control, the editor does error processing.

In MAXCIX, the pre-compiler stored the displacement into LBUFF at which it detected the error. The Program Editor changes the character at this location to inverse video.

The character in inverse video may not be the point of error from your point of view, but it is where the pre-compiler detected an error. For example, assume you entered $X=$ YAND Z . You probably meant to enter $\mathrm{X}=\mathrm{Y}$ AND Z , and therefore would consider the error to be between Y and AND. However, since YAND is a valid variable name, $\mathrm{X}=\mathrm{YAND}$ is a valid BASIC statement.

The pre-compiler doesn't know there is an error until it encounters $B$. The value of highlighting the error with inverse
video is that it gives the user an approximation of where the error is. This can be a big advantage, especially if the input line contained multiple statements or complex expressions.

The next thing the editor does when a syntax error has been detected is set a value in DIRFLG to indicate this fact for future reference. Since the DIRFLG byte also indicates whether this is a direct statement, the error indicator of $\$ 40$ is ORed with the value already in DIRFLG.

The editor takes the value that it saved in STMSTRT and puts it into CIX so that CIX now points to the start of the first statement in the input line in LBUFF. STMLBD is set to indicate the location of the first statement length byte in OUTBUFF. (A length will be stored into OUTBUFF at this displacement at a later time.)

The editor sets the index into OUTBUFF (COX) to indicate the Statement Name Token of the first statement in OUTBUFF, and stores a token at that location to indicate that this line has a syntax error. The entire line (after the line number) is moved into OUTBUFF. At this point COX indicates the end of the line in OUTBUFF. (Later, the contents of OUTBUFF will be moved to the Statement Table.)

This is the end of the special processing for an erroneous line. The process that follows is done for both correct and erroneous lines.

## Final Statement Processing

During initial line processing, the Program Editor saved in STMLBD a value that represents the location in OUTBUFF at which the statement length (displacement to the next statement) should be stored. The Program Editor now retrieves that value from STMLBD. Using this value as an index, the editor stores the value from COX in OUTBUFF as the displacement to the next statement.

The Program Editor checks the next character in LBUFF. If this character is not a carriage return (indicating end of the line), then the statement processing is repeated. When the carriage return is found, COX will be the displacement to the next line. The Program Editor stores COX as the line length at a displacement of two into OUTBUFF.

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## Statement Table Processing

The final tokenized form of the line exists in OUTBUFF at this point. The Program Editor's next task is to insert or replace the line in the Statement Table.

The Program Editor first needs to create the correct size hole in the Statement Table. The editor calls the GETSTMT routine (\$A9A2) to find the address where this line should go in the Statement Table. If a line with the same line number already exists, the routine returns with the address in STMCUR and with the 6502 carry flag off. Otherwise, the routine puts the address where the new line should be inserted in the Statement Table into STMCUR and turns on the 6502 carry flag. (See Chapter 6.)

If the line does not exist in the Statement Table, the editor loads zero into the 6502 accumulator. If the line does exist, the editor calls the GETLL routine (\$A9DD) to put the line length into the accumulator. The editor then compares the length of the line already in the Statement Table (old line) with the length of the line in OUTBUFF (new line).

If more room is needed in the Statement Table, the editor calls the EXPLOW (\$A87F; see Chapter 3). If less space is needed for the new line, it calls a routine to point to the next line (GNXTL, at location \$A9D0; see Chapter 6), and then calls the CONTLOW (\$A8FB; see Chapter 3).

Now that we have the right size hole, the tokenized line is moved from OUTBUFF into the Statement Table at the location indicated by STMCUR.

## Line Wrap-up

After the line has been added to the Statement Table, the editor checks DIRFLG for the syntax error indicator. If the second most significant bit (\$40) is on, then there is an error.

## Error Wrap-up

If there is an error, the editor removes any variables that were added by this line by getting the number of bytes that were added to the Variable Name Table and the Variable Value Table from SVVNTP and SVVVTE. It then calls CONTRACT (\$A8FD) to remove the bytes from each table.

Next, the editor lists the line. The Statement Name Token, which was set to indicate an error, causes the word "ERROR"'

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to be printed. An inverse video character indicates where the error was detected. The editor now waits for you to enter another line.

## Handling Correct Lines

If the line was syntactically correct, the editor again examines DIRFLG. In earlier processing, the most significant bit (\$80) of this byte was set on if the line was a direct statement. If it is not a direct statement, then the editor is finished with the line, and it waits for another input line.

If the line is a direct statement, earlier processing already assigned it a line number of 32768 (\$8000), one larger than the largest line number a user can enter. Since lines are arranged in the Statement Table in ascending numerical order, this line will have been inserted at the end of the table. The current statement pointer (STMCUR-\$8A, \$8B) points to this line.

The Program Editor transfers control to a Program Executor routine, Execution Control (EXECNL at location \$A95F), which will handle the execution of the direct statement. (See Chapter 6.)

## The Pre-compiler

The symbols and symbol-combining rules of Atari BASIC are coded into Syntax Tables, which direct the Program Precompiler in examining source code and producing tokens. The information in the Syntax Tables is a transcription of a metalanguage definition of Atari BASIC.

## The Atari BASIC Meta-language

A meta-language is a language which describes or defines another language. Since a meta-language is itself a language, it also has symbols and symbol-combining rules - which define with precision the symbols and symbol-combining rules of the subject language.

Atari BASIC is precisely defined with a specially developed meta-language called the Atari BASIC Meta-language, or ABML. (ABML was derived from a commonly used compilertechnology meta-language called BNF.) The symbols and symbol-combining rules of ABML were intentionally kept very simple.

## Making Up a Language

To show you how ABML works, we'll create an extremely simple language called SAP, for Simple Arithmetic Process. SAP symbols consist of variables, constants, and operators.

- Variables: The letters A, B, and C only.
- Constants: The numbers $1,2,3,4,5,6,7,8$, and 9 only.
- Operators: The characters $+,-, *, l$, and ! only. Of course, you already know the functions of all the operators except "!". The character! is a pseudo-operator of the SAP language used to denote the end of the expression, like the period that ends this sentence.
The grammar of the SAP language is precisely defined by the ABML definition in Figure 5-1.


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Figure 5-1. The SAP Language Expressed in ABML

The ABML symbols used to define the SAP language in Figure 5-1 are:
$:=$ is defined as Whatever is on the left of := is defined as consisting of whatever is on the right of :=, and in that order.
The symbol | allows choices for what something is defined as. For instance, in the sixth line < variable> can be A or B or C. If | does not appear between two symbols, then there is no choice. For example, in the second line <expression > must have both <value> and <operation>, in that order, to be valid.
$<>$ label Whatever comes between $<$ and $>$ is an ABML label. All labels, as non-terminal symbols, must be defined at some point, though the definitions can be circular notice that <operation> is part of the definition of <expression> in the second line, while in the third line < expression > is part of the definition of <operation>.
terminal Symbols used in definitions, which are not symbols enclosed by $<$ and $>$ and are also not one of the ABML symbols, are terminal symbols in the language being defined by ABML. In SAP, some terminal symbols are A, !, B, *, and 1. They cannot be defined as consisting of other symbols - they are themselves the symbols that the SAP language manipu-
lates, and must appear exactly as they are shown to be valid in SAP. In effect, they are the vocabulary of the SAP language.

## Statement Generation

The ABML description of SAP can be used to generate grammatically correct statements in the SAP language. To do this, we merely start with the first line of the definition and replace the non-terminal symbols with the definitions of those symbols. The replacement continues until only terminal symbols remain. These remaining terminal symbols constitute a grammatically correct SAP statement.

Since the or statement requires that one and only one of the choices be used, we will have to arbitrarily replace the nonterminal with the one valid choice.

Figure 5-2 illustrates the ABML statement generation process.

## Figure 5-2. The Generation of One Possible SAP Statement

(1) SAP := <expression>!
(2) SAP := < value ><operation>!
(3) SAP := <variable><operation >!
(4) $\mathrm{SAP}:=\mathrm{B}<$ operation $>$ !
(5) SAP := $\mathrm{B}<$ operator $><$ expression $>$ !
(6) SAP $:=\mathrm{B}^{*}<$ expression $>$ !
(7) SAP := $\mathrm{B}^{*}<$ value $><$ operation $>$ !
(8) SAP $:=\mathrm{B}^{*}<$ constant $><$ operation $>$ !
(9) SAP $:=\mathrm{B}^{*} 4<$ operation $>$ !
(10) SAP $:=\mathrm{B}^{*} 4<$ operator $><$ expression $>$ !
(11) SAP $:=\mathrm{B}^{*} 4+<$ expression $>$ !
(12) SAP : $=\mathrm{B}^{*} 4+<$ value $><$ operation $>$ !
(13) SAP := $\mathrm{B}^{*} 4+<$ variable $><$ operation $>$ !
(14) SAP := $\mathrm{B}^{*} 4+\mathrm{C}<$ operation $>$ !
(15) SAP $:=B^{*} 4+C$ !

In (2), < value > < operation > replaces < expression > because the ABML definition of SAP (Figure 5-1) defines <expression> as < value> <operation>.

In (3), the non-terminal <value $>$ is replaced with
<variable>. The definition of < value> gives two choices for the substitution of $<$ value $>$. We happened to choose $<$ variable>.

In (4), we reach a terminal symbol, and the process of defining <value> ends. We happened to choose $\boldsymbol{B}$ to replace $<$ variable $>$.

In (5), we go back and start defining <operation>. There are two choices for the replacement of <operation>, either <operator> <expression> or nothing at all (since there is nothing to the right of $\mid$ in the second line of Figure 5-1). If nothing had been chosen, then (5) would have been: SAP :=B! The statement $B$ ! has no further non-terminals; the process would have been finished, and a valid statement would have been produced. Instead we happened to choose <operator> <expression>.

The SAP definition for < expression > is <value> <operation>. If we replace <operation> with its definition we get:
<expression> :=<value><operator> <expression>

The definition of <expression> includes <expression> as part of its definition. If the <operator> <expression> choice were always made for <operation> , then the process of replacement would never stop. A SAP statement can be infinitely long by definition. The only thing which prevents us from always having an infinitely long SAP statement is that there is a second choice for the replacement of <operation> : nothing.

The replacements in (5) and (10) reflect the repetitive choices of defining <expression > in terms of itself. The choice in (15) reflects the nothing choice and thus finishes the replacement process.

## Computerized Statement Generation

If we slightly modify our procedure for generating statements, we will have a process that could be easily programmed into a computer. Instead of arbitrarily replacing the definition of nonterminals, we can think of the non-terminal as a GOSUB. When we see $\langle\mathrm{X}\rangle:=\langle\mathrm{Y}\rangle\langle\mathrm{Z}\rangle$, we can think of $\langle\mathrm{Y}\rangle$ as being a subroutine-type procedure:
(a) Go to the line that has $\langle Y\rangle$ on the left side.
(b) Process the definition (right side) of $\langle\mathrm{Y}\rangle$.

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(c) If while processing the definition of $\langle\mathrm{Y}\rangle$, other nonterminals are found, GOSUB to them.
(d) If while processing the definition of $\langle Y\rangle$ we encounter a terminal, output the terminal symbol as the next symbol of the generated statement.
(e) When the definition of $\langle Y\rangle$ is finished, return to the place that $\langle Y\rangle$ was called from and continue.
Since ABML is structured so that it can be programmed, a fascinating exercise is to design a simple English sentence grammar with ABML, then write a BASIC program to generate valid English sentences at random. The randomness of the sentences would be derived by using the RND function to select from the definitions or choices. An example of such a grammar is shown in Figure 5-3. (The programming exercise is left to you.)

## Figure 5-3. A Simple English Sentence Grammar in ABML

$$
\begin{aligned}
& \text { SENTENCE }:=\text { <subject }><\text { adverb }><\text { verb }>\text { <object }>. \\
& \text { < subject }>:=\text { The <adjective } \ll \text { noun }> \\
&<\text { verb }>:=\text { eats } \mid \text { sleeps } \mid \text { drinks } \mid \text { talks } \mid \text { hugs } \\
& \text { <adverb> }>:=\text { quickly } \mid \text { silently } \mid \text { slowly } \mid \text { viciously } \mid \\
& \text { lovingly } \mid \text { sadly } \mid
\end{aligned}
$$

## Syntactical Analysis

The process of examining a language statement for grammatical correctness is called syntactical analysis, or syntaxing.

Statement verification is similar to statement generation. Instead of arbitrarily choosing which or definition to use, however, the choices are already made, and we must check to see whether the statement symbols are used in valid patterns. To do this, we must process through each or definition until we find a matching valid terminal symbol.

The result of statement generation is a valid, grammatically correct statement, but the result of statement verification is a

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statement validity indication, which is a simple yes or no. Either the statement is grammatically correct or it is not. Failure occurs when some statement symbol cannot be matched with a valid terminal symbol under the rules of the grammar.

## The Reporting System

To use the pass/fail result of statement verification, we must build a reporting system into the non-terminal checking process. Whenever we, in effect, GOSUB to a non-terminal definition, that non-terminal definition must report its pass/fail status.

A fail status is generated and returned by a non-terminal definition when it finds no matching terminal for the current statement symbol. If the current statement symbol is B and the <constant> definition in the SAP language is called, then <constant> would report a fail status to the routine that called it.

A pass status is returned when a terminal symbol is found which matches the current statement symbol. If our current statement symbol had been 7 instead of B, then $<$ constant > would have reported pass.

Whenever such a match does occur, we return to the statement, and the next symbol to the right becomes the new current symbol for examination and verification.

## Cycling Through the Definitions

In SAP, the <constant > definition is called from the < value > definition. If <constant> reports fail, then we examine the next or choice, which is <variable>. The current symbol is B, so < variable> reports pass.

Since at least one of the or choices of < value $>$ has reported pass, <value> will report pass to its caller. If both <constant> and < variable> had reported fail, then <value> would report fail to its caller.

The caller of <value> is <expression> . If < value> reports pass, <operation> is called. If <operation> reports pass, then <expression > can report pass to its caller. If either <value> or <operation> reports fail, then <expression> must report fail, since there are no other or choices for <expression>.

The definition of <operation> contains a special pass/fail property. If either <operator> or <expression> reports fail,

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then the or choice must be examined. In this case the or choice is nothing. The or nothing means something special: report pass, but do not advance to the next symbol.

The final pass/fail report is generated from the first line of the definition. If <expression > reports pass and the next symbol is !, then SAP reports pass. If either one of these conditions has a fail status, then SAP must report fail to whatever called SAP from outside the language.

## Backing Up

Sometimes it is necessary to back up over symbols which have already been processed. Let's assume that there was a definition of the type $\langle X\rangle:=\langle Y\rangle \mid\langle Z\rangle$. It is possible that while $\langle Y\rangle$ is attempting to complete its definition, it will find a number of valid matching terminal symbols before it discovers a symbol that it cannot match. In this case, $\langle Y\rangle$ would have consumed a number of symbols before it decided to report fail. All of the symbols that $\langle\mathrm{Y}\rangle$ consumed must be unconsumed before $\langle\mathrm{Z}\rangle$ can be called, since $\langle\mathrm{Z}\rangle$ will need to check those same symbols.

The process of unconsuming symbols is called backup. Backup is usually performed by the caller of $\langle Y\rangle$, which remembers which source symbol was current when it called $<\mathrm{Y}\rangle$. If $\langle\mathrm{Y}\rangle$ reports fail, then the caller of $\langle\mathrm{Y}\rangle$ restores the current symbol pointer before calling $\langle\mathrm{Z}\rangle$.

## Locating Syntax Error

When a final report of fail is given for a statement, it is often possible to guess where the error occurred. In a left-to-right system, the symbol causing the failure is usually the symbol which follows the rightmost symbol found to be valid. If we keep track of the rightmost valid symbol during the various backups, we can report a best guess as to where the failurecausing error is located. This is exactly what Atari BASIC does with the inverse video character in the ERROR line.

For simplicity, our example was coded for SAP, but the syntactical analysis we have just described is essentially the process that the Atari BASIC pre-compiler uses to verify the grammar of a source statement. The Syntax Tables are an ABML description of Atari BASIC. The pre-compiler, also known as the syntaxer, contains the routines which verify BASIC statements.

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## Statement Syntax Tables

There is one entry in the Syntax Tables for each BASIC statement. Each statement entry in the Syntax Table is a transcription of an ABML definition of the grammar for that particular statement. The starting address of the table entry for a particular statement is pointed to by that statement's entry in the Statement Name Table.

The data in the Syntax Tables is very much like a computer machine language. The pseudo-computer which executes this pseudo-machine language is the pre-compiler code. Like any machine language, the pseudo-machine language of the Syntax Tables has instructions and instruction operands. For example, an ABML non-terminal symbol is transcribed to a code which the pre-compiler executes as a type of "GOSUB and report pass/fail"' command.

Here are the pseudo-instruction codes in the Syntax Tables; each is one byte in length.

## Absolute Non-Terminal Vector

Name: ANTV
Code: \$00
This is one of the forms of the non-terminal GOSUB. It is followed by the address, minus 1 , of the non-terminal's definition within the Syntax Table. The address is two bytes long, with the least significant byte first.

## External Subroutine Call

Name: ESRT
Code: \$01
This instruction is a special type of terminal symbol checker. It is followed by the address, minus 1 , of a 6502 machine language routine. The address is two bytes long, with the least significant byte first. The ESRT instruction is a deus ex machina - the "god from the machine" who solved everybody's problems at the end of classical Greek plays. There are some terminals whose definition in ABML would be very complex and require a great many instructions to describe. In these cases, we go outside the pseudo-machine language of the Syntax Tables and get help from 6502 machine language routines - the deus ex machina that quickly gives the desired

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result. A numeric constant is one example of where this outside help is required.

## ABML or

Name: OR
Value: \$02
This is the familiar ABML or symbol ( | ). It provides for an alternative definition of a non-terminal.

## Return

Name: RTN
Value: \$03
This code signals the end of an ABML definition line. When we write an ABML statement on paper, the end of a definition line is obvious - there is no further writing on the line. When ABML is transcribed to machine codes, the definitions are all pushed up against each other. Since the function that is performed at the end of a definition is a return, the end of definition is called return (RTN).
Unused (Codes \$04 through \$0D are unused.)

## Expression Non-Terminal Vector

Name: VEXP
Value: \$0E
The ABML definition for an Atari BASIC expression is located at \$A60D. Nearly every BASIC statement definition contains the possibility of having <expression> as part of it. VEXP is a single-byte call to <expression > , to avoid wasting the two extra bytes that ANTV would take. The pseudomachine understands that this instruction is the same as an ANTV call to <expression> at SA60D.

## Change Last Token

Name: CHNG
Value: \$0F
This instruction is followed by a one-byte change to token value. The operator token instructions cause a token to be placed into the output buffer. Sometimes it is necessary to change the token that was just produced. For example, there are several $=$ operators. One $=$ operator is for the assignment

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statement LET X=4. Another = operator is for comparison operations like IF $Y=5$. The pseudo-machine will generate the assignment $=$ token when it matches $=$. The context of the grammar at that point may have required a comparison $=$ token. The CHNG instruction rectifies this problem.

## Operator Token

Name: (many)
Value: $\$ 10$ through \$7F
These instructions are terminal codes for the Atari BASIC Operators. The code values are the values of each operator token. The values, value names, and operator symbols are defined in the Operator Name Table (see Chapter 2).

When the pseudo-machine sees these terminal symbol representations, it compares the symbol it represents to the current symbol in the source statement. If the symbols do not match, then fail status is generated. If the symbols match, then pass status is generated, the token (instruction value) is placed in the token output buffer, and the next statement source symbol becomes the current symbol for verification.

## Relative Non-Terminal Vectors

Name: (none)
Value: $\$ 80-\$$ BF (Plus)
$\$ \mathrm{CO}-\$ F F$ (Minus)
This instruction is similar to ANTV, except that it is a single byte. The upper bit is enough to signal that this one-byte code is a non-terminal GOSUB. The destination address of the GOSUB is given as a position relative to the current table location. The values $\$ 80$ through $\$ \mathrm{BF}$ correspond to an address which is at the current table address plus $\$ 00$ through $\$ 3 \mathrm{~F}$. The values $\$ C 0$ through $\$ F F$ correspond to an address which is at the current table address minus $\$ 01$ through $\$ 3 \mathrm{~F}$.

## Pre-compiler Main Code Description

The pre-compiler, which starts at SYNENT (\$A1C3), uses the pseudo-instructions in the Syntax Tables to verify the correctness of the source line and to generate the tokenized statements.

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## Syntax Stack

The pre-compiler uses a LIFO stack in its processing. Each time a non-terminal vector ("GOSUB'") is executed, the precompiler must remember where the call was made from. It must also remember the current locations in the input buffer (source statement) and the output buffer (tokenized statement) in case the called routine reports fail and backup is required. This LIFO stack is called the Syntax Stack.

The Syntax Stack starts at $\$ 480$ at the label SIX. The stack is 256 bytes in size. Each entry in the stack is four bytes long. The stack can hold 64 levels of non-terminal calls. If a sixty-fifth stack entry is attempted, the LINE TOO LONG error is reported. (This error should be called LINE TOO COMPLEX, but the line is most likely too long also.)

The first byte of each stack entry is the current input index (CIX). The second byte is the current output index (COX). The final two bytes are the current address within the syntax tables.

The current stack level is managed by the STKLVL (\$A9) cell. STKLVL maintains a value from $\$ 00$ to $\$ F C$, which is the displacement to the current top of the stack entry.

## Initialization

The editor has saved an address in SRCADR (\$96). This address is the address, minus 1 , of the current statement's ABML instructions in the Syntax Tables. The current input index (CIX) and the current output index (COX) are also preset by the editor.

The initialization code resets the syntax stack manager (STKLVL) to zero and loads the first stack entry with the values in CIX, COX, and CPC - the current program counter, which holds the address of the next pseudo-instruction in the Syntax Tables.

## PUSH

Values are placed on the stack by the PUSH routine (\$A228). PUSH is entered with the new current pseudo-program counter value on the CPU stack. PUSH saves the current CIX, COX, and CPC on the syntax stack and increments STKLVL. Next, it sets a new CPC value from the data on the CPU stack. Finally, PUSH goes to NEXT.

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## POP

Values are removed from the stack with the POP routine (\$A252). POP is entered with the 6502 carry flag indicating pass/fail. If the carry is clear, then pass is indicated. If the carry is set, then fail is indicated.

POP first checks STKLVL. If the current value is zero, then the pre-compiler is done. In this case, POP returns to the editor via RTS. The carry bit status informs the editor of the pass/fail status.

If STKLVL is not zero, POP decrements STKLVL.
At this point, POP examines the carry bit status. If the carry is clear (pass), POP goes to NEXT. If the carry is set (fail), POP goes to FAIL.

## NEXT and the Processes It Calls

After initialization is finished and after each Syntax Table instruction is processed, NEXT is entered to process the next syntax instruction.

NEXT starts by calling NXSC to increment CPC and get the next syntax instruction into the A register. The instruction value is then tested to determine which syntax instruction code it is and where to go to process it.

If the Syntax Instruction is OR (\$02) or RTN (\$03), then exit is via POP. When POP is called due to these two instructions, the carry bit is always clear, indicating pass.
ERNTV. If the instruction is RNTV ("GOSUB'" $\$ 80-\$ F F$ ), then ERNTV (\$A201) is entered. This code calculates the new CPC value, then exits via PUSH.
GETADR. If the instruction is ANTV ( $\$ 00$ ) or the deus ex machina ESRT (\$01) instruction, then GETADR is called. GETADR obtains the following two-byte address from the Syntax Table.

If the instruction was ANTV, then GETADR exits via PUSH.

If the instruction was ESRT, then GETADR calls the external routine indicated. The external routine will report pass/fail via the carry bit. The pass/fail condition is examined at $\$ \mathrm{~A} 1 \mathrm{~F} 0$. If pass is indicated, then NEXT is entered. If fail is indicated, then FAIL is entered.
TERMTST. If the instruction is VEXP (\$0E), then the code at \$A1F9 will go to TERMTST (\$A2A9), which will cause the code

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at \$A2AF to be executed for VEXP. This code obtains the address, minus 1, of the ABML for the <expression > in the Syntax Table and exits via PUSH.
ECHNG. If the instruction was CHNG (\$0F), then ECHNG (\$A2BA) is entered via tests at \$A1F9 and \$A2AB. ECHNG will increment CPC and obtain the change-to token which will then replace the last previously generated token in OUTBUFF. ECHNG exits via RTS, which will take control back to NEXT.
SRCONT. The Operator Token Instructions ( $\$ 10-\$ 7 \mathrm{~F}$ ) are handled by the SRCONT routine. SRCONT is called via tests at \$A1F9 and \$A2AD. SRCONT will examine the current source symbol to see if it matches the symbol represented by the operator token. When SRCONT has made its determination, it will return to the code at \$A1FC. This code will examine the pass/fail (carry clear/set) indicator returned by SRCONT and take the appropriate action. (The SRCONT routine is detailed on the next page.)

FAIL
If any routine returns a fail indicator, the FAIL code at \$A26C will be entered. FAIL will sequentially examine the instructions, starting at the Syntax Table address pointed to by CPC, looking for an OR instruction.

If an OR instruction is found, the code at \$A27D will be entered. This code first determines if the current statement symbol is the rightmost source symbol to be examined thus far. If it is, it will update MAXCIX. The editor will use MAXCIX to set the inverse video flag if the statement is erroneous. Second, the code restores CIX and COX to their before-failure values and goes to NEXT to try this new OR choice.

If, while searching for an OR instruction, FAIL finds a RTN instruction, it will call POP with the carry set. Since the carry is set, POP will re-enter FAIL once it has restored things to the previous calling level.

All instruction codes other than OR and RTN are skipped over by FAIL.

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## Pre-compiler Subroutine Descriptions

## SRCONT (\$A2E6)

The SRCONT code will be entered when an operator token instruction is found in the Syntax Tables by the main precompiler code. The purpose of the routine is to determine if the current source symbol in the user's line matches the terminal symbol represented by the operator token. If the symbols match, the token is placed into the output buffer and pass is returned. If the symbols do not match, fail is returned.

SRCONT uses the value of the operator token to access the terminal symbol name in the Operator Name Table. The characters in the source symbol are compared to the characters in the terminal symbol. If all the characters match, pass is indicated.

## TNVAR, TSVAR (\$A32A)

These deus ex machina routines are called by the ESRT instruction. The purpose of the routines is to determine if the current source symbol is a valid numeric (TNVAR) or string (TSVAR) variable. If the source symbol is not a valid variable, fail is returned.

When pass is indicated, the routine will put a variable token into the output buffer. The variable token ( $\$ 80-\$ \mathrm{FF}$ ) is an index into the Variable Name Table and the Variable Value Table, plus $\$ 80$.

The Variable Name Table is searched. If the variable is already in the table, the token value for the existing variable is used. If the variable is not in the table, it will be inserted into both tables and a new token value will be used.

A source symbol is considered a valid variable if it starts with an alphabetic character and it is not a symbol in the Operator Name Table, which includes all the reserved words.

The variable is considered to be a string if it ends with $\$$; otherwise it is a numeric variable. If it is a string variable, $\$$ is stored with the variable name characters.

The routine also determines if the variable is an array by looking for (. If the variable is an array, (is stored with the variable name characters in the Variable Name Table. As a result, $A B C, A B C \$$, and $A B C(n)$ are all recognized as different variables.

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## TNCON (\$A400)

TNCON is called by the ESRT instruction. Its purpose is to examine the current source symbol for a numeric constant, using the floating point package. If the symbol is not a numeric constant, the routine returns fail.

If the symbol is a numeric constant, the floating point package has converted it to a floating point number. The resulting six-byte constant is placed in the output buffer preceded by the $\$ 0 \mathrm{E}$ numeric constant token. The routine then exits with pass indicated.

## TSCON (\$A428)

TSCON is called by the ESRT instruction. Its purpose is to examine the current symbol for a string constant. If the symbol is not a string constant, the routine returns fail.

If the first character of the symbol is " $"$, the symbol is a string constant. The routine will place the string constant token (\$0F) into the output buffer, followed by a string length byte, followed by the string characters.

The string constant consists of all the characters that follow the starting double quote up to the ending double quote. If the EOL character (\$9B) is found before the ending double quote, an ending double quote is assumed. The EOL is not part of the string. The starting and ending double quotes are not saved with the string. All 256 character codes except $\$ 9 \mathrm{~B}$ (EOL) and $\$ 22$ (' ${ }^{\prime \prime}$ ) are allowed in the string.

## SEARCH (\$A462)

This is a general purpose table search routine used to find a source symbol character string in a table.

The table to be searched is assumed to have entries which consist of a fixed length part ( 0 to 255 bytes) followed by a variable length ATASCII part. The last character of the ATASCII part is assumed to have the most significant bit ( $\$ 80$ ) on. The last table entry is assumed to have the first ATASCII character as $\$ 00$.

Upon entry, the $X$ register contains the length of the fixed part of the table ( 0 to 255 ). The A, Y register pair points to the start of the table to be searched. The source string for comparison is pointed to by INBUFF plus the value in CIX.

Upon exit, the 6502 carry flag is clear if a match was found, and set if no match was found. The $X$ register points to the end

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of the symbol, plus 1, in the buffer. The SRCADR (\$95) twobyte cell points to the matched table entry. STENUM (\$AF) contains the number, relative to zero, of the matched table entry.

## SETCODE (A2C8)

The SETCODE routine is used to place a token in the next available position in the output (token) buffer. The value in COX determines the current displacement into the token buffer. After the token is placed in the buffer, COX is incremented by one. If COX exceeds 255 , the LINE TOO LONG error message is generated.

## Execution Overview

During the editing and pre-compiling phase, the user's statements were checked for correct syntax, tokenized, and put into the Statement Table. Then direct statements were passed to the Program Executor for immediate processing, while program statements awaited later processing by the Program Executor.

We now enter the execution phase of Atari BASIC. The Program Executor consists of three parts: routines which simulate the function of individual statement types; an expression execution routine which processes expressions (for example, $\mathrm{A}+\mathrm{B}+3, \mathrm{~A}(1,3)$, "HELP", $\mathrm{A}(3)+7.26 \mathrm{E}-13)$; and the Execution Control routine, which manages the whole process.

## Execution Control

Execution Control is invoked in two situations. If the user has entered a direct statement, Execution Control does some initial processing and then calls the appropriate statement execution routine to simulate the requested operation. If the user has entered RUN as a direct statement, the statement execution routine for RUN instructs Execution Control to start processing statements from the beginning of the statement table.

When the editor has finished processing a direct statement, it initiates the Execution Control routine EXECNL (\$A95F). Execution Control's job is to manage the process of statement simulation.

The editor has saved the address of the statement it processed in STMCUR and has put the statement in the Statement Table. Since this is a direct statement, the line number is $\$ 8000$, and the statement is saved as the last line in the Statement Table.

The fact that a direct statement is always the last statement in the Statement Table gives a test for the end of a user's program.

The high-order byte of the direct statement line number ( $\$ 8000$ ) has its most significant bit on. Loading this byte ( $\$ 80$ )

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into the 6502 accumulator will set the minus flag on. The line number of any program statement is less than or equal to \$7FFF. Loading the high order byte (\$7F or less) of a program line number into the accumulator will set the 6502 minus flag off. This gives a simple test for a direct statement.

## Initialization

Execution Control uses several parameters to help it manage the task of statement execution.

STMCUR holds the address in the Statement Table of the line currently being processed.

LLNGTH holds the length of the current line.
NXTSTD holds the displacement in the current line of the next statement to process.

STMCUR already contains the correct value when Execution Control begins processing. SETLN1 (\$B81B) is called to store the correct values into LLNGTH and NXTSTD.

## Statement Execution

Since the user may have changed his or her mind about execution, the routine checks to see if the user hit the break key. If the user did hit BREAK, Execution Control carries out XSTOP (\$B793), the same routine that is executed when the STOP statement is encountered. At the end of its execution, the XSTOP routine gives control to the beginning of the editor.

If the user did not hit BREAK, Execution Control checks to see whether we are at the end of the tokenized line. Since this is the first statement in the line, we can't be at the end of the line. So why do the test? Because this part of the routine is executed once for each statement in the line in order to tell us when we do reach the end of the line. (The end-of-line procedure will be discussed later in this chapter.)

The statement length byte (the displacement to the next statement in the line) is the first byte in a statement. (See Chapter 3.) The displacement to this byte was saved in NXTSTD. Execution Control now loads this new statement's displacement using the value in NXTSTD.

The byte after the statement length in the line is the statement name token. Execution Control loads the statement name token into the A register. It saves the displacement to the next byte, the first of the statement's tokens, in STINDEX for the use of the statement simulation routines.

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The statement name token is used as an index to find this statement's entry in the Statement Execution Table. Each table entry consists of the address, minus 1 , of the routine that will simulate that statement. This simulation routine is called by pushing the address from the table onto the 6502 CPU stack and doing an RTS. Later, when a simulation routine is finished, it can do an RTS and return to Execution Control. (The name of most of the statement simulation routines in the BASIC listing is the statement name preceded by an X : XFOR , XRUN, XLIST.)

Most of the statement simulation routines return to Execution Control after processing.

Execution Control again tests for BREAK and checks for the end of the line. As long as we are not at end-of-line, it continues to execute statements. When we reach end-of-line, it does some end-of-line processing.

## End-of-line Handling in a Direct Statement

When we come to the end of the line in a direct statement, Execution Control has done its job and jumps to SNX3. The READY message is printed and control goes back to the Program Editor.

## End-of-line Handling during Program Execution

Program execution is initiated when the user types RUN. Execution Control handles RUN like any other direct statement. The statement simulation routine for RUN initializes STMCUR, NXTSTD, and LLNGTH to indicate the first statement of the first line in the Statement Table, then returns to Execution Control. Execution Control treats this first program statement as the next statement to be executed, picking up the statement name tokens and calling the simulation routines.

Usually, Execution Control is unaware of whether it is processing a direct statement or a program statement. End-ofline is the only time the routine needs to make a distinction.

At the end of every program line, Execution Control gets the length of the current line and calls GNXTL to update the address in STMCUR to make the next line in the Statement Table the new current line. Then it calls TENDST (\$A9E2) to test the new line number to see if it is another program line or a direct statement. If it is a direct statement, we are at the end of the user's program.

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Since the direct statement includes the RUN command that started program execution, Execution Control does not execute the line. Instead, Execution Control calls the same routine that would have been called if the program had contained an END statement (XEND, at \$B78D). XEND does some end-ofprogram processing, causes READY to be printed, and returns to the beginning of the editor.

If we are not at the end of the user's program, processing continues with the new current line.

## Execution Control Subroutines

## TENDST (\$A9E2)

Exit parameters: The minus flag is set on if we are at the end of program.

This routine checks for the end of the user's program in the Statement Table.

The very last entry in the Statement Table is always a direct statement. Whenever the statement indicated by STMCUR is the direct statement, we have finished processing the user's program.

The line number of a direct statement is $\$ 8000$. The line number of any other statement is \$7FFF or less. TENDST determines if the current statement is the direct statement by loading the high-order byte of the line number into the A register. This byte is at a displacement of one from the address in STMCUR. If this byte is $\$ 80$ (a direct statement), loading it turns the 6502 minus flag on. Otherwise, the minus flag is turned off.

## GETSTMT (\$A9A2)

Entry parameters: TSLNUM contains the line number of the statement whose address is required.

Exit parameters: If the line number is found, the STMCUR contains the address of the statement and the carry flag is set off (clear). If the line number does not exist, STMCUR contains the address where a statement with that line number should be, and the carry flag is set on (set).

The purpose of this routine is to find the address of the statement whose line number is contained in TSLNUM.

The routine saves the address currently in STMCUR into SAVCUR and then sets STMCUR to indicate the top of the

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Statement Table. The line whose address is in STMCUR is called the current line or statement.

GETSTMT then searches the Statement Table for the statement whose line number is in TSLNUM. The line number in TSLNUM is compared to the line number of the current line. If they are equal, then the required statement has been found. Its address is in STMCUR, so GETSTMT clears the 6502 carry flag and is finished.

If TSLNUM is smaller than the current statement line number, GETSTMT gets the length of the current statement by executing GETLL (\$A9DD). GNXTL (\$A9D0) is executed to make the next line in the statement table the current statement by putting its address into STMCUR. GETSTMT then repeats the comparison of TSLNUM and the line number of the current line in the same manner.

If TSLNUM is greater than the current line number, then a line with this line number does not exist. STMCUR already points to where the line should be, the 6502 carry flag is already set, and the routine is done.

## GETLL (\$A9DD)

Entry parameters: STMCUR indicates the line whose length is desired.

Exit parameters: Register A contains the length of the current line.

GETLL gets the length of the current line (that is, the line whose address is in STMCUR).

The line length is at a displacement of two into the line. GETLL-loads the length into the A register and is done.

## GNXTL (\$A9D0)

Entry parameters: STMCUR contains the address of the current line, and register A contains the length of the current line.

Exit parameters: STMCUR contains the address of the next line.

This routine gets the next line in the statement table and makes it the current line.

GNXTL adds the length of the current line (contained in the A register) to the address of the current line in STMCUR. This process yields the address of the next line in the statement table, which replaces the value in STMCUR.

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## SETLN1 (\$B81B)

Entry parameters: STMCUR contains the address of the current line.

Exit parameters: LLNGTH contains the length of the current line. NXTSTD contains the displacement in the line to the next statement to be executed (in this case, the first statement in the line).

This routine initializes several line parameters so that Execution Control can process the line.

The routine gets the length of the line, which is at a displacement of two from the start of the line.

SETLN1 loads a value of three into the Y register to indicate the displacement into the line of the first statement and stores the value into NXTSTD as the displacement to the next statement for execution.

## SETLINE (\$B818)

Entry parameters: TSLNUM contains the line number of a statement.

Exit parameters: STMCUR contains the address of the statement whose line number is in TSLNUM. LLNGTH contains the length of the line. NXTSTD contains the displacement in the line to the next statement to be executed (in this case, the first statement in the line). Carry is set if the line number does not exist.

This routine initializes several line parameters so that execution control can process the line.

SETLINE first calls GETSTMT (\$A9A2) to find the address of the line whose number is in TSLNUM and put that address into STMCUR. It then continues exactly like SETLN1.

## Execute Expression

The Execute Expression routine is entered when the Program Executor needs to evaluate a BASIC expression within a statement. It is also the executor for the LET and implied LET statements.

Expression operators have an order of precedence; some must be simulated before others. To properly evaluate an expression, Execute Expression rearranges it during the evaluation.

## Expression Rearrangement Concepts

Operator precedence rules in algebraic expressions are so simple and so unconscious that most people aren't aware of following them. When you evaluate a simple expression like $\mathrm{Y}=\mathrm{AX}^{2}+\mathrm{BX}+\mathrm{C}$, you don't think: "Exponentiation has a higher precedence than multiplication, which has a higher precedence than addition; therefore, I will first square the X , then perform the multiplication." You just do it.

Computers don't develop habits or common sense - they have to be specifically commanded. It would be nice if we could just type $Y=A X^{2}+B X+C$ into our machine and have the computer understand, but instead we must separate all our variables with operators. We also have to learn a few new operators, such as * for multiply and $\wedge$ for exponentiation.

Given that we are willing to adjust our thinking this much, we enter $Y=A^{*} X^{\wedge} 2+B^{*} X+C$. The new form of expression does not quite have the same feel as $Y=A X^{2}+B X+C$; we have translated normal human patterns halfway into a form the computer can use.

Even the operation $X^{\wedge} 2$ causes another problem for the computer. It would really prefer that we give it the two values first, then tell it what to do with them. Since the computer still needs separators between items, we should write $X^{\wedge} 2$ as $X, 2, \wedge$.

Now we have something the computer can work with. It can obtain the two values $X, 2$, apply the operator $\wedge$, and get a result without having to look ahead.

If we were to transcribe $X^{\wedge} 2^{*} A$ in the same manner, we would have $X, 2, \wedge, A,{ }^{*}$. The value returned by $X, 2, \wedge$ is the first value to multiply, so the value pair for multiplication is ( $\mathrm{X}, 2, \wedge$ ) and A. Again we have two values followed by an operator, and the computer can understand.

If we continue to transcribe the expression by pairing values and operators, we find that we don't want to add the value $X^{\wedge} 2^{*} A$ to $B$; we want to add the value $X^{\wedge} 2^{*} A$ to $B^{*} X$. Therefore, we need to tell the computer $\mathrm{X}, 2, \wedge, A,{ }^{*}, \mathrm{~B}, \mathrm{X},{ }^{*},+$. The value pair for the operator + is $\left(X, 2, \wedge, A,{ }^{*}\right)$ and $\left(B, X,{ }^{*}\right)$.

The value pair for the final operation, $=$, is ( $\mathrm{X}, 2, \wedge, \mathrm{~A},{ }^{*}, \mathrm{~B}, \mathrm{X}$, $*,+, C,+)$ and $Y$. So the complete translation of $Y=A X^{2}+B X+$ C is $\mathrm{X}, 2, \wedge, \mathrm{~A},{ }^{*}, \mathrm{~B}, \mathrm{X},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=$.

Very few people other than Forth programmers put up with this form of expression transcription. Therefore, Atari BASIC was designed to perform this translation for us, provided we use the correct symbols, like * and $\wedge$.

## The Expression Rearrangement Algorithm

The algorithm for expression rearrangement requires two LIFO stacks for temporary storage of the rearranged terms. The Operator Stack is used for temporarily saving operators; the Argument Stack is used for saving arguments. Arguments are values consisting of variables, constants, and the constant-like values resulting from previous expression operations.

## Operator Precedence Table

The Atari BASIC User's Manual lists the operators by precedence. The highest-precedence operators, like $<,>$, and $=<$, are at the top of the list; the lowest-precedence operator, OR, is at the bottom. The operators at the top of the list get executed before the operators at the bottom of the list.

The operators in the precedence table are arranged in the same order as the Operator Name Table. Thus the token values can be used as direct indices to obtain an operator precedence value.

The entry for each operator in the Operator Precedence Table contains two precedence values, the go-onto-stack precedence and the come-off-stack precedence. When a new operator has been plucked from an expression, its go-ontostack precedence is tested in relation to the top-of-stack operator's come-off-stack precedence.

## Expression Rearrangement Procedure

The symbols of the expression (the arguments and the operators) are accessed sequentially from left to right, then rearranged into their correct order of precedence by the following procedure:

1. Initialize the Operator Stack with the Start Of Expression (SOE) operator.
2. Get the next symbol from the expression.
3. If the symbol is an argument (variable or constant), place the argument on the top of the Argument Stack. Go to step 2.
4. If the symbol is an operator, save the operator in the temporary save cell, SAVEOP.
5. Compare the go-onto-stack precedence of the operator in SAVEOP to the come-off stack precedence of the operator on the top of the Operator Stack.
6. If the top-of-stack operator's precedence is less than the precedence of the SAVEOP operator, then the SAVEOP operator is pushed onto the Operator Stack. When the push is done, go back to step 2.
7. If the top-of-stack operator's precedence is equal to or greater than the precedence of the SAVEOP operator, then pop the top-of-stack operator and execute it. When the execution is done, go back to step 5 and continue.
The Expression Rearrangement Procedure has one apparent problem. It seems that there is no way to stop it. There are no exits for the "evaluation done" condition. This problem is handled by enclosing the expression with two special operators: the Start Of Expression (SOE) operator, and the End Of Expression (EOE) operator. Remember that SOE was the first operator placed on the Operator Stack, in step 1. Execution code for the SOE operator will cause the procedure to be exited in step 7, when SOE is popped and executed. The EOE operator is never executed. EOE's function is to force the execution of SOE.

The precedence values of SOE and EOE are set to insure that SOE is executed only when the expression evaluation is finished. The SOE come-off-stack precedence is set so that its value is always less than all the other operators' go-onto-stack precedence values. The EOE go-onto-stack precedence is set so that its value is always equal to or less than all the other

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operators' (including SOE's) come-off-stack precedence values.

Because SOE and EOE precedence are set this way, no operator other than EOE can cause SOE to be popped and executed. Second, EOE will cause all stacked operators, including SOE, to be popped and executed. Since SOE is always at the start of the expression and EOE is always at the end of the expression, SOE will not be executed until the expression is fully evaluated.

In actual practice, the SOE operator is not physically part of the expression in the Statement Table. The Expression Rearrangement Procedure initializes the Operator Stack with the SOE operator before it begins to examine the expression.

There is no single operator defined as the End Of Expression (EOE) operator. Every BASIC expression is followed by a symbol like :, THEN, or the EOL character. All of these symbols function as operators with precedence equivalent to the precedence of our phantom EOE operator. The THEN token, for example, serves a dual purpose. It not only indicates the THEN action, but also acts as the EOE operator when it follows an expression.

## Expression Rearrangement Example

To illustrate how the expression evaluation procedure works, including expression rearrangement, we will evaluate our $Y=A^{*} X^{\wedge} 2+B^{*} X+C$ example and see how the expression is rearranged to $\mathrm{X}, 2, \wedge, \mathrm{~A},{ }^{*}, \mathrm{~B}, \mathrm{X},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=$ with a correct result. To work our example, we need to establish a precedence table for the operators. The values in Figure 7-1 are similar to the actual values of these operators in Atari BASIC. The lowest precedence value is zero; the highest precedence value is $\$ 0 \mathrm{~F}$.

Figure 7-1. Example Precedence Table

| operator <br> symbol | go-on-stack <br> precedence | come-off-stack <br> precedence |
| :---: | :---: | :---: |
| SOE | NA | $\$ 00$ |
| + | $\$ 09$ | $\$ 09$ |
| $*$ | $\$ 0 \mathrm{~A}$ | $\$ 0 \mathrm{~A}$ |
| $\wedge$ | $\$ 0 \mathrm{C}$ | $\$ 0 \mathrm{C}$ |
| $=$ | $\$ 0 \mathrm{~F}$ | $\$ 01$ |
| $!(\mathrm{EOE})$ | $\$ 00$ | NA |

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Symbol values and notations. In the example steps, the term PS $n$ refers to step $n$ in the Expression Rearrangement Procedure (page 57). Step 5, for instance, will be called PS5.

In the actual expression, the current symbol will be underlined. If $B$ is the current symbol, then the actual expression will appear as $Y=A^{*} X \quad 2+\underline{B}^{*} X+C$. In the rearranged expression, the symbols which have been evaluated up to that point will also be underlined.

The values of the variables are:

$$
\begin{array}{ll}
\mathrm{A}=2 & \mathrm{C}=6 \\
\mathrm{~B}=4 & \mathrm{X}=3
\end{array}
$$

The variable values are assumed to be accessed when the variable arguments are popped for operator execution.

The end-of-expression operator is represented by !.

## Example step 1.

Actual Expression: $Y=A^{*} X \wedge 2+B^{*} X+C$ !
Rearranged Expression: $\mathrm{X}, 2, \wedge, \mathrm{~A},{ }^{*}, \mathrm{~B}, \mathrm{X},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=$, !
Argument Stack:
Operator Stack: SOE
SAVEOP:
PS1 has been executed. The Operator Stack has been initialized with the SOE operator. We are ready to start processing the expression symbols.

## Example step 2.

Actual Expression: $\underline{\mathbf{Y}}=A^{*} X \wedge 2+B^{*} X+C$ !
Rearranged Expression: $\mathrm{X}, 2, \wedge, \mathrm{~A},{ }^{*}, \mathrm{~B}, \mathrm{X},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=$, !
Argument Stack: Y
Operator Stack: SOE
SAVEOP:
The first symbol, Y , has been obtained and stacked in the Argument Stack according to PS2 and PS3.

## Example step 3.

Actual Expression: $\mathrm{Y} \equiv \mathrm{A}^{*} \mathrm{X} \wedge 2+\mathrm{B}^{*} \mathrm{X}+\mathrm{C}$ !
Rearranged Expression: $\mathrm{X}, 2, \wedge, \mathrm{~A},{ }^{*}, \mathrm{~B}, \mathrm{X},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=$, !
Argument Stack: Y
Operator Stack: $\mathrm{SOE},=$
SAVEOP: =

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Operator $=$ has been obtained via PS2. The relative precedences of SOE $(\$ 00)$ and $=(\$ 0 \mathrm{~F})$ dictate that the $=$ be placed on the Operator Stack via PS6.

## Example step 4.

Actual Expression: $Y=\underline{A^{*}} X \wedge 2+B^{*} X+C$ !
Rearranged Expression: $\mathrm{X}, 2, \wedge, \mathrm{~A},{ }^{*}, \mathrm{~B}, \mathrm{X},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=$, !
Argument Stack: Y,A
Operator Stack: SOE , $=$
SAVEOP:
The next symbol is A. This symbol is pushed onto the Argument Stack via PS3.

## Example step 5.

Actual Expression: $Y=A * X \wedge 2+B^{*} X+C$ !
Rearranged Expression: $\mathrm{X}, 2, \wedge, \mathrm{~A},{ }^{*}, \mathrm{~B}, \mathrm{X},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=$ !
Argument Stack: Y,A
Operator Stack: SOE,=,*
SAVEOP:
The next symbol is the operator *. The relative precedence of * and = dictates that * be pushed onto the Operator Stack.

## Example step 6.

Actual Expression: $\quad Y=A^{*} \underline{X} \wedge 2+B^{*} X+C$ !
Rearranged Expression: $\mathrm{X}, 2, \wedge, \mathrm{~A},{ }^{*}, \mathrm{~B}, \mathrm{X},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=$ !
Argument Stack: Y,A,X
Operator Stack: SOE, =,*
SAVEOP:
The next symbol is the variable $X$. This symbol is stacked on the Argument Stack according to PS3.

## Example step 7.

Actual Expression: $Y=A^{*} X \_2+B^{*} X+C$ !
Rearranged Expression: $\mathrm{X}, 2, \wedge, \mathrm{~A},{ }^{*}, \mathrm{~B}, \mathrm{X},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=$, !
Argument Stack: Y,A,X
Operator Stack: SOE, =,*, ^
SAVEOP: ^
The next symbol is $\wedge$. The relative precedence of the and the * dictate that $\wedge$ be stacked via PS6.

## Example step 8.

Actual Expression: $Y=A^{*} X^{\wedge} \underline{2}+B^{*} X+C$ !
Rearranged Expression: $\mathrm{X}, 2, \wedge, \mathrm{~A},{ }^{*}, \mathrm{~B}, \mathrm{X},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=,!$
Argument Stack: Y,A,X,2
Operator Stack: SOE $=,{ }^{*}, \wedge$
SAVEOP:
The next symbol is 2 . This symbol is stacked on the Argument Stack via PS3.

## Example step 9.

Actual Expression: $Y=A^{*} X^{\wedge} 2 \pm B^{*} X+C$ !
Rearranged Expression: $\underline{\mathbf{X}, 2, \mathcal{A}^{\prime}, \mathrm{A},{ }^{*}, \mathrm{~B}, \mathrm{X},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=\text { ! }}$
Argument Stack: Y,A,9
Operator Stack: $\mathrm{SOE}_{\boldsymbol{\prime}}=$, ${ }^{*}$
SAVEOP:+
The next symbol is the operator + . The precedence of the 'sperator that was at the top of the stack, $\wedge$, is greater than the precedence of + . PS7 dictates that the top-of-stack operator be popped and executed.

The $\wedge$ operator is popped. Its execution causes arguments $X$ and 2 to be popped from the Argument Stack, replacing the variable with the value that it represents and operating on the two values yielded: $X \wedge 2=3 \wedge 2=9$. The resulting value, 9 , is pushed onto the Argument Stack. The + operator remains in SAVEOP. We continue at PS5.

Note that in the rearranged expression the first symbols, $X, 2, \wedge$, have been evaluated according to plan.

## Example step 10.

Actual Expression: $Y=A^{*} X^{\wedge} 2 \pm B^{*} X+C$ !
Rearranged Expression: $\mathbf{X}, \mathbf{2}, \mathbf{\Lambda}, \mathbf{A},{ }^{*}, \mathrm{~B}, \mathrm{X},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=$, ! Argument Stack: Y,18
Operator Stack: SOE,=
SAVEOP: +
This step originates at PS5. The SAVEOP operator, +, has a precedence that is less than the operator which was at the top of the stack, *. Therefore, according to PS7, the * is popped and executed.

The execution of * results in $\mathrm{A}^{*} 9=2^{*} 9=18$. The resulting value is pushed onto the Argument Stack.

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## Example step 11.

Actual Expression: $Y=A^{*} X^{\wedge} 2 \pm B^{*} X+C$ !
Rearranged Expression: $\mathbf{X}, \mathbf{2}, \mathbf{\Lambda}, \mathbf{A},{ }^{*}, \mathrm{~B}, \mathrm{X},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=$ !
Argument Stack: Y,18
Operator Stack: $\mathrm{SOE},=,+$
SAVEOP:
When step 10 finished, we went to PS5. The operator in SAVEOP was + . Since + has a higher precedence than the top-of-stack operator, $=$, the + operator was pushed onto the Operator Stack via PS6.

## Example step 12.

Actual Expression: $Y=A^{*} X \wedge 2+\underline{B}^{*} X+C$ !
Rearranged Expression: $\underline{\mathbf{X}, \mathbf{2}, \mathbf{A}, \mathbf{A},{ }^{*}, \mathrm{~B}, \mathrm{X},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=, \text { ! }, ~}$
Argument Stack: Y,18,B
Operator Stack: $\mathrm{SOE},=,+$
SAVEOP:
The next symbol is the variable $B$, which is pushed onto the Argument Stack via PS3.

## Example step 13.

Actual Expression: $Y=A^{*} X \wedge 2+B_{-}^{*} X+C$ !
Rearranged Expression: $\mathbf{X}, \mathbf{2}, \mathbf{1}, \mathbf{A},{ }^{*}, \mathrm{~B}, \mathrm{X},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=$, ! Argument Stack: Y,18,B
Operator Stack: $\mathrm{SOE},=,+$, ${ }^{*}$
SAVEOP
The next symbol is the operator *. Since * has a higher precedence than the top-of-stack,+ * is pushed onto the stack via PS6.

## Example step 14.

Actual Expression: $Y=A^{*} X \wedge 2+B^{*} \underline{X}+C$ !
Rearranged Expression: $\underline{X}, 2, \boldsymbol{\Lambda}, \mathbf{A},{ }^{*}, \mathrm{~B}, \mathrm{X},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=$,!
Argument Stack: Y,18,B,X
Operator Stack: $\mathrm{SOE},=,+$, *
SAVEOP:
The variable X is pushed onto the Argument Stack via PS3.

## Example step 15.

Actual Expression: $Y=A^{*} X \wedge 2+B^{*} X \pm C$ !
Rearranged Expression: $\mathbf{X}, \mathbf{2}, \mathbf{\wedge}, \mathbf{A},{ }^{*}, \overline{\mathbf{B}}, \mathbf{X},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=$ !

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Argument Stack: Y,18,12
Operator Stack: SOE, $=,+$
SAVEOP: +
The operator + is retrieved from the expression. Since + has a lower precedence than * which is at the top of the stack, * is popped and executed.

The execution of * causes $B^{*} X=4 * 3=12$. The resulting value of 12 is pushed onto the Argument Stack. We will continue at PS5 via the PS7 exit rule.

## Example step 16.

Actual Expression: $Y=A^{*} X \wedge 2+B^{*} X \pm C$ !
Rearranged Expression: $\mathbf{X}, \mathbf{2}, \mathbf{\Lambda}, \mathbf{A},{ }^{*}, \mathbf{B}, \mathbf{X},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=$, Argument Stack: Y,30
Operator Stack: $\mathrm{SOE}_{,}=$
SAVEOP: +
This step starts at PS5. The SAVEOP operator, +, has precedence that is equal to the precedence of the top-of-stack operator, also + . Therefore, + is popped from the operator stack and executed. The results of the execution cause $18+12$, or 30 , to be pushed onto the Argument Stack. PS5 is called.

## Example step 17.

Actual Expression: $Y=A^{*} X \wedge 2+B^{*} X \pm C!$
Rearranged Expression: $\mathbf{X}, \mathbf{2}, \mathbf{A}, \mathbf{A},{ }^{*}, \overline{\mathbf{B}, \mathbf{X}},{ }^{*},+, \mathrm{C},+, \mathrm{Y},=$ !
Argument Stack: Y,30
Operator Stack: $\mathrm{SOE},=,+$
SAVEOP:
This step starts at PS5. The SAVEOP is + . The top-of-stack operator, $=$, has a lower precedence than + ; therefore, + is pushed onto the stack via PS6.

## Example step 18.

Actual Expression: $Y=A^{*} X \wedge 2+B^{*} X+C$ !
Rearranged Expression: $\mathbf{X}, \mathbf{2}, \mathbf{\Lambda}, \mathbf{A},{ }^{*}, \mathbf{B}, \mathbf{X},{ }^{*},+, C,+, Y,=$ !
Argument Stack: Y,30,C
Operator Stack: $\mathrm{SOE},=,+$
SAVEOP:
The variable $C$ is pushed onto the Argument Stack via PS3.

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## Example step 19.

Actual Expression: $Y=A^{*} X \wedge 2+B^{*} X+C!$
Rearranged Expression: $\mathbf{X}, \mathbf{2}, \boldsymbol{\wedge}, \mathbf{A},{ }^{*}, \mathbf{B}, \mathbf{X}^{-}{ }^{*},+, \mathbf{C},+, \mathrm{Y},=$ ! Argument Stack: Y,36
Operator Stack: SOE,=
SAVEOP:
The EOE operator ! is plucked from the expression. The EOE has a lower precedence than the top-of-stack + operator. Therefore, + is popped and executed. The resulting value of $30+6,36$, is pushed onto the Argument Stack. PS5 will execute next.
Example step 20.
Actual Expression: $Y=A^{*} X \wedge 2+B^{*} X+C$ !
Rearranged Expression: $\underline{\mathbf{X}, 2, \wedge, \mathbf{A},{ }^{*}, \mathbf{B}, \mathbf{X},{ }^{*},+, \mathbf{C},+, \mathbf{Y},=,!}$
Argument Stack:
Operator Stack: SOE
SAVEOP: !
This step starts at PS5. The ! operator has a lower precedence than the top-of-stack = operator, which is popped and executed. The execution of = causes the value 36 to be assigned to Y. This leaves the Argument Stack empty. PS5 will be executed next.

## Example step 21.

Actual Expression: $Y=A * X \wedge 2+B^{*} X+C!$
Rearranged Expression: $\mathbf{X}, 2, \mathbf{\Lambda}, \mathbf{A},{ }^{*}, \mathbf{B}, \mathbf{X},{ }^{*},+, \mathbf{C},+, \mathbf{Y},=,!$
Argument Stack:
Operator Stack:
SAVEOP: !
The ! operator in SAVEOP causes the SOE operator to be popped and executed. The execution of SOE terminates the expression evaluation.

Note that the rearranged expression was executed exactly as predicted.

## Mainline Code

The Execute Expression code implements the Expression Rearrangement Procedure. The mainline code starts at the EXEXPR label at \$AAE0. The input to EXEXPR starts at the current token in the current statement. STMCUR points to the

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current statement. STINDEX contains the displacement to the current token in the STMCUR statement. The output of EXEXPR is whatever values remain on the top of the argument stack when the expression evaluation is finished.

In the following discussion, $\mathrm{PS} n$ refers to the procedure step $n$ in the Expression Rearrangement Procedure.

PS1, initialization, occurs when EXEXPR is entered.
EXPINT is called to initialize the operator and argument stacks.
EXPINT places the SOE operator on the operator stack.
PS2, which obtains the next token, directly follows initialization at EXNXT (\$AAE3). The code calls EGTOKEN to get the next expression symbol and classify it. If the token is an argument, the carry will be set. If the token is an operator, the carry will be clear.

If the token is an argument, PS3 is implemented via a call to ARGPUSH. After the argument is pushed onto the argument stack, EXNXT (PS2) will receive control.

If the token was an operator, then the code at EXOT (\$AAEE) will be executed. This code implements PS4 by saving the token in EXSVOP.

PS5, which compares the precedents of the EXSVOP token and the top-of-stack token, follows EXOT at EXPTST (\$AAFA). This code also executes the SOE operator. If SOE is popped, then Execute Expression finishes via RTS.

If the top-of-stack operator precedence is less than the EXSVOP operator precedence, PS6 is implemented at EOPUSH (\$AB15). EOPUSH pushes EXSVOP onto the operator stack and then goes to EXNXT (PS2).

If the top-of-stack operator precedence is greater than or equal to the EXSVOP operator precedence, then PS7 is implemented at EXOPOP (\$AB0B). EXOPOP will pop the top-of-stack operator and execute it by calling EXOP. When EXOP is done, control passes to EXPTST (PS5).

## Expression Evaluation Stacks

The two expression evaluation stacks, the Argument Stack and the Operator Stack, share a single 256-byte memory area. The Argument Stack grows upward from the lower end of the 256byte area. The Operator Stack grows downward from the upper end of the 256 -byte area.

The 256-byte stack area is the multipurpose buffer at the start of the RAM tables. The buffer is pointed to by the

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ARGSTK (also ARGOPS) zero-page pointer at $\$ 80$. The current index into the Argument Stack is maintained by ARSLVL (\$AA). When the Argument Stack is empty, ARSLVL is zero.

The OPSTKX cell maintains the current index into the Operator Stack. When the Operator Stack is initialized with the SOE operator, OPSTKX is initialized to \$FF. As operators are added to the Operator Stack, OPSTKX is decremented. As arguments are added to the Argument Stack, ARSLVL is incremented.

Since the two stacks share a single 256-byte memory area, there is a possibility that the stacks will run into each other. The code at \$ABC1 is used to detect a stack collision. It does this by comparing the values in ARSLVL and OPSTKX. If ARSLVL is greater than or equal to OPSTKX, then a stack collision occurs, sending the STACK OVERFLOW error to the user.

## Operator Stack

Each entry on the Operator Stack is a single-byte operator-type token. Operators are pushed onto the stack at EXOPUSH (\$AB15) and are popped from the stack at EXOPOP (\$AB0B).

## Argument Stack

Each entry on the Argument Stack is eight bytes long. The format of these entries is described in Figures 7-2, 7-3, and 7-4, and are the same as the formats for entries in the Variable Value Table.

Unlike the Variable Value Table, the Argument Stack must deal with both variables and constants. In Figure 7-2, we see that VNUM is used to distinguish variable entries from constant entries.

The SADR and AADR fields in the entries for strings and arrays are of special interest. (See Figures 7-3 and 7-4.) When a string or array variable is dimensioned, space for the variable is created in the string/array space. The displacement to the start of the variable's area within the string/array space is placed in the SADR/AADR fields at that time. A displacement is used rather than an absolute address because the absolute address can change if any program changes are made after the DIM statement is executed.

Execute Expression needs these values to be absolute address values within the 6502 address space. When a string/array variable is retrieved from the Variable Value Table,

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the displacement is transformed to an absolute address. When (and if) the variable is put back into the Variable Value Table, the absolute address is converted back to a displacement.

The entries for string constants also deserve some special attention. String constants are the quoted strings within the user program. These strings become part of the tokenized statements in the Statement Table. When Execute Expression gets a string token, it will create a string constant Argument Stack entry. This entry's SADR is an absolute address pointer to the string in the Statement Table. SLEN and SDIM are set to the actual length of the quoted string.

## Argument Work Area

An argument which is currently being examined by Execute Expression is kept in a special zero-page Argument Work Area (AWA). The AWA starts at the label VTYPE at \$D2.

Figure 7-2. Argument Stack Entry
0
12 8

| VTYPE | VNUM | DATA |
| :---: | :---: | :---: |
| \begin{tabular}{\|c|c|}
\hline
\end{tabular} |  |  |

Data Field. Format depends on VTYPE.
If $\mathrm{VNUM}=0$, the entry is a constant. If VNUM $>0$, the entry is a variable. In this case, the VNUM value is the entry number in the Variable Value Table. The token representing this variable is VNUM+\$80.

- $\$ 00=$ Data is a six-byte floating point constant.
$\$ 80=$ Data represents an undimensioned string.
$\$ 81=$ Data represents a dimensioned string with a relative address pointer.
$\$ 83=$ Data represents a dimensioned string with an absolute address pointer.
$\$ 40=$ Data represents an undimensioned array.
$\$ 41=$ Data represents a dimensioned array with a relative address pointer.
- \$43 = Data represents a dimensioned array with an absolute address pointer.


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Figure 7-3. Argument Stack String Entry


Figure 7-4. Argument Stack Array Entry


## Chapter Seven

## Operator Executions

An operator is executed when it is popped from the Operator Stack. Execute Expression calls EXOP at \$AB20 to start this execution. The EXOP routine uses the operator token value as an index into the Operator Execution Table (\$AA70). The operator execution address from this table, minus 1 , is placed on the 6502 CPU stack. An RTS is then executed to begin executing the operator's code.

The names of the operator execution routines all begin with the characters $X P$.

All the Atari BASIC functions, such as PEEK, RND, and $A B S$, are executed as operators.

Most routines for the execution of the operators are very simple and straightforward. For example, the * operator routine, XPMUL (\$AC96), pops two arguments, multiplies them via the floating point package, pushes the result onto the argument stack, and returns.

## String, Array, DIM, and Function Operations

Any array reference in an expression may be found in one of two forms: $\mathrm{A}(x)$ or $\mathrm{A}(x, y)$. The indices $x$ and $y$ may be any valid expression. The intent of the indices is to reference a specific array element.

Before the specific element reference can take place, the $x$ and/or $y$ index expressions must be fully evaluated. To do this, the characters ' $(\text { ' }, \text { ' } \text { and ' })^{\prime}$ ' are made operators. The precedence of these operators forces things to happen in the correct sequence. Figure 7-5 shows the relative precedence of these operators for an array.

Figure 7-5. Array Operator Precedence

| operator <br> symbol | go-on-stack <br> precedence | come-off-stack |
| :---: | :---: | :---: |
| ( | $\$ 0 \mathrm{~F}$ | precedence |
| (comma) | $\$ 04$ | $\$ 02$ |
| ) | $\$ 04$ | $\$ 03$ |
|  | $\$ 0 \mathrm{E}$ |  |

As a result of these precedence values, ( has a high enough precedence to go onto the stack, no matter what other operator is on the top of the stack.

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The comma's go-on-stack precedence will force all operators except ( to be popped and executed. As a result, the $x$ index sub-expression, in the expression $\mathrm{A}(x, y)$, will be fully evaluated and the final $x$ index value will be pushed onto the Argument Stack.

The comma will then be placed onto the Operator Stack. Its come-off-stack precedence is such that no other operator, except ), will pop it off.

The ) operator precedence will force any $y$ index expression to be fully evaluated and the $y$ index result value to be placed onto the Argument Stack.

It will then force the comma operator to be popped and executed. This action results in a comma counter being incremented.

The ) will then force the ( to be popped and executed. The execution of ( results in the proper array element being referenced. The ( operator will pop the indices from the Argument Stack. The number of indices (either zero or one) to be popped is governed by the comma counter, which was incremented by one for each comma that was popped and executed.

Atari BASIC has numerous ( tokens, and each causes a different ( routine to be executed. These ( operators are array (CALPRN), string (CSLPRN), array DIM (CDLPRN), string DIM (CDSLPR), function (CFLPRN), and the expression grouping CLPRN operator. The Syntax Table pseudo-instruction CHNG is used to change the CLPRN token to the other ( tokens in accordance with the context of the grammar.

The expression operations for each of these various ( operators in relation to commas and (is exactly the same. When ( is executed, the comma count will show how many arguments the operator's code must pop from the argument stack. Each of these arguments will have been evaluated down to a single value in the form of a constant.

# Execution Boundary Conditions 

BASIC Language statements can be divided into groups with related functions. The execution boundary statements, RUN, STOP, CONT and END, cause a BASIC program to start or stop executing. The routines which simulate these statements are XRUN, XSTOP, XCONT, and XEND.

## Program Termination Routines

Any BASIC statement can be used as either a direct statement or a program statement, but some only make sense in one mode. The STOP statement has no real meaning when entered as a direct statement. When the statement simulation routine for STOP is asked to execute in direct mode, it does as little processing as possible and exits. Useful processing occurs only when STOP is a program statement.
STOP (\$B7A7). The XSTOP and XEND routines are similar and perform some of the same tasks. The tasks common to both are handled by the STOP routine.

If this statement is not a direct statement, the STOP routine saves the line number of the current line in STOPLN. This line number is used later for printing the STOPed message. It is also used by the CONT simulation routine (XCONT) to determine where to restart program execution. (Since XEND also uses this routine, it is possible to CONTinue after an END statement in the middle of a program.)

The STOP routine also resets the LIST and ENTER devices to the screen and the keyboard.
XSTOP (\$B793). XSTOP does the common STOP processing and then calls :LPRTOKEN(\$B535) to print the STOPed message. It then calls one of the error printing routines, :ERRM2 (\$B974), to output the AT LINE nnn portion. The :ERRM2 routine will not print anything if this was a direct statement. When :ERRM2 is finished, it jumps back to the start of the editor.

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XEND (\$B78D). XEND calls the STOP routine to save the current line number. It then transfers to the start of the editor via the SNX1 entry point. This turns off the sound, closes any open IOCBs, and prints the READY message. XEND also leaves values on the 6502 CPU stack. These values are thrown away when the editor resets the stack.
END OF PROGRAM. A user may have neglected to include an END statement in his program. In this case, when Execution Control comes to the end of the Statement Table it calls XEND, and the program is terminated exactly as if the last statement in the program were an END.

## Program Initiation Routines

The statements that cause a user's program to begin execution are RUN and CONT. These statements are simulated by XRUN and XCONT.
XCONT (\$B7BE). The CONT statement has no meaning when encountered as a program statement, so its execution has no effect.

When the user enters CONT as a direct statement, XCONT uses the line number that was saved in STOPLN to set Execution Control's line parameters (STMCUR, NXTSTD, and LLNGTH). This results in the current line being the line following the one whose line number is in STOPLN. This means that any statement following STOP or END on a line will not be executed; therefore, STOP and END should always be the last statement in the line.

If we are at the end of the Statement Table, XCONT terminates as if an END statement had been encountered in the program. If there are more lines to process, XCONT returns to Execution Control, which resumes processing at the line whose address was just put into STMCUR.
XRUN (\$B74D). The RUN statement comes in two formats, RUN and RUN < filespec > . In the case of RUN < filespec>, XRUN executes XLOAD to load a saved program, which replaces the current one in memory. The process then proceeds like RUN.

XRUN sets up Execution Control's line pointers to indicate the first line in the Statement Table. It clears some flags used to control various other BASIC statements; for example, it resets STOPLN to 0. It closes all IOCBs and executes XCLR to reset all

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the variables to zero and get rid of any entries in the String/Array Table or the Runtime Stack.

If there is no program, so the only thing in the Statement Table is the direct statement, then XRUN does some clean-up, prints READY, and returns to the start of the editor, which resets the 6502 CPU stack.

If there is a program, XRUN returns to Execution Control, which starts processing the first statement in the table as the current statement.

When RUN < filespec > is used as a program statement, it performs the useful function of chaining to a new program, but if RUN alone is used as a program statement, an infinite loop will probably result.

## Error Handling Routine

There are other conditions besides the execution boundary statements that terminate a program's execution. The most familiar are errors.

There are two kinds of errors that can occur during execution: Input/Output errors and BASIC language errors.

Any BASIC routine that does I/O calls the IOTEST routine (\$BCB3) to check the outcome of the operation. If an error that needs to be reported to the user is indicated, IOTEST gets the error number that was returned by the Operating System and joins the Error Handling Routine, ERROR (\$B940), which finishes processing the error.

When a BASIC language error occurs, the error number is generated by the Error Handling Routine. This routine calculates the error by having an entry point for every BASIC language error. At each entry point, there is a 6502 instruction that increments the error number. By the time the main routine, ERROR, is reached, the error number has been generated.

The Error Handling Routine calls STOP (\$B7A7) to save the line number of the line causing the error in STOPLN. It tests TRAPLN to see if errors are being TRAPed. The TRAP option is on if TRAPLN contains a valid line number. In this case, the Error Handler does some clean-up and joins XGOTO, which transfers processing to the desired line.

If the high-order byte of the line number is $\$ 80$ (not a valid line number), then we are not TRAPing errors. In this case, the Error Handler prints the four-part error message, which

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consists of ERROR, the error number, AT LINE, and finally the line number. If the line in error was a direct statement, the AT LINE part is not printed. The error handler resets ERRNUM to zero and is finished.

The Error Handling Routine does not do an orderly return, but jumps back to the start of the editor at the SYNTAX entry point where the 6502 stack is reset, clearing it of the nowunwanted return addresses.

# Program Flow Control Statements 

Execution Control always processes the statement in the Statement Table that follows the one it thinks it has just finished. This means that statements in a BASIC program are usually processed in sequential order.

Several statements, however, can change that order: GOTO, IF, TRAP, FOR, NEXT, GOSUB, RETURN, POP, and ON. They trick Execution Control by changing the parameters that it maintains.

## Simple Flow Control Statements

## XGOTO (\$B6A3)

The simplest form of flow control transfer is the GOTO statement, simulated by the XGOTO routine.

Following the GOTO token in the tokenized line is an expression representing the line number of the statement that the user wishes to execute next. The first thing the XGOTO routine does is ask Execute Expression to evaluate the expression and convert it to a positive integer. XGOTO then calls the GETSTMT routine to find this line number in the Statement Table and change Execution Control's line parameters to indicate this line.

If the line number does not exist, XGOTO restores the line parameters to indicate the line containing the original GOTO, and transfers to the Error Handling Routine via the ERNOLN entry point. The Error Handling Routine processes the error and jumps to the start of the editor.

If the line number was found, XGOTO jumps to the beginning of Execution Control (EXECNL) rather than returning to the point in the routine from which it was called. This leaves garbage on the 6502 CPU stack, so XGOTO first pulls the return address off the stack.

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## XIF (\$B778)

The IF statement changes the statement flow based on a condition. The simulation routine, XIF, begins by calling a subroutine of Execute Expression to evaluate the condition. Since this is a logical (rather than an arithmetic) operation, we are only interested in whether the value is zero or non-zero. If the expression was false (non-zero), XIF modifies Execution Control's line parameters to indicate the end of this line and then returns. Execution Control moves to the next line, skipping any remaining statements on the original IF statement line.

If the expression is true (zero), things get a little more complicated. Back during syntaxing, when a statement of the form IF < expression > THEN < statement > was encountered, the pre-compiler generated an end-of-statement token after THEN. XIF now tests for this token. If we are at the end of the statement, XIF returns to Execution Control, which processes what it thinks is the next statement in the current line, but which is actually the THEN < statement > part of the IF statement.

If XIF does not find the end-of-statement token, then the statement must have had the form IF < expression > THEN <line number>. XIF jumps to XGOTO, which finishes processing by changing Execution Control's line parameters to indicate the new line.

## XTRAP (\$B7E1)

The TRAP statement does not actually change the program flow when it is executed. Instead, the XTRAP simulation routine calls a subroutine of Execute Expression to evaluate the line number and then saves the result in TRAPLN (\$BC).

The program flow is changed only if there is an error. The Error Handling Routine checks TRAPLN. If it contains a valid line number, the error routine does some initial set-up and joins the XGOTO routine to transfer to the new line.

## Runtime Stack Routines

The rest of the Program Flow Control Statements use the Runtime Stack. They put items on the stack, inspect them, and/or remove them from the stack.

Every item on the Runtime Stack contains a four-byte header. This header consists of a one-byte type indication, a

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two-byte line number, and a one-byte displacement to the Statement Name Token. (See pages 18-19.) The type byte is the last byte placed on the stack for each entry. This means that the pointer to the top of the Runtime Stack (RUNSTK) points to the type byte of the most recent entry on the stack. A zero type byte indicates a GOSUB-type entry. Any non-zero type byte represents a FOR-type entry.

A GOSUB entry consists solely of the four-byte header. A FOR entry contains twelve additional bytes: a six-byte limit value and a six-byte step value.

Several routines are used by more than one of the statement simulation routines.
PSHRSTK (\$B683) This routine expands the Runtime Stack by calling EXPLOW and then storing the type byte, line number, and displacement of the Statement Name Token on the stack.
POPRSTK (\$B841) This routine makes sure there really is an entry on the Runtime Stack. POPRSTK saves the displacement to the statement name token in SVDISP, saves the line number in TSLNUM, and puts the type/variable number in the 6502 accumulator. It then removes the entry by calling the CONTLOW routine.
:GETTOK (\$B737) This routine first sets up Execution Control's line parameters to point to the line whose number is in the entry just pulled from the Runtime Stack. If the line was found, :GETTOK updates the line parameters to indicate that the statement causing this entry is now the current statement. Finally, it loads the 6502 accumulator with the statement name token from the statement that created this entry and returns to its caller.

If the line number does not exist, :GETTOK restores the current statement address and exits via the ERGFDEL entry point in the Error Handling Routine.

Now let's look at the simulation routines for the statements that utilize the Runtime Stack.

## XFOR (\$B64B)

XFOR is the name of the simulation routine which executes a FOR statement.

In the statement FOR I = 1 TO 10 STEP 2:
I is the loop control variable

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1 is its initial value
10 is the limit value
2 is the step value
XFOR calls Execute Expression, which evaluates the initial value and puts it in the loop control variable's entry in the Variable Value Table.

Then it calls a routine to remove any currently unwanted stack entries - for example, a previous FOR statement that used the same loop control variable as this one.

XFOR calls a subroutine of Execute Expression to evaluate the limit and step values. If no step value was given, a value of 1 is assigned. It expands the Runtime Stack using EXPLOW and puts the values on the stack.

XFOR uses PSHRSTK to put the header entry on the stack. It uses the variable number of the loop control variable (machine-language ORed with \$80) as the type byte. XFOR now returns to Execution Control, which processes the statement following the FOR statement.

The FOR statement does not change program flow. It just sets up an entry on the Runtime Stack so that the NEXT statement can change the flow.

XNEXT (\$B6CF)
The XNEXT routine decides whether to alter the program flow, depending on the top Runtime Stack entry. XNEXT calls the POPRSTK routine repeatedly to remove four-byte header entries from the top of the stack until an entry is found whose variable number (type) matches the NEXT statement's variable token. If the top-of-stack or GOSUB-type entry is encountered, XNEXT transfers control to an Error Handling Routine via the ERNOFOR entry point.

To compute the new value of the loop variable, XNEXT calls a subroutine of Execute Expression to retrieve the loop control variable's current value from the Variable Value Table, then gets the step value from the Runtime Stack, and finally adds the step value to the variable value. XNEXT again calls an Execute Expression subroutine to update the variable's value in the Variable Value Table.

XNEXT gets the limit value from the stack to determine if the variable's value is at or past the limit. If so, XNEXT returns to Execution Control without changing the program flow, and the next sequential statement is processed.

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If the variable's value has not reached the limit, XNEXT returns the entry to the Runtime Stack and changes the program flow. POPRSTK already saved the line number of the FOR statement in TSLNUM and the displacement to the statement name token in SVDISP. XNEXT calls the :GETTOK routine to indicate the FOR statement as the current statement.

If the token at the saved displacement is not a FOR statement name token, then the Error Handling Routine is given control at the ERGFDEL entry point. Otherwise, XNEXT returns to Execution Control, which starts processing with the statement following the FOR statement.

## XGOSUB (\$B6A0)

The GOSUB statement causes an entry to be made on the Runtime Stack and also changes program flow.

The XGOSUB routine puts the GOSUB-type indicator (zero) into the 6502 accumulator and calls PSHRSTK to put a four-byte header entry on the Runtime Stack for later use by the simulation routine for RETURN. XGOSUB then processes exactly like XGOTO.

## XRTN (\$B719)

The RETURN statement causes an entry to be removed from the Runtime Stack. The XRTN routine uses the information in this entry to determine what statement should be processed next.

The XRTN first calls POPRSTK to remove a GOSUB-type entry from the Runtime Stack. If there are no GOSUB entries on the stack, then the Error Handling Routine is called at ERBRTN. Otherwise, XRTN calls: GETTOK to indicate that the statement which created the Runtime Stack entry is now the current statement.

If the statement name token at the saved displacement is not the correct type, then XRTN exits via the Error Handling Routine's ERGFDEL entry point. Otherwise, control is returned to the caller. When Execution Control was the caller, then GOSUB must have created the stack entry, and processing will start at the statement following the GOSUB.

Several other statements put a GOSUB-type entry on the stack when they need to mark their place in the program. They do not affect program flow and will be discussed in later chapters.

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## XPOP (\$B841)

The XPOP routine uses POPRSTK to remove an entry from the Runtime Stack. A user might want to do this if he decided not to RETURN from a GOSUB.

## XON (\$B7ED)

The ON statement comes in two versions: ON-GOTO and ONGOSUB. Only ON-GOSUB uses the Runtime Stack.

The XON routine evaluates the variable and converts it to an integer (MOD 256). If the value is zero, XON returns to Execution Control without changing the program flow.

If the value is non-zero and this is an ON-GOSUB statement, XON puts a GOSUB-type entry on the Runtime Stack for RETURN to use later.

From this point, ON-GOSUB and ON-GOTO perform in exactly the same manner. XON uses the integer value calculated earlier to index into the tokenized statement line to the correct GOTO or GOSUB line number. If there is no line number corresponding to the index, XON returns to Execution Control without changing program flow. Otherwise, XON joins XGOTO to finish processing.

# Tokenized Program Save and Load 

The tokenized program can be saved to and reloaded from a peripheral device, such as a disk or a cassette. The primary statement for saving the tokenized program is SAVE. The saved program is reloaded into RAM with the LOAD statement. The CSAVE and the CLOAD statements are special versions of SAVE and LOAD for use with a cassette.

## Saved File Format

The tokenized program is completely contained within the Variable Name Table, the Variable Value Table, and the Statement Table. However, since these tables vary in size, we must also save some information about the size of the tables.

The SAVE file format is shown in Figure 10-1. The first part consists of seven fields, each of them two bytes long, which tell where each table starts or ends. Part two contains the saved program's Variable Name Table (VNT), Variable Value Table (VVT), and Statement Table (ST).

The displacement value in all the part-one fields is actually the displacement plus 256. We must subtract 256 from each displacement value to obtain the true displacement.

The VNT starts at relative byte zero in the file's second part. The second field in part one holds that value plus 256.

The DVVT field in part one contains the displacement, minus 256, of the VVT from the start of part two.

The DST value, minus 256, gives the displacement of the Statement Table from the start of part two.

The DEND value, minus 256, gives the end-of-file displacement from the start of part two.

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Figure 10-1. SAVE File Format
PART 1

## XSAVE (\$BB5D)

The code that implements the SAVE statement starts at the XSAVE (\$BB5D) label. Its first task is to open the specified output file, which it does by calling ELADVC.

The next operation is to move the first seven RAM table pointers from $\$ 80$ to a temporary area at $\$ 500$. While these pointers are being moved, the value contained in the first pointer is subtracted from the value in each of the seven pointers, including the first.

Since the first pointer held the absolute address of the first RAM table, this results in a list of displacements from the first RAM table to each of the other tables. These seven two-byte displacements are then written from the temporary area to the file via IO3. These are the first fourteen bytes of the SAVE file. (See Figure 10-1.)

The first RAM table is the 256-byte buffer, which will not be SAVEd. This is why the seven two-byte fields at the beginning of the SAVEd file hold values exactly 256 more than the true

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displacement of the tables they point to. (The LOAD procedure will resolve the 256-byte discrepancy.)

The next operation is to write the three needed RAM tables. The total length of these tables is determined from the value in the seventh entry in the displacement list, minus 256 . To write the three entries, we point to the start of the Variable Name Table and call IO4, with the length of the three tables. This saves the second part of the file format.

The file is then closed and XSAVE returns to Execution Control.

## XLOAD (\$BAFB)

The LOAD statement is implemented at the XLOAD label located at \$BAFB.

XLOAD first opens the specified load file for input by calling ELADVC. BASIC reads the first fourteen bytes from the file into a temporary area starting at $\$ 500$. These fourteen bytes are the seven RAM table displacements created by SAVE.

The first two bytes will always be zero, according to the SAVE file format. (See Figure 10-1.) BASIC tests these two bytes for zero values. If these bytes are not zero, BASIC assumes the file is not a valid SAVE file and exits via the ERRNSF, which generates error code 21 (Load File Error).

If this is a valid SAVE file, the value in the pointer at $\$ 80$ (Low Memory Address) is added to each of the seven displacements in the temporary area. These values will be the memory addresses of the three RAM tables, if and when they are read into memory.

The seventh pointer in the temporary area contains the address where the end of the Statement Table will be. If this address exceeds the current system high memory value, the routine exits via ERRPTL, which generates error code 19 (Load Program Too Big).

If the program will fit, the seven addresses are moved from the temporary area to the RAM table pointers at $\$ 80$. The second part of the file is then loaded into the area now pointed to by the Variable Name Table pointer $\$ 82$. The file is closed, CLR is executed, and a test for RUN is made.

If RUN called XLOAD, then a value of \$FF was pushed onto the CPU stack. If RUN did not call XLOAD, then $\$ 00$ was pushed onto the CPU stack. If RUN was the caller, then an RTS is done.

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If XLOAD was entered as a result of a LOAD or CLOAD statement, then XLOAD exits directly to the Program Editor, not to Execution Control.

## CSAVE and CLOAD

The CSAVE and CLOAD statements are special forms of SAVE and LOAD. These two statements assume that the SAVE/LOAD device is the cassette device.

CSAVE is not quite the same as SAVE " $\mathrm{C}:$ ". Using SAVE with the "C:" device name will cause the program to be saved using long cassette inter-record gaps. This is a time waster, and CSAVE uses short inter-record gaps.

CSAVE starts at XCSAVE (\$BBAC). CLOAD starts at XCLOAD (\$BBA4).

# The LIST and ENTER Statements 

LIST can be used to store a program on an external device and ENTER can retrieve it. The difference between LOAD-SAVE and LIST-ENTER is that LOAD-SAVE deals with the tokenized program, while LIST-ENTER deals with the program in its source (ATASCII) form.

## The ENTER Statement

BASIC is in ENTER mode whenever a program is not RUNning. By default the Program Editor looks for lines to be ENTERed from the keyboard, but the editor handles all ENTERed lines alike, whether they come from the keyboard or not.

## The Enter Device

To accomplish transparency of all input data (not just ENTERed lines), BASIC maintains an enter device indicator, ENTDTD (\$B4). When a BASIC routine (for example, the INPUT simulation routine) needs data, an I/O operation is done to the IOCB specified in ENTDTD. When the value in ENTDTD is zero, indicating IOCB 0 , input will come from the keyboard. When data is to come from some other device, ENTDTD contains a number indicating the corresponding IOCB. During coldstart initialization, the enter device is set to IOCB 0 . It is also reset to 0 at various other times.

## XENTER (\$BACB)

The XENTER routine is called by Execution Control to simulate the ENTER statement. XENTER opens IOCB 7 for input using the specified <filespec $>$, stores a 7 in the enter device ENTDTD, and then jumps to the start of the editor.

## Entering from a Device

When the Program Editor asks GLGO, the get line routine (\$BA92), for the next line, GLGO tells CIO to get a line from the

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device specified in ENTDTD - in this case, from IOCB 7. The editor continues to process lines from IOCB 7 until an end-offile error occurs. The IOTEST routine detects the EOF condition, sees that we are using IOCB 7 for ENTER, closes device 7, and jumps to SNX2 to reset the enter device (ENTDTD) to 0 and print the READY message before restarting at the beginning of the editor.

## The LIST Statement

The routine which simulates the LIST statement, XLIST, is actually another example of a language translator, complete with symbols and symbol-combining rules. XLIST translates the tokens generated by Atari BASIC back into the semiEnglish BASIC statements in ATASCII. This translation is a much simpler task than the one done by the pre-compiler, since XLIST can assume that the statement to be translated is syntactically correct. All that is required is to translate the tokens and insert blanks in the appropriate places.

## The List Device

BASIC maintains a list device indicator, LISTDTD (\$B5), similar to the enter device indicator discussed earlier. When a BASIC routine wants to output some data (an error message, for example), the I/O operation is done to the device (IOCB) specified in LISTDTD.

During coldstart initialization and at various other times, LISTDTD is set to zero, representing IOCB 0 , the editor, which will place the output on the screen. Routines such as XPRINT or XLIST can change the LIST device to indicate some other IOCB. Thus the majority of the BASIC routines need not be concerned about the output's destination.

Remember that IOCB 0 is always open to the editor, which gets input from the keyboard and outputs to the screen. IOCB 6 is the $S$ : device, the direct access to graphics screen, which is used in GRAPHICS statements. Atari BASIC uses IOCB 7 for I/O commands that allow different devices, like SAVE, LOAD, ENTER, and LIST.

## XLIST (\$B483)

The XLIST routine considers the output's destination in its initialization process and then forgets about it. It looks at the first expression in the tokenized line. If it is the < filespec >

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string, XLIST calls a routine to open the specified device using IOCB 7 and to store a 7 in LISTDTD. All of XLIST's other processing is exactly the same, regardless of the LISTed data's final destination.

XLIST marks its place in the Statement Table by calling a subroutine of XGOSUB to put a GOSUB type entry on the Runtime Stack. Then XLIST steps through the Statement Table in the same way that Execution Control does, using Execution Control's line parameters and subroutines. When XLIST is finished, Execution Control takes the entry off the Runtime Stack and continues.

The XLIST routine, assuming it is to LIST all program statements, sets default starting and ending line numbers of 0 (in TSLNUM) and \$7FFF (in LELNUM).

XLIST then determines whether line numbers were specified in the tokenized line that contained the LIST statement. XLIST compares the current index into the line (STINDEX) to the displacement to the next statement (NXTSTD). If STINDEX is not pointing to the next statement, at least one line number is specified. In this case, XLIST calls a subroutine of Execute Expression to evaluate the line number and convert it to a positive integer, which XLIST stores in TSLNUM as the starting line number.

If a second line number is specified, XLIST calls Execute Expression again and stores the value in LELNUM as the final line to LIST. If there is no second line number, then XLIST makes the ending line number equal to the starting line number, and only one line will be LISTed. If no line numbers were present, then TSLNUM and LELNUM still contain their default values, and all the program lines will be LISTed.

XLIST gets the first line to be LISTed by calling the Execution Control subroutine GETSTMT to initialize the line parameters to correspond to the line number in TSLNUM. If we are not at the end of the Statement Table, and if the current line's number is less than or equal to the final line number to be LISTed, XLIST calls a subroutine :LLINE to list the line.

After LISTing the line, XLIST calls Execution Control's subroutines to point to the next line. LISTing continues in this manner until the end of the Statement Table is reached or until the final line specified has been printed.

When XLIST is finished, it exits via XRTN at $\$ B 719$, which makes the LIST statement the current statement again and then returns to Execution Control.

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## LIST Subroutines

:LLINE (\$B55C)
The :LLINE routine LISTs the current line (the line whose address is in STMCUR).
:LLINE gets the line number from the beginning of the tokenized line. The floating point package is called to convert the integer to floating point and then to printable ATASCII. The result is stored in the buffer indicated by INBUFF. :LLINE calls a subroutine to print the line number and then a blank.

For every statement in the line, :LLINE sets STINDEX to point to the statement name token and calls the :LSTMT routine (\$B590) to LIST the statement. When all statements have been LISTed, :LLINE returns to its caller, XLIST.

## :LSTMT (\$B590)

The :LSTMT routine LISTs the statement which starts at the current displacement (in STINDEX) into the current line. This routine does the actual language translation from tokens to BASIC statements.
:LSTMT uses two subroutines, :LGCT and :LGNT, to get the current and next token, respectively. If the end of the statement has been reached, these routines both pull the return address of their caller off the 6502 CPU stack and return to :LSTMT's caller, :LLINE. Otherwise, they return the requested token from the tokenized statement line.

The first token in a statement is the statement name token. :LSTMT calls a routine which prints the corresponding statement name by calling :LSCAN to find the entry and :LPRTOKEN to print it.

In the discussion of the Program Editor we saw that an erroneous statement was given a statement name of ERROR and saved in the Statement Table. If the current statement is this ERROR statement or is REM or DATA, :LSTMT picks up each remaining character in the statement and calls PRCHAR (\$BA9F) to print the character.

Each type of token is handled differently. :LSTMT determines the type (variable, numeric constant, string constant, or operator) and goes to the proper code to translate it.

Variable Token. A variable token has a value greater than or equal to $\$ 80$. When :LSTMT encounters a variable token, it

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turns off the most significant bit to get an index into the Variable Name Table. :LSTMT asks the :LSCAN routine to get the address of this entry. :LSTMT then calls :LPRTOKEN (\$B535) to print the variable name. If the last character of the name is (, the next token is an array left parenthesis operator, and :LSTMT skips it.
Numeric Constant Token. A numeric constant is indicated by a token of $\$ 0 \mathrm{E}$. The next six bytes are a floating point number. :LSTMT moves the numeric constant from the tokenized line to FRO (\$D4) and asks the floating point package to convert it to ATASCII. The result is in a buffer pointed to by INBUFF. :LSTMT moves the address of the ATASCII number to SRCADR and tells :LPRTOKEN to print it.
String Constant Token. A string constant is indicated by a token of $\$ 0 \mathrm{~F}$. The next byte is the length of the string followed by the actual string data. Since the double quotes are not stored with a string constant, :LSTMT calls PRCHAR (\$BA9F) to print the leading double quote. The string length tells :LSTMT how many following characters to print without translation.
:LSTMT repeatedly gets a character and calls PRCHAR to print it until the whole string constant has been processed. It then asks PRCHAR to print the ending double quote.
Operator Token. An operator token is any token greater than or equal to $\$ 10$ and less than $\$ 80$. By subtracting $\$ 10$ from the token value, :LSTMT creates an index into the Operator Name Table. :LSTMT calls :LSCAN to find the address of this entry. If the operator is a function (token value greater than or equal to \$3D), :LPROTOKEN is called to print it. If this operator is not a function but its name is alphabetic (such as AND), the name is printed with a preceding and following blank. Otherwise, :LPRTOKEN is called to print just the operator name.

## :LSCAN (\$B50C)

This routine scans a table until it finds the translation of a token into an ATASCII name. A token's value is based on its table entry number; therefore, the entry number can be derived by modifying the token. For example, a variable token is created by machine-language ORing the table entry number of the variable name with $\$ 80$. The entry number can be produced by ANDing out the high-order bit of the token. It is this entry number, stored in SCANT, that the :LSCAN routine uses.

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## :LPRTOKEN (\$B535)

This routine's task is to print the string of ATASCII characters whose address is in SCRADR. :LPRTOKEN makes sure the most significant bit is off (except for a carriage return) and prints the characters one at a time until it has printed the last character in the string (the one with its most significant bit on).

# Atari Hardware Control Statements 


#### Abstract

The Atari Hardware Control Statements allow easy access to some of the computer's graphics and audio capabilities. The statements in this group are COLOR, GRAPHICS, PLOT, POSITION, DRAWTO, SETCOLOR, LOCATE, and SOUND.


## XGR (\$BA50)

The GRAPHICS statement determines the current graphics mode. The XGR simulation routine executes the GRAPHICS statement. The XGR routine first closes IOCB 6. It then calls an Execute Expression subroutine to evaluate the graphics mode value and convert it to an integer.

XGR sets up to open the screen by putting the address of a string " $\mathrm{S}:$ :" into INBUFF. It creates an AUX1 and AUX2 byte from the graphics mode integer. XGR calls a BASIC I/O routine which sets up IOCB 6 and calls CIO to open the screen for the specified graphics mode. Like all BASIC routines that do I/O, XGR jumps to the IOTEST routine, which determines what to do next based on the outcome of the I/O.

## XCOLOR (\$BA29)

The COLOR statement is simulated by the XCOLOR routine. XCOLOR calls a subroutine of Execute Expression to evaluate the color value and convert it to an integer. XCOLOR saves this value (MOD 256) in BASIC memory location COLOR (\$C8). This value is later retrieved by XPLOT and XDRAWTO.

## XSETCOLOR (\$B9B7)

The routine that simulates the SETCOLOR statement, XSETCOLOR, calls a subroutine of Execute Expression to evaluate the color register specified in the tokenized line. The Execute Expression routine produces a one-byte integer. If the value is not less than 5 (the number of color registers), XSETCOLOR exits via the Error Handling Routine at entry point ERVAL. Otherwise, it calls Execute Expression to get two more integers from the tokenized line.

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To calculate the color value, XSETCOLOR multiplies the first integer (MOD 256) by 16 and adds the second (MOD 256). Since the operating system's five color registers are in consecutive locations starting at \$2C4, XSETCOLOR uses the register value specified as an index to the proper register location and stores the color value there.

## XPOS (\$BA16)

The POSITION statement, which specifies the $X$ and $Y$ coordinates of the graphics cursor, is simulated by the XPOS routine.

XPOS uses a subroutine of Execute Expression to evaluate the $X$ coordinate of the graphics window cursor and convert it to an integer value. The two-byte result is stored in the operating system's X screen coordinate location (SCRX at \$55). This is the column number or horizontal position of the cursor.

XPOS then calls another Execute Expression subroutine to evaluate the $Y$ coordinate and convert it to a one-byte integer. The result is stored in the $Y$ screen coordinate location (SCRY at $\$ 54)$. This is the row number, or vertical position.

## XLOCATE (\$BC95)

XLOCATE, which simulates the LOCATE statement, first calls XPOS to set up the $X$ and $Y$ screen coordinates. Next it initializes IOCB 6 and joins a subroutine of XGET to do the actual I/O required to get the screen data into the variable specified.

## XPLOT (\$BA76)

XPLOT, which simulates the PLOT statement, first calls XPOS to set the $X$ and $Y$ coordinates of the graphics cursor. XPLOT gets the value that was saved in COLOR (\$C8) and joins a PUT subroutine (PRCX at \$BAA1) to do the I/O to IOCB 6 (the screen).

## XDRAWTO (\$BA31)

The XDRAWTO routine draws a line from the current $\mathrm{X}, \mathrm{Y}$ screen coordinates to the $X, Y$ coordinates specified in the statement. The routine calls XPOS to set the new X, Y coordinates. It places the value from BASIC's memory location COLOR into OS location SVCOLOR (\$2FB). XDRAWTO does some initialization of IOCB 6 specifying the draw command (\$11). It then calls a BASIC I/O routine which finishes the
initialization of IOCB 6 and calls CIO to draw the line. Finally, XDRAWTO jumps to the IOTEST routine, which will determine what to do next based on the outcome of the I/O.

## XSOUND (\$B9DD)

The Atari computer hardware uses a set of memory locations to control sound capabilities. The SOUND statement gives the user access to some of these capabilities. The XSOUND routine, which simulates the SOUND statement, places fixed values in some of the sound locations and user specified values in others.

The XSOUND routine uses Execute Expression to get four integer values from the tokenized statement line. If the first integer (voice) is greater than or equal to 4, the Error Handling Routine is invoked at ERVAL.

The OS audio control bits are all turned off by storing a 0 into $\$$ D208. Any bits left on from previous serial port usage are cleared by storing 3 in \$D20F.

The Atari has four sound registers (one for each voice) starting at \$D200. The first byte of each two-byte register determines the pitch (frequency). In the second byte, the four most significant bits are the distortion, and the four least significant bits are the volume.

The voice value mentioned earlier is multiplied by 2 and used as an index into the sound registers. The second value from the tokenized line is stored as the pitch in the first byte of one of the registers (\$D200, \$D202, \$D204, or \$D206), depending on the voice index. The third value from the tokenized line is multiplied by 16 and the fourth value is added to it to create the value to be stored as distortion/volume. The voice, times 2, is again used as an index to store this value in the second byte of a sound register (\$D201, \$D203, \$D205, or \$D207). The XSOUND routine then returns to Execution Control.

## External Data I/O Statements

The external data I/O statements allow data which is not part of the BASIC source program to flow into and out of BASIC. External data can come from the keyboard, a disk, or a cassette. BASIC can also create external information by sending data to external devices such as the screen, a printer, or a disk.

The INPUT and GET statements are the primary statements used for obtaining information from external devices. The PRINT and PUT statements are the primary statements for sending data to external devices.

XIO, LPRINT, OPEN, CLOSE, NOTE, POINT and STATUS are specialized I/O statements. LPRINT is used to print a single line to the " $\mathrm{P}:$ " device. The other statements assist in the I/O process.

## XINPUT (\$B316)

The execution of the INPUT statement starts at XINPUT (\$B316).
Getting the Input Line. The first action of XINPUT is to read a line of data from the indicated device. A line is any combination of up to 255 characters terminated by the EOL character (\$9B). This line will be read into the buffer located at $\$ 580$.

If the INPUT statement contained was followed by \# < expression>, the data will be read from the IOCB whose number was specified by <expression>. If there was no \# <expression>, IOCB 0 will be used. IOCB 0 is the screen editor and keyboard device (E:). If IOCB 0 is indicated, the prompt character (?) will be displayed before the input line request is made; otherwise, no prompt is displayed.
Line Processing. Once the line has been read into the buffer, processing of the data in that line starts at XINA (\$B335). The input line data is processed according to the tokens in the INPUT (or READ) statements. These tokens are numeric or string variables separated by commas.

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Processing a Numeric Variable. If the new token is a numeric variable, the CVAFP routine is called to convert the next characters in the input line to a floating point number. If this conversion does not report an error, and if the next input line character is a comma or an EOL, the floating point value is processed.

The processing of a valid numeric input value consists of calling RTNVAR to return the variable and its new value to the Variable Value Table.

If there is an error, INPUT processing is aborted via the ERRINP routine. If there is no error, but the user has hit BREAK, the process is aborted via XSTOP. If there is no abort, XINX (\$B389) is called to continue with INPUT's next task.
Processing a String Variable. If the next statement token is a string variable, it is processed at XISTR (\$B35E). This routine is also used by the READ statement. If the calling statement is INPUT, then all input line characters from the current character up to but not including the EOL character are considered to be part of the input string data. If the routine was called by READ, all characters up to but not including the next comma or EOL are considered to be part of the input string.

The process of assigning the data to the string variable is handled by calling RISASN (\$B386). If RISASN does not abort the process because of an error like DIMENSION TOO SMALL, XINX is called to continue with INPUT's next task.
XINX. The XINX (\$B389) routine is entered after each variable token in an INPUT or a READ statement is processed.

If the next token in the statement is an EOL, the INPUT/READ statement processing terminates at XIRTS (\$B3A1). XIRTS restores the line buffer pointer (\$80) to the RAM table buffer. It then restores the enter device to IOCB 0 (in case it had been changed to some other input device). Finally, XIRTS executes an RTS instruction.

If the next INPUT/READ statement token is a comma, more input data is needed. If the next input line character is an EOL, another input line is obtained. If the statement was INPUT, the new line is obtained by entering XIN0 (\$B326). If the statement was READ, the new line is obtained by entering XRD3 (\$B2D0).

The processing of the next INPUT/READ statement variable token continues at XINA.

## XGET (\$BC7F)

The GET statement obtains one character from some specified device and assigns that character to a scalar (non-array) numeric variable.

The execution of GET starts at XGET (\$BC7F) with a call to GIODVC. GIODVC will set the I/O device to whatever number is specified in the \# < expression > or to IOCB zero if no \# <expression> was specified. (If the device is IOCB 0 (E:), the user must type RETURN to force E: to terminate the input.)

The single character is obtained by calling IO3. The character is assigned to the numeric variable by calling ISVAR1 (\$BD2D). ISVAR1 also terminates the GET statement processing.

## PRINT

The PRINT statement is used to transmit text data to an external device. The arguments in the PRINT statement are a list of numeric and/or string expressions separated by commas or semicolons. If the argument is numeric, the floating point value is converted to text form. If the argument is a string, the string value is transmitted as is.

If an argument separator is a comma, the arguments are output in tabular fashion: each new argument starts at the next tab stop in the output line, with blanks separating the arguments.

If the argument separator is a semicolon, the transmitted arguments are appended to each other without separation.

The transmitted line is terminated with an EOL, unless a semicolon or comma directly precedes the statement's EOL or statement separator (:).
XPRINT (\$B3B6). The PRINT routine begins at XPRINT (\$B3B6). The tab value is maintained in the PTABW (\$C9) cell. The cell is initialized with a value of ten during BASIC's cold start, so that commas in the PRINT line cause each argument to be displaced ten positions after the beginning of the last argument. The user may POKE PTABW to set a different tab value.

XPRINT copies PTABW to SCANT (\$AF). SCANT will be used to contain the next multiple-of-PTABW output line displacement - the column number of the next tab stop.

COX is initialized to zero and is used to maintain the current output column or displacement.

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XPRO. XPRINT examines the next statement token at XPR0 (\$B3BE), classifies it, and executes the proper routine.
\# Token. If the next token is \#, XPRIOD (\$B437) is entered. This routine modifies the list device to the device specified in the \# < expression > . XPR0 is then entered to process the next token.
, Token. The XPTAB (\$B419) routine is called to process the , token. Its job is to tab to the next tab column.

If COX (the current column) is greater than SCANT, we must skip to the next available tab position. This is done by continuously adding PTABW to SCANT until COX is less than or equal to SCANT. When COX is less than SCANT, blanks (\$20) are transmitted to the output device until COX is equal to SCANT.

The next token is then examined at XPRO.
EOL and : Tokens. The XPEOS (\$B446) routine is entered for EOL and : tokens. If the previous token was a ; or , token, PRINT exits at XPRTN (\$B458). If the previous token was not a; or, token, an EOL character is transmitted before exiting via XPRTN.
; Token. No special action is taken for the ; token except to go to XPR0 to examine the next token.
Numbers and Strings. If the next token is not one of the above tokens, Execute Expression is called to evaluate the expression. The resultant value is popped from the argument stack and its type is tested for a number or a string.

If the argument popped was numeric, it will be converted to text form by calling CVFASC. The resulting text is transmitted to the output device from the buffer pointed to by INBUFF (\$F3). XPR0 is then entered to process the next token.

If the argument popped was a string, it will be transmitted to the output device by the code starting at :XPSTR (\$B3F8). This code examines the argument parameters to determine the current length of the string. When the string has been transmitted, XPR0 is entered to process the next token.

## XLPRINT (\$B464)

LPRINT, a special form of the PRINT statement, is used to print a line to the printer device (P:).

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The XLPRINT routine starts at \$B464 by opening IOCB 7 for output to the P: device. XPRINT is then called to do the printing. When the XPRINT is done, IOCB 7 is closed via CLSYS1 and LPRINT is terminated.

## XPUT (\$BC72)

The PUT statement sends a single byte from the expression in the PUT statement to a specified external device.

Processing starts at XPUT (\$BC72) with a call to GIODVC. GIODVC sets the I/O device to the IOCB specified in \# < expression > . If a \# < expression > does not exist, the device will be set to IOCB zero (E:).

The routine then calls GETINT to execute PUT's expression and convert the resulting value to a two-byte integer. The least significant byte of this integer is then sent to the PUT device via PRCX. PRCX also terminates the PUT processing.

## XXIO (\$BBE5)

The XIO statement, a general purpose I/O statement, is intended to be used when no other BASIC I/O statement will serve the requirements. The XIO parameters are an IOCB I/O command, an IOCB specifying expression, an AUX1 value, an AUX2 value, and finally a string expression to be used as a filespec parameter.

XIO starts at XXIO (\$BBE5) with a call to GIOCMD. GIOCMD gets the IOCB command parameter. XIO then continues at XOP1 in the OPEN statement code.

## XOPEN (\$BBEB)

The OPEN statement is used to open an external device for input and/or output. OPEN has a \# < expression > , the open type parameter (AUX1), an AUX2 parameter, and a string expression to be used as a filespec.

OPEN starts at XOPEN at \$BBEB. It loads the open command code into the A register and continues at XOP1.
XOP1. XOP1 continues the OPEN and XIO statement processing. It starts at $\$$ BBED by storing the A register into the IOCMD cell. Next it obtains the AUX1 (open type) and AUX2 values from the statement.

The next parameter is the filespec string. In order to insure that the filespec has a proper terminator, SETSEOL is called to place a temporary EOL at the end of the string.

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The XIO or OPEN command is then executed via a call to IO1. When IO1 returns, the temporary EOL at the end of the string is replaced with its previous value by calling RSTSEOL.

OPEN and XIO terminate by calling IOTEST to insure that the command was executed without error.

## XCLOSE (\$BC1B)

The CLOSE statement, which closes the specified device, starts at XCLOSE (\$BC1B). It loads the IOCB close command code into the A register and continues at GDVCIO.
GDVCIO. GDVCIO (\$BC1D) is used for general purpose device I/O. It stores the A register into the IOCMD cell, calls GIODVC to get the device from \# < expression > , then calls IO7 to execute the I/O. When IO7 returns, IOTEST is called to test the results of the I/O and terminate the routine.

## XSTATUS (\$BC28)

The STATUS statement executes the IOCB status command. Processing starts at XSTATUS (\$BC28) by calling GIODVC to get the device number from \# < expression> . It then calls IO8 with the status command in the A register. When IO8 returns, the status returned in the IOCB status cell is assigned to the variable specified in the STATUS statement by calling ISVAR1. ISVAR1 also terminates the STATUS statement processing.

## XNOTE (\$BC36)

The NOTE statement is used specifically for disk random access. NOTE executes the Disk Device Dependent Note Command, $\$ 26$, which returns two values representing the current position within the file for which the IOCB is open. NOTE begins at XNOTE at $\$$ BC36. The code loads the command value, $\$ 26$, into the A register and calls GDVCIO to do the I/O operation. When GDVCIO returns, the values are moved from AUX3 and AUX4 to the first variable in the NOTE statement. The next variable is assigned the value from AUX5.

## XPOINT (\$BC4D)

The POINT statement is used to position a disk file to a previously NOTEd location. Processing starts at XPOINT (\$BC4D). This routine converts the first POINT parameter to an integer and stores the value in AUX3 and AUX4. The second parameter is then converted to an integer and its value stored
in AUX5. The POINT command, $\$ 25$, is executed by calling GDIO1, which is part of GDVCIO.

## Miscellaneous I/O Subroutines

IOTEST. IOTEST (\$BCB3) is a general purpose routine that examines the results of an I/O operation. If the I/O processing has returned an error, IOTEST processes that error.

IOTEST starts by calling LDIOSTA to get the status byte from the IOCB that performed the last I/O operation. If the byte value is positive (less than 128), IOTEST returns to the caller.

If the status byte is negative, the I/O operation was abnormal and processing continues at SICKIO.

If the I/O aborted due to a BREAK key depression, BRKBYT (\$11) is set to zero to indicate BREAK. If a LOAD was in progress when BREAK was hit, exit is via COLDSTART; otherwise IOTEST returns to its caller.

If the error was not from IOCB 7 (the device BASIC uses), the error status value is stored in ERRNUM and ERROR is called to print the error message and abort program execution.

If the error was from IOCB 7, then IOCB 7 is closed and ERROR is called with the error status value in ERRNUM unless ENTER was being executed, and the error was an end-of-file error. In this case, IOCB 7 is closed, the enter device is reset to IOCB 0, and SNX2 is called to return control to the Program Editor.
I/O Call Routine. All I/O is initiated from the routine starting at IO1 (\$BD0A). This routine has eight entry points, IO1 through IO8, each of which stores predetermined values in an IOCB. All IOn entry points assume that the X register contains the IOCB value, times 16 .

IO1 sets the buffer length to 255 .
IO2 sets the buffer length to zero.
IO3 sets the buffer length to the value in the Y register plus a most-significant length byte of zero.

IO4 sets the buffer length from the values in the $\mathrm{Y}, \mathrm{A}$ register pair, with the A register being the most-significant value.

IO5 sets the buffer address from the value in the INBUFF cell (\$F3).

IO6 sets the buffer address from the Y,A register pair. The A register contains the most significant byte.

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IO7 sets the I/O command value from the value in the IOCMD cell.

IO8 sets the I/O command from the value in the A register.
All of this is followed by a call to the operating system CIO entry point. This call executes the I/O. When CIO returns, the general I/O routine returns to its caller.

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# Internal I/O Statements 

The READ, DATA, and RESTORE statements work together to allow the BASIC user to pass predetermined information to his or her program. This is, in a sense, internal I/O.

## XDATA (\$A9E7)

The information to be passed to the BASIC program is stored in one or more DATA statements. A DATA statement can occur any place in the program, but execution of a DATA statement has no effect.

When Execution Control encounters a DATA statement, it expects to process this statement just like any other. Therefore an XDATA routine is called, but XDATA simply returns to Execution Control.

## XREAD (\$B283)

The XREAD routine must search the Statement Table to find DATA. It uses Execution Control's subroutines and line parameters to do this. When XREAD is done, it must restore the line parameters to point to the READ statement. In order to mark its place in the Statement Table, XREAD calls a subroutine of XGOSUB to put a GOSUB-type entry on the Runtime Stack.

The BASIC program may need to READ some DATA, do some other processing, and then READ more DATA. Therefore, XREAD needs to keep track of just where it is in which DATA statement. There are two parameters that provide for this. DATALN (\$B7) contains the line number at which to start the search for the next DATA statement. DATAD (\$B6) contains the displacement of the next DATA element in the DATALN line. Both values are set to zero as part of RUN and CLR statement processing.

XREAD calls Execution Control's subroutine GETSTMT to get the line whose number is stored in DATALN. If this is the first READ in the program and a RESTORE has not set a

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different line number, DATALN contains zero, and GETSTMT will get the first line in the program. On subsequent READs, GETSTMT gets the last DATA statement that was processed by the previous READ.

After getting its first line, XREAD calls the XRTN routine to restore Execution Control's line parameters.

The current line number is stored in DATALN. XREAD steps through the line, statement by statement, looking for a DATA statement. If the line contains no DATA statement, then subsequent lines and statements are examined until a DATA statement is found.

When a DATA statement has been found, XREAD inspects the elements of the DATA statement until it finds the element whose displacement is in DATAD.

If no DATA is found, XREAD exits via the ERROOD entry point in the Error Handling Routine. Otherwise, a flag is set to indicate that a READ is being done, and XREAD joins XINPUT at :XINA. XINPUT handles the assignment of the DATA values to the variables. (See Chapter 13.)

## XREST (\$B26B)

The RESTORE statement allows the BASIC user to re-READ a DATA statement or change the order in which the DATA statements are processed. The XREST routine simulates RESTORE.

XREST sets DATALN to the line number given, or to zero if no line number is specified. It sets DATAD to zero, so that the next READ after a RESTORE will start at the first element in the DATA line specified in DATALN.

# Miscellaneous Statements 

## XDEG (\$B261) and XRAD (\$B266)

The transcendental functions such as SIN or COS will work with either degrees or radians, depending on the setting of RADFLG (\$FB). The DEG and RAD statements cause RADFLG to be set. These statements are simulated by the XDEG and XRAD routines, respectively.

The XDEG routine stores a six in RADFLG. XRAD sets it to zero. These particular values were chosen because they aid the transcendental functions in their calculations.

RADFLG is set to zero during BASIC's initialization process and also during simulation of the RUN statement.

## XPOKE (\$B24C)

The POKE statement is simulated by the XPOKE routine. XPOKE calls a subroutine of Execute Expression to get the address and data integers from the tokenized line. XPOKE then stores the data at the specified address.

## XBYE (\$A9E8)

The XBYE routine simulates the BYE statement. XBYE closes all IOCBs (devices and files) and then jumps to location \$E471 in the Operating System. This ends BASIC and causes the memo pad to be displayed.

## XDOS (\$A9EE)

The DOS statement is simulated by the XDOS routine. The XDOS routine closes all IOCBs and jumps to whatever address is stored in location $\$ 0 \mathrm{~A}$. This will be the address of DOS if DOS has been loaded. If DOS has not been loaded, \$0A will point to the memo pad.

## XLET (\$AAE0)

The LET and implied LET statements assign values to variables. They both invoke the XLET routine, which consists of the Execute Expression routines. (See Chapter 7.)

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## XREM (\$A9E7)

The REM statement is for documentation purposes only and has no effect on the running program. The routine which simulates REM, XREM, simply executes an RTS instruction to return to Execution Control.

## XERR (\$B91E)

When a line containing a syntax error is entered, it is given a special statement name token to indicate the error. The entire line is flagged as erroneous no matter how many previously good statements are in the line. The line is then stored in the Statement Table.

The error statement is processed just like any other. Execution Control calls a routine, XERR, which is one of the entry points to the Error Handling Routine. It causes error 17 (EXECUTION OF GARBAGE).

## XDIM (\$B1D9)

The DIMension statement, simulated by the XDIM routine, reserves space in the String/Array Table for the DIMensioned variable.

The XDIM routine calls Execute Expression to get the variable to be DIMensioned from the Variable Value Table. The variable entry is put into a work area. In the process, Execute Expression gets the first and second DIMension values and sets a default of zero if only one value is specified.

XDIM checks to see if the variable has already been
DIMensioned. If the variable was already DIMensioned, XDIM exits via the ERRDIM entry point in the Error Handling Routine. If not, a bit is set in the variable type byte in the work area entry to mark this variable as DIMensioned.

Next, XDIM calculates the amount of space required. This calculation is handled differently for strings and arrays.
DIMensioning an Array. XDIM first increments both dimension values by one and then multiplies them together to get the number of elements in the array. XDIM multiplies the result by 6 (the length of a floating point number) to get the number of bytes required. EXPAND is called to expand the String/Array Table by that amount.

XDIM must finish building the variable entry in the work area. It stores the first and second dimension values in the entry. It also stores the array's displacement into the

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String/Array Table. It then calls an Execute Expression subroutine to return the variable to the Variable Value Table. (See Chapter 3.)
DIMensioning a String. Reserving space for a string in the String/Array Table is much simpler. XDIM merely calls the EXPAND routine to expand by the user-specified size.

XDIM must also build the Variable Value Table entry in the work area. It sets the current length to 0 and the maximum length to the DIMensioned value. The displacement of the string into the String/Array Table is also stored in the variable. XDIM then calls a subroutine of Execute Expression to return the variable entry to the Variable Value Table. (See Chapter 3.)

## Initialization

When the Atari computer is powered up with the BASIC cartridge in place, the operating system does some processing and then jumps to a BASIC routine. Between the time that BASIC first gets control and the time it prints the READY message, initialization takes place. This initialization is called a cold start. No data or tables are preserved during a cold start.

Initialization is repeated if things go terribly awry. For example, if there is an I/O error while executing a LOAD statement, BASIC is totally confused. It gives up and begins all over again with the COLDSTART routine.

Sometimes a less drastic partial initialization is necessary. This process is handled by the WARMSTART routine, in which some tables are preserved.

Entering the NEW statement, simulated by the XNEW routine, has almost the same effect as a cold start.

## COLDSTART (\$A000)

Two flags, LOADFLG and WARMFLG, are used to determine if a cold or warm start is required.

The load flag, LOADFLG (\$CA), is zero except during the execution of a LOAD statement. The XLOAD routine sets the flag to non-zero when it starts processing and resets it to zero when it finishes. If an I/O error occurs during that interval, IOTEST notes that LOADFLG is non-zero and jumps to COLDSTART.

The warm-start flag, WARMFLG (\$08), is never set by BASIC. It is set by some other routine, such as the operating system or DOS. If WARMFLG is zero, a cold start is done. If it is non-zero, a warm start is done. During its power-up processing, before BASIC is given control, OS sets WARMFLG to zero to request a cold start. During System Reset processing, OS sets the flag to non-zero, indicating a warm start is desired.

If DOS has loaded any data into BASIC's program area during its processing, it will request a cold start.

The COLDSTART routine checks both WARMFLG and LOADFLG to determine whether to do a cold or warm start. If a cold start is required, COLDSTART initializes the 6502 CPU

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stack and clears the decimal flag. The rest of its processing is exactly the same as if the NEW statement had been entered.

## XNEW (\$A00C)

The NEW statement is simulated by the XNEW routine. XNEW resets the load flag, LOADFLG, to zero. It initializes the zeropage pointers to BASIC's RAM tables. It reserves 256 bytes at the low memory address for the multipurpose buffer and stores its address in the zero-page pointer located at $\$ 80$. Since none of the RAM tables are to retain any data, their zeropage pointers ( $\$ 82$ through $\$ 90$ ) are all set to low memory plus 256.

The Variable Name Table is expanded by one byte, which is set to zero. This creates a dummy end-of-table entry.

The Statement Table is expanded by three bytes. The line number of the direct statement ( $\$ 8000$ ) is stored there along with the length (three). This marks the end of the Statement Table.

A default tab value of 10 is set for the PRINT statement.

## WARMSTART (\$A04D)

A warm start is the least drastic of the three types of initialization. Everything the WARMSTART routine does is also done by COLDSTART and XNEW.

The stop line number (STOPLN), the error number (ERRNUM), and the DATA parameters (DATALN and DATAD) are all set to zero. The RADFLG flag is set to zero, indicating that transcendental functions are working in radians. The break byte (BRKBYT) is set off and \$FF is stored in TRAPLN to indicate that errors are not being trapped.

All IOCBs (devices and files) are closed.
The enter and list devices (ENTDTD and LISTDTD) are set to zero to indicate the keyboard and the screen, respectively.

Finally, the READY message is printed and control passes to the Program Editor.

Part Two

## Directly <br> Accessing <br> Atari BASIC

# Introduction to Part Two 

Congratulations! If you have read all of Part 1, you are through the hard stuff. In Part 2, we hope to teach you how to use at least some of the abundance of information presented in the Source Listing and in Part 1. In particular, we will show you how to examine the various RAM and ROM tables used by BASIC.

The examples and suggestions will be written in Atari BASIC. But those of you who are true-blue assembly language fanatics should have little trouble translating the concepts to machine code, especially with the source listing to guide you.

Would that we could present an example program or concept for each possible aspect of the BASIC interpreter, but space does not allow it - nor would it be appropriate. For example, although we will present here a program to list all keywords and token values used by BASIC, we will not explore the results (usually disastrous) of changing token values within a BASIC program.

Part 2 begins with a pair of introductory chapters. If you are experienced at hexadecimal-to-decimal conversions and with the concepts of word and byte PEEKs and POKEs, you may wish to skip directly to Chapter 3 .

# Hexadecimal Numbers 

The word hexadecimal means, literally, " of six and ten. ${ }^{\text {' }}$ It implies, however, a number notation which uses 16 as its base instead of 10. Hexadecimal notation is used as a sort of shorthand for the eight-digit binary numbers that the 6502 understands. If Atari BASIC understood hexadecimal numbers and we all had eight fingers on each hand, there would be no need for this chapter. Instead, to use this book you have to make many conversions back and forth between hexadecimal ("hex") and decimal notation. Many BASIC users have never had to learn that process.

Virtually all the references to addresses and other values in this book are given in hexadecimal notation (or simply "hex" to us insiders). For example, we learn that the Atari BASIC ROM cartridge has $\$ A 000$ for its lowest address and that location $\$ 80$ contains a pointer to BASIC's current LOMEM. But what does all that mean?

First of all, if you are not familiar with 6502 assembly language, let me point out that there is a convention that a number preceded by a dollar sign ( $\$ 80$ ) is a hexadecimal number, even if it contains only decimal digits. Also, notice that in the Source Listing all numbers in the first three columns are hexadecimal, even though the dollar sign is not present. (To the right of those columns, though, only those numbers preceded by a dollar sign are in hex.)

Now, suppose I wanted to look at the contents of location \$A4AF (SNTAB in the listing). Realistically, the only way to look at a memory location from BASIC is via the PEEK function (and see the next chapter if you are not sure how to use PEEK in this situation). But BASIC's language syntax requires a decimal number with PEEK - for instance, PEEK (15).

Obviously, we need some way to convert from hexadecimal to decimal. Aside from going out and buying one of the calculators made just for this purpose, the best way is probably to let your computer help you. And the computer can help you

## Chapter One

even if you only understand BASIC. As an example, here's a BASIC program that will convert hex to decimal notation:

```
1Ø DIM HEX$ (23),NUM$ (4)
2\emptyset HEX$="@ABCDEFGHI#######JKLMNO"
3\emptyset CVHEX=9\emptyset\emptyset\emptyset
IØ\varnothing PRINT :PRINT "GIVE ME A HEX NUMBER ";
llØ INPUT NUM$
12\emptyset GOSUB CVHEX
13\emptyset PRINT "HEX ";NUMS;" = DECIMAL ";NUM
14\emptyset GOTO 1Ø\emptyset
9Ø\emptyset\emptyset REM THE CONVERT HEX TO DECIMAL ROUTINE
901\emptyset NUM=\emptyset
902\emptyset FOR I=1 TO LEN(NUM$)
9030 NUM=NUM*16+ASC(HEX$(ASC(NUM$(I))-47))-64
9\emptyset4\emptyset NEXT I:RETURN
```

Now, while this program might be handy for a few purposes, it would be much neater if we could simply use its capabilities anytime we wanted to examine or change a location (or its contents) referred to by a hex address or data. And so shall it be used.

If we remove lines 100 through 140, inclusive, then any BASIC program which incorporates the rest of the program may change a hex number into decimal by simply

1. placing the ATASCII form of the hex number in the variable NUM\$,
2. calling the convert routine at line 9000 (via GOSUB CVHEX), and
3. using the result, which is returned in the variable NUM.

In the next chapter, we will immediately begin to make use of this routine. If you are not used to hex notation, you might do well to type in and play with this program before proceeding.

Finally, before we leave this subject, let's examine a routine which will allow us to go the other way - that is, convert decimal to hex:

```
4\emptyset DIM DEC$(16):DEC$="\emptyset123456789ABCDEF"
5\emptyset CVDEC=91Ø\emptyset
IØ\emptyset PRINT :PRINT "GIVE ME A DECIMAL NUMBER ";
```


## Chapter One

```
11\emptyset INPUT DEC:NUM=DEC
12\emptyset GOSUB CVDEC:REM 'NUM' is destroyed by this
13\emptyset PRINT DEC;" Decimal = ";NUM$;" Hex"
14Ø GOTO 1Ø\emptyset
91Ø\emptyset REM CONVERT DECIMAL TO HEX ROUTINE
911\emptyset DIV=4Ø96
912\emptyset FOR I=1 TO 4
9130 N=INT(NUM/DIV):NUM$(I,I)=DEC$(N+1)
9140 NUM=NUM-DIV*N:DIV=DIV/16
9150 NEXT I
9160 RETURN
```

These lines are meant to be added to the previous program, though they can be used alone if you simply add this line:

10 DIM NUM\$(4)
We will use portions of these programs in later chapters, but we may compress some of the code into fewer lines simply to save wear and tear on our fingers. If you study these routines, you'll recognize them in their transformed versions.

## PEEKing and POKEing

In contrast to languages which include direct machine addressing capability, like " C " and Forth, and in contrast to "school" languages like Pascal and Fortran, which specifically prevent such addressing, BASIC provides a sort of halfway measure in machine accessibility.

POKE is a BASIC statement. Its syntax is POKE <address > , < data > . Naturally, both <address> and <data> may be constants, variables, or even full-blown expressions:

POKE 82,0: REM change left screen margin to zero produces the same result as

LEFTMARGIN $=82:$ POKE LEFTMARGIN, 0
PEEK, on the other hand, is a BASIC function. It cannot stand alone as a statement. To use PEEK, we either PRINT the value (contents) of a PEEKed location, assign a PEEKed value to a variable, or test the value for some condition:

POKE 82, PEEK(82) +1 : REM move the left margin in a space

PRINT PEEK(106) : REM where is the top of system memory?

IF PEEK(195) = 136 THEN PRINT "End of File"
In the first example, the number POKEd into 82 will be whatever number was stored before, plus 1. As explained in Part 1, the PEEK function is executed before the POKE.

An aside: Just where did I get those addresses I used in the PEEKs and POKEs? One way to find them is to peruse the listings of Atari's operating system, available in Atari's technical manuals set, and the listing of BASIC in this book. Another way would be to use a book (like COMPUTE! Books' Mapping the Atari) or a reference card designed specifically to tell you about such addresses.

And one more thing to consider before moving on. If we counted all of the bit patterns possible in a single 8 -bit byte (like

01010101, 11110000, and 00000001, where each 1 or 0 represents a single on or off bit), we would discover that there are 256 unique combinations, ranging in value from 0 to 255 . Since each memory location can hold only one byte, it is not surprising to learn that the PEEK function will always return a number from 0 to 255 ( $\$ 00$ to $\$ F F$ ). Similarly, BASIC will only POKE a data value that is an integer from 0 to 255. In fact, BASIC will convert any data to be POKEd to an integer number, rounding off any fractional parts.

So far so good. But suppose we want to examine a location which is actually a two-byte word, such as the line number where the last TRAPped error occurred, stored starting at location \$BA hex or 186 decimal. PEEK only lets us look at one byte at a time. How do we look at two bytes? Simple: one byte at a time.

In most cases, words in a 6502-based machine are stored in memory with the least significant byte stored first. This means that the second byte of each word is a count of the number of 256's there are in its value, and the first byte is the leftovers. (Or we can more properly say that the first byte contains "the word's value modulo 256.") Confused? Let's try restating that.

In decimal arithmetic, we can count from 0 to 9 in a single digit. To go beyond 9, we have a convention that says the digit second from the right represents the number of 10 's in the number, and so on.

If we consider bytes to be a computer's digits, which in many ways they are, and if we remember that each byte may represent any number from 0 to 255 (or $\$ 00$ to $\$ F F$ ), then it is logical to say that the next byte is a count of the number of 256 's in the number. The only thing illogical is that the higher byte comes after the lower byte (like reading 37 as " 7 tens and 3 ones' ${ }^{\prime}$ instead of what we are used to).

Some examples might help:

| a 6502 word <br> in memory | as written <br> in assembler | think of <br> it as | decimal <br> value |
| :---: | :---: | :---: | ---: |
| 0100 | $\$ 0001$ | $0 * 256+1$ | 1 |
| 0001 | $\$ 0100$ | $1 * 256+0$ | 256 |
| 0204 | $\$ 0402$ | $4^{*} 256+2$ | 1026 |
| FF FF | $\$ F F F F$ | $255^{*} 256+255$ | 65535 |

So let's examine that error line location: PRINT PEEK(186) + 256 * PEEK(187)

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Do you see it? Since the second byte is a count of the number of 256 's in the value, we must multiply it by 256 to calculate its true value.

Now, in the case of line numbers, it is well and good that we print out a decimal value, since that is how we are used to thinking of them. But suppose you wished to print out some of BASIC's tables? You might very well wish to see the hex representations. The program presented here allows you to specify a hex address. It then presents you with the contents of the byte and the word found at that address, in both decimal and hex form.

```
10 DIM HEX$(23),NUM$(4)
2\emptyset HEX$="@ABCDEFGHI#######JKLMNO"
3\varnothing CVHEX=9øøø
4Ø DIM DEC$(16):DEC$="Ø123456789ABCDEF"
5\varnothing CVDEC=91ø\emptyset
1\emptyset\emptyset PRINT :PRINT "WHAT ADDRESS TO VIEW ";
11\emptyset INPUT NUM$:PRINT
12Ø PRINT "Address ";NUM$;" contains:"
13\emptyset GOSUB CVHEX:ADDR=NUM
14ø NUM=PEEK(ADDR):GOSUB CVDEC
15\emptyset PRINT ,"byte ";PEEK(ADDR);" = $";NUM$(3)
16\emptyset WORD=PEEK(ADDR)+256*PEEK(ADDR+1)
17\emptyset NUM=WORD:GOSUB CVDEC
18Ø PRINT ,"word ";WORD;" = $";NUM$
19Ø GOTO 1ø\emptyset
9øø\emptyset REM THE CONVERT HEX TO DECIMAL ROUTINE
901\varnothing NUM=\varnothing
9ø2\emptyset FOR I=1 TO LEN(NUM$)
903ø NUM=NUM*16+ASC(HEX$(ASC(NUM$(I))-47))-64
9040 NEXT I:RETURN
91ø\emptyset REM CONVERT DECIMAL TO HEX ROUTINE
911\varnothing DIV=4096
912\emptyset FOR I=l TO 4
9130 N=INT(NUM/DIV):NUM$(I,I)=DEC$(N+1)
9140 NUM=NUM-DIV*N:DIV=DIV/l6
915\emptyset NEXT I
9160 RETURN
```

You may have noticed that lines 10 through 50 and lines 9000 to the end are the same as those used in the example

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programs in the last chapter. And did you see line 160, where we obtained the word value by multiplying by 256 ?

As the last point of this chapter, we need to discuss how to change a word value. Obviously, in Atari BASIC we can't POKE both bytes of a word at once any more than we could retrieve both bytes at once (although BASIC A + can, by using the DPOKE statement and DPEEK function). So we must invent a mechanism to do a double POKE.

Given that the variable ADDR contains the address at which we wish to POKE a word, and given that the variable WORD contains the value (in decimal) of the desired word, the following code fragment will perform the double POKE:

POKE ADDR + 1,INT(WORD/256)
POKE ADDR,WORD-256*PEEK(ADDR + 1)
This is kind of sneaky code, but calculating the most significant byte and POKEing the value in byte location ADDR +1 first allows us to also use it as a kind of temporary variable in calculating the least significant byte. By PEEKing the location that already holds the high-order byte, we can subtract it from the original value. The remainder is WORD modulo 256 - the low-order byte.

And that's about it. Hopefully, if you were not familiar with PEEK and POKE before, you now at least will not approach their use with too much caution. Generally, PEEKs will never harm either your running program or the machine, but don't be surprised if a stray POKE or two sends your computer off into never-never land. After all, you may have just told BASIC to start putting your program into ROM, or worse.

On the other hand, if you have removed your diskettes and turned off your cassette recorder, the worst that can happen from an erring POKE is that you'll have to turn the power off and back on again. So have at it. Happy PEEKing and POKEing.

# Listing Variables in Use 

Chapter 3 of Part 1 described the layout of the Variable Name Table and the Variable Value Table. In particular, we read that the Variable Name Table was built in a very simple fashion: Each new variable name, as it is encountered upon program entry, is simply added to the end of the list of names. The most significant bit of the last character of the name is turned on, to signal the end of that name. The contents of VNTP point to the beginning of the list of names, and the content of VNTD is the address of the byte after the end of the list.

Now, what does all that mean? What does it imply that we can do? Briefly, it implies that we can look at BASIC's memory and find out what variable names are in current use. Here's a program that will do exactly that:

| 32700 | Q |
| :---: | :---: |
| 32710 | FOR Q=PEEK (13Ø) +256*PEEK (131) TO PE EK (132) +256*PEEK(133)-1 |
| 32720 | IF PEEK $(Q)<128$ THEN PRINT CHR\$ (PEEK (Q)) ;: NEXT Q:STOP |
| 32730 | PRINT CHR\$ (PEEK (Q)-128): QQ=QQ+l:PRI NT QQ,:NEXT Q:STOP |

Actually, this is not so much a program as it is a program fragment. It is intended that you will type NEW, type in the above fragment, and then LIST the fragment to a disk file (LIST "D:LVAR") or to a cassette (LIST "C:'"). Then type NEW again and ENTER or LOAD the program whose variables you want to list. Finally, use ENTER to re-enter the fragment from disk (ENTER "D:LVAR") or cassette (ENTER "C:'"). Then type GOTO 32700 to obtain your Variable Name Table listing.

Of course, if you had OPENed a channel to the printer (OPEN \#1,8,0, "P:'"), you could change the PRINTs to direct the listing to the printer (PRINT \#1; CHR\$ (<expression > ) ).

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How does the fragment work? The reason for the start and end limits for the FOR loop are simple: word location 130 (\$82) contains the pointer to the beginning of the Variable Name Table and word location 132 (\$84) contains the pointer to the end of that same table, plus 1 . So we simply traipse through that table, printing characters as we encounter them - except that when we encounter a character with its most significant bit on (IF PEEK $(Q)>127$ ), we turn off that bit before printing it and start the next name on a new line.

Notice that we use the variable QQ to allow us to print out the token value for each variable name. We will use this information in some later chapters.

Also note that the variable names $Q Q$ and $Q$ will appear in your variable name listing. Sorry. We can write a program which would accomplish the same thing without using variables, but it would be two or three times as big and much harder to understand. Of course, if you consistently use certain variable names, such as $I$ and $J$ in FOR-NEXT loops, you could use those names here instead, thus not affecting the count of variables in use.

Incidentally, the STOP at the end of the third line should be unnecessary, since the table is supposed to end with a character with its upper bit on. But I've learned not to take chances - things don't always go as they're supposed to.

## Variable Values

In this chapter, we will show how you can determine the value of any variable by inspecting the Variable Value Table. Actually, in many respects this is a waste of effort. After all, if I need to know the value of the variable TOTAL, I can just type PRINT TOTAL.

But this book is supposed to be a guide, and there are a few uses for this information, particularly in assembly language subroutines, and it is instructive in that it gives us an inkling of what BASIC goes through to evaluate a variable reference.

It will probably be better to present the program first, and then explain what it does. Before doing so, though, note that the program fragment expects you to give it a valid variable token (128 through 255). No checks are made on the validity of that number, since we are all intelligent humans here and since we want to save program space. Enough. The program:

| Ø | PRINT : PRINT "WHAT VARIABLE NU ";:INPUT Q |
| :---: | :---: |
| 32505 | ```Q=PEEK(134)+256*PEEK(135)+(Q-128)* 8``` |
| 32510 | PRINT : PRINT "VARIABLE NUMBER ";PE EK (Q+1), |
| 32515 | ON INT(PEEK (Q)/64) GOTO 326øø,3265 Ø |
| 32 | PRINT "IS A NUMBER, ":PRINT ,"VALU E"; |
| 32525 | QEXP=PEEK $(Q+2): I F$ QEXP $>127$ THEN PR INT "-";:QEXP=QEXP-128 |
| 32530 | QNUM=Ø:FOR $\mathrm{QQ}=\mathrm{Q}+3 \mathrm{TO} \mathrm{Q}+7$ |
| 32535 | ```QNUM=QNUM*I\emptyset\emptyset+PEEK(QQ)-6*INT( QQ)/16):NEXT QQ``` |
| 32540 | QEXP=QEXP-68:IF QEXP=Ø THEN 32555 |
| 32545 | FOR QQ=QEXP TO SGN(QEXP) STEP -SGN (QEXP) |

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| 32550 | $\begin{aligned} & \text { QNUM }=(\text { QEXP }>\varnothing) * \\ & \text { UM/ } \varnothing \varnothing: \text { NEXT } Q Q \end{aligned}$ |
| :---: | :---: |
| 32555 | PRINT QNUM:PRINT : GOTO 325øø |
| 32570 | IF PEEK (Q)/2 < > INT (PEEK (Q)/2) THEN 32580 |
| 32575 | PRINT ,"AND IS NOT YET DIMENSIONED ": POP:GOTO 325øø |
| 32580 | PRINT ,"ADDRESS IS ";PEEK (Q+2)+256 <br> *PEEK ( $\mathrm{Q}+3$ ) : RETURN |
| 32600 | PRINT "IS AN ARRAY, ":GOSUB 32570 |
| 32610 | PRINT , "DIM 1 IS "; PEEK (Q+4)+256*P EEK (Q+5) |
| 32615 | PRINT ,"DIM 2 IS ";PEEK (Q+6)+256*P EEK (Q+7) |
| 32620 | GOTO 325øø |
| 32650 | PRINT "IS A STRING, ":GOSUB 3257Ø |
| 32660 | PRINT , "LENGTH IS ";PEEK (Q+4)+256* PEEK ( $Q+5$ ) |
| 32665 | PRINT ,"\{3 SPACES\}DIM IS ";PEEK(Q+ $6)+256 * \operatorname{PEEK}(Q+7)$ |

Did you get lost in all of that? I got lost several times as I wrote it, but it seems to work well. Shall we discuss it?

The first place where confusion may arise is when I ask you to give a variable token from 128 to 255 , and then reveal that the entry in the Variable Value Table thinks variable numbers range from 0 to 127 . Actually, there is no anomaly here. The variable token that you input is the token value of the variable in your program. The number in the table is its relative position. The numbers differ only in their uppermost bit.

The program uses the number you specify to form an address of an entry somewhere within the Variable Value Table. It then displays the internal variable number and examines the flag byte of the variable entry. Recall that the uppermost bit ( $\$ 80$, or 128 ) of the flag byte is on, if this variable is a string. The next bit ( $\$ 40$, or 64 ) is on if the variable is an array. If neither is on, the variable is a normal floating point number (or scalar, as it is sometimes called, to distinguish it from a floating point array). All this is decided and acted upon in line 32515.

Before examining what happens if the number is a scalar, let's look at strings and arrays. Both start out (lines 32600 and 32650) by identifying themselves and calling a subroutine which determines if the variable has been DIMensioned yet. If not, the subroutine tells us so, removes the GOSUB entry from the stack, and starts the whole shebang over again. If the variable is DIMensioned, though, we print its address before returning. Note that the address printed is the relative address within the String/Array Table.

If the DIMension check subroutine returns, both string and array variables have their vitals printed out before the program asks you for another variable number. In the case of a string, we see the current length (as would be obtained by the LENgth function) and its dimension. For an array, we see both dimensions. Note that array dimensions here are always one greater than the user program specified, so that a zero dimension value means "this dimension is unused."

Point of interest: this program will never print a zero for an array dimension. Why? Because Atari BASIC never places a zero in either dimension when the DIM statement is executed. In a way, this is a "feature" (a feature is a documented bug). It implies that we may code DIM $X X(7)$ and yet use something like PRINT XX(N,0). In other words, a singly dimensioned array in Atari BASIC is exactly equivalent to a doubly dimensioned array with a 0 as the second subscript in the DIM statement.

Back to the listing. Fairly straightforward up until now. But look what happens if the variable is a scalar, a single floating point number.

First, we obtain the exponent byte; if its upper bit is on, the number is negative, so we print the minus sign before turning the bit off.

Second, we must loop through the five bytes of the mantissa, accumulating a value. The really strange part here is line 32535 , so let's examine it closely. As we get each byte, we must multiply what we have gotten so far by 100 (remember, floating point numbers are in BCD format, so each byte represents a power of 100). Then, what we really want to do is add in 10 times the higher digit in the byte, plus the lower digit. We could have gotten those numbers as follows:

NEWBCDVALUE = OLDBCDVALUE*100
HIGHER $=\operatorname{INT}($ PEEK $(Q Q) / 16)$

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LOWER $=$ PEEK (QQ) $-16 *$ HIGHER
BYTEVALUE $=10 *$ HIGHER + LOWER
NEWBCDVALUE $=$ NEWBCDVALUE + BYTEVALUE
OLDBCDVALUE = NEWBCDVALUE
Hopefully, your algebra is up to understanding how line 32535 is just a simplification of all that. If not, don't worry about it. It works.

But we still haven't accounted for the exponent. Now, exponents in the Atari floating point format are powers of 100 in "excess 64" notation, which simply means that you subtract 64 from the exponent to get the real power of 100 . But wait! The implied decimal point is all the way to the left of the number. So we must bias our "excess 64" by the five multiplies-by-100 we did in deriving the BCD value. All that is done in line 32540.

Finally, we simply count the exponent down to one or up to minus one, depending on what it started at. And line 32545 is tricky, but not too much so. I will leave its inner workings as an exercise for you, the reader.

And, hard though it may be to believe, we arrive at line 32555 with the number in hand. Then we PRINT it.

Did we really have to go through all that? Not really, but perhaps it gives you an idea of what BASIC's GETTOK routine (\$AB3E) does when it encounters a variable name.

Finally, to test all this out, you should type it in, LIST it to disk or cassette, use NEW, and then enter or load your favorite program. Finally, re-ENTER this program fragment from disk or cassette and type GOTO 32500. Just for fun, you might try finding the variable values for the following program:


Type this little guy in, ENTER the variable value printer, and RUN the whole thing. Answer the variable number prompt with numbers from 128 to 135 and see what you get. It's interesting!

# Examining the Statement Table 

If you will recall, Chapter 3 in Part 1 discussed the various user tables that existed in Atari BASIC's RAM memory space. Specifically, it discussed the Variable Name Table, Variable Value Table, Statement Table, String/Array Table, and Runtime Stack.

In the last two chapters, we investigated the Variable Name Table and the Variable Value Table, showing how Atari BASIC can examine itself. So what is more logical than to now use Atari BASIC to display the contents of the Statement Table?

While we could write a program that would examine the tokenized program and produce source text, there is little incentive to do so. The task would be both very difficult and very redundant: BASIC's LIST command performs the same task very nicely, thank you.

What we can do, though, is write a program which will show the actual hex tokens used in a logical and almost readable form. Again, let's look at the program before decoding what it does.


## Chapter Five

| $\begin{aligned} & 32 ø \varnothing \emptyset \\ & 32 ø 1 \varnothing \end{aligned}$ |  |
| :---: | :---: |
|  | $\begin{aligned} & \mathrm{Q}=\mathrm{PEEK}(\mathrm{QQ})+256 * \operatorname{PEEK}(\mathrm{QQ}+1): Q S=Q Q: Q Q \\ & =Q Q+3 \end{aligned}$ |
| 32015 | IF Q>32767 THEN PRINT "--END--":ST |
|  | OP |
| 32020 | QL=PEEK (QQ-1) +QS:PRINT "LINE NUMBE |
|  | R ";Q,"LINE LENGTH ";PEEK(QQ-1) |
| 32030 | QT=PEEK (QQ+l):PRINT "\{2 SPACES ${ }^{\text {STM }}$ |
|  | T LENGTH "; PEEK (QQ), "STMT CODE "; |
|  | EEK (QQ+l) |
| 32040 | $Q=P E E K(Q Q)+Q S: Q Q=Q Q$ |
| 32050 | IF QQ<Q THEN 32ø8ø |
| 32060 | IF Q<QL THEN PRINT :GOTO 32Ø3ø |
| 32070 | PRINT :GOTO 3201ø |
| 32080 | IF QT>1 AND QT<55 THEN 3212ø |
| 32090 | PRINT "\{2 SPACES\}UNTOKENIZED: : "; |
| $321 \varnothing \varnothing$ | PRINT CHR\$ (PEEK (QQ)) ; : $\mathrm{QQ=QQ}$ |
|  | Q<Q THEN 321øø |
| 32110 | PRINT :GOTO 3201ø |
| 32120 | NUM=PEEK (QQ) : GOSUB CVDEC |
| 32125 | IF PEEK (QQ) > 127 THEN PRINT " V=";N UM\$ (3):GOTO 322øø |
| 32130 | IF PEEK (QQ) >15 THEN PRINT " ";NUM\$ (3): : GOTO $322 \varnothing \varnothing$ |
| 32140 | IF PEEK (QQ) $=14$ THEN GOTO $3217 \emptyset$ |
| 32150 | $Q Q=Q Q+1: Q N=\operatorname{PEEK}(Q Q): N U M=Q N: G O S U$ |
|  | VDEC |
| 32155 | PRINT " S,";NUM\$ (3);"=";:IF QN=ø |
|  | HEN 322øø |
| 32160 | FOR $Q Q=Q Q+1$ TO QQ+QN-1:PRINT CHR\$( |
|  | $\operatorname{PEEK}(Q Q)) ;:$ NEXT QQ:GOTO 3219め |
| 32170 | PRINT " $\mathrm{N}=$ |
| 32180 | $F O R \quad Q Q=Q Q+1$ TO $Q Q+5: \mathrm{NUM}=\mathrm{PEEK}(\mathrm{QQ}): \mathrm{G}$ |
|  | OSUB CVDEC:PRINT NUM\$ (3);:NEXT QQ |
| 32190 | $Q Q=Q Q-1: P R I N T$ |
| 322 øø | $Q Q=Q Q+1: I F ~ Q Q<Q ~ T H E N ~ 3212 \emptyset ~$ |
| 32210 | PRINT : IF QQ<QL THEN 32ø3ø |
| 32220 | PRINT :GOTO 32ø1ø |

Now, even if you don't want to type all that in, there are a few points to be made about it. First, note that lines 10 through 50 and 9100 through 9160 are the decimal-to-hex converter from

## Chapter Five

Chapter 2. Then, let's start with line 32000 and do a functional description, with the line numbers denoting the portion we are examining.
32000. Decimal 136 is hex $\$ 88$, the location of STMTAB, the pointer to the user's program space.

32010, 32020. In each line, the first two bytes are the line number; the next byte is the line length (actually, the offset to next line). Remember, line 32768 is actually the direct statement.

32030, 32040. Within a line, each statement begins with a statement length (the offset to the next statement from the beginning of the line) and a statement token.

32050-32070. Boundary conditions are checked for.
32080-32110. REM becomes statement token 0, DATA is token 1 and the error token is 55 (\$37). All three of them simply store the user's input unchanged.
32120. Remember, any token with its upper bit on indicates a variable number token. They really don't need to be special cased in this program, but we do so for readability.
32130. Operator tokens have values of 16 to 127 ( $\$ 10$ to \$7F).

32140-32160. For string constants (also called string literals), we simply print out the string length and its contents (the characters between the quote signs).

32170-32180. For numeric constants, we simply print the hex values of all six bytes.

32190-32200. Clean-up. We ensure that we return for all remaining tokens (if any) in each statement and for all remaining statements (if any) in each line.

Observe the FOR-NEXT loop controls in line 32180. Why $Q Q+1$ TO QQ + 5 if we want six values printed out? Ah, but this is a trick. Note that the loop termination value $(Q Q+5)$ involves the loop variable (QQ). The problem is, though, that the loop variable is changed by the prior implied assignment $(\mathrm{QQ}=\mathrm{QQ}+1)$ when the assignment takes place - which is, of course, before the determination of the value of " $\mathrm{QQ}+5$ " takes place.

In other words, by the time we are ready to evaluate $Q Q+5$, the variable $Q Q$ has already been changed from its original value to its new, loop controlling value ( $\mathrm{QQ}+1$ ).

Quite possibly, the proper general solution to using a FOR loop's variable in its own termination (or STEP) values is to

## Chapter Five

assign it to a temporary variable, thusly:

$$
\text { QTEMP = QQ:FOR QQ = QTEMP + } 1 \text { TO QTEMP + } 6
$$

Did you notice that line 32160 actually has the same problem? Notice that we solved it there by adding -1 to the termination value to compensate for the +1 in the initialization assignment.

One last comment before leaving the subject of strange FOR-NEXT loops. In Atari BASIC (and, indeed, in virtually all microcomputer BASICs), the termination (TO) value and the STEP value are determined when the FOR statement is first executed and are NOT changeable. Example:
$10 \mathrm{X}=7: \mathrm{Y}=2$
$2 \emptyset$ FOR $\mathrm{I}=1$ TO X STEP Y
$30 \mathrm{X}=\mathrm{X}+1$
$40 \mathrm{Y}=\mathrm{Y}+\mathrm{X}$
50 NEXT I
This FOR loop will execute exactly four times ( $\mathrm{I}=1,3,5$, and 7). The fact that $X$ and $Y$ change within the loop has no effect on the actual loop execution.

## Viewing the

The Runtime Stack is the last of the user RAM tables that we will discuss in Part 2.

Perhaps you noticed that we left out a discussion of the String/Array Table in Part 2. The omission was on purpose: there seems little purpose in PEEKing the contents of this table when BASIC's PRINT statement does an admirable job of letting you see all variable values. However, if you are so inclined, you could use the general purpose memory PEEKer program of Chapter 2 to view any portion of any memory, including the String/Array Table.

On the other hand, looking at the Runtime Stack is kind of fun and enlightening. And the program we will present here might even find use on occasion. If you are having trouble tracing a program's flow, through various GOSUBs and/or FOR loops, simply drop in the routine below and GOSUB to it at an appropriate place in your program. It will print out a LIFO (Last In, First Out) listing of all active GOSUB calls and FORNEXT loop beginnings.

```
1\emptyset FOR J=1 TO 3
2\emptyset GOSUB 3\emptyset
3\emptyset FOR K=1 TO 5
4\emptyset GOSUB 5\emptyset
5\emptyset JUNK=7:FOR Q=1 TO 2:GOSUB 324\emptyset\emptyset
3240\varnothing QQ=PEEK(144)+256*PEEK(145)
3241\emptyset IF QQ<=PEEK(142)+256*PEEK(143) THE
    N PRINT "--END OF STACK--":STOP
3242\emptyset PRINT "AT LINE ";PEEK(QQ-3)+256*PE
    EK(QQ-2);
3243\emptyset PRINT ", OFFSET ";PEEK(QQ-1);
3244\emptyset IF PEEK(QQ-4)=\emptyset THEN PRINT ", GOSU
    B":QQ=QQ-4:GOTO 3241Ø
3245\emptyset PRINT ", FOR (#";PEEK(QQ-4);")":QQ
        =QQ-16:GOTO 3241Ø
```


## Chapter Six

The first thing you might notice about this little routine is that, in contrast to all the programs we have used so far, it examines its portion of user RAM backward. That is, it starts at the top (high address) of the Runtime Stack area and works downward toward the bottom.

Again, nothing surprising. If you will recall the description of entries on this stack (pages 18-19 and 133-34), you will remember that every entry, whether a GOSUB or FOR, has a four-byte header. And, while FOR statements also have twelve bytes of termination and step value added, the four bytes are always at the top of each entry - they are the last items put on the stack.

Thus, we start at the top of the stack and examine four bytes. If the type byte is zero, it is a GOSUB entry, and all we must do is display the line number and statement offset. If we remove the four-byte header by subtracting 4 from our stack pointer, we are ready to examine the next entry.

In the case of a FOR entry, we similarly display the line number and statement offset. However, each FOR entry also has a variable token associated with it, so we also display that token's value. With the variable name lister of Chapter 2, you can find out which variable is controlling this FOR loop. Finally, note that after displaying a FOR loop entry, we remove sixteen bytes (the four-byte header and the two six-byte floating point values) in preparation for the next entry.

Incidentally, lines 10 through 50 are present as examples only. Add lines 32400 to 32450 to your own programs and see where you've come from.

## Fixed Tokens

In the last chapter, we discussed the last of the tables in user RAM. Now we will see how and where BASIC stores its internal ROM-based tables.

As we noted in Chapter 5 of Part 1(and viewed via the listing program of Chapter 5 in this Part), there are four kinds of tokens in an Atari BASIC program: (1) statement name tokens, (2) operator tokens, (3) variable tokens, and (4) constant tokens (string and numeric constants). Also, we learned in Part 1 how the tokenizing process works, converting the user's ATASCII source code into tokens. What we didn't learn, though, was exactly what token replaces what BASIC keyword.

In this chapter, we present a program which will list all of the fixed tokens (those in ROM). Actually, the program presents three listings, each consisting of a list of token values with their associated ATASCII strings. But wait a moment! Three listings? There are only two ROM-based tables - SNTAB and OPNTAB.

Yes, but it seems that this program is also capable of listing the Variable Name Table. Why list it again, when we did it so well in Chapter 3? Because we wanted to show you how BASIC itself does it. In many ways, this program emulates the functions of the SEARCH routine at address \$A462 in the source listing. And, yes, BASIC uses a single routine to search all three of these same tables. You might want to examine BASIC's SEARCH routine at the same time you peruse this listing.

```
1\emptyset\emptyset REM we make use of the general purpose
ll\varnothing REM token lister three times:
2ø\varnothing PRINT :PRINT "A LIST OF VARIABLE TOKENS"
21\emptyset ADDR=PEEK(13Ø)+256*PEEK(131)
22ø SKIP=\emptyset:TOKEN=128:GOSUB 1øø\emptyset
3ø\emptyset PRINT :PRINT "A LIST OF STATEMENT TOKENS"
31\varnothing ADDR=42159:SKIP=2:TOKEN=\emptyset:GOSUB 1øø\emptyset
4\emptyset\emptyset PRINT :PRINT "A LIST OF OPERATOR TOKENS"
41\varnothing ADDR=42979:SKIP=\varnothing:TOKEN=16:GOSUB 1\varnothingø\varnothing
420 STOP
1\varnothing\varnothing\varnothing REM a general purpose token listing routine
```


## Chapter Seven

```
1Ø01 REM
10Ø2 REM On entry to this routine, the following
lØ\emptyset3 REM variables have meanings:
lØ\emptyset4 REM ADDR = address of beginning of table
10\emptyset5 REM SKIP = bytes per entry to skip
1Ø\emptyset6 REM TOKEN = starting token number
10\emptyset7 REM
11Ø\emptyset ADDR=ADDR+SKIP:IF PEEK(ADDR)=\emptyset THEN RETURN
111\emptyset PRINT TOKEN,:TOKEN=TOKEN+1
112\emptyset IF PEEK(ADDR)>127 THEN 114\emptyset
1130 PRINT CHR$(PEEK(ADDR));:ADDR=ADDR+1:GOTO 112\emptyset
114\emptyset PRINT CHR$(PEEK(ADDR)-128):ADDR=ADDR+1:GOTO 11Ø\emptyset
```

The main routine is actually lines 1100 through 1140 (while lines 1000 through 1007 simply explain it all). It's actually fairly simple. Each table is assumed to consist of a fixed number of bytes followed by a variable number of ATASCII bytes, the last of which has its upper bit on.

In line 1100, we skip over the fixed bytes (if any) and check for the end of the table. After that, we simply print the token value followed by the name.

Worth examining, though, are lines 200 through 420, where we call the main subroutine. First, note that the Variable Name Table has no bytes to skip and is located via its zero-page pointer. Naturally, the first variable token value is 128.

Each entry in the Statement Name Table (SNTAB, at location \$A4AF) has two leading bytes (actually, the two-byte address, minus 1 , of the syntax table entry for this statement). Statement name token values begin at zero, and 42159 is the decimal address of SNTAB.

Finally, the smallest-numbered operator token is 16 decimal (except for string and numeric constants, which are special cased). There are no leading bytes in the Operator Name Table, and it starts at location 42979 decimal (OPNTAB, at \$A7E3).

## What Takes Precedence?

There was one other ROM-based table mentioned in Part 1 which deserves some attention here. You may recall that when an expression is executed, the execution operators are given particular precedences, so that in BASIC, $2+3 * 4$ equals 14, not 20. Chapter 7 of Part 1 does a particularly thorough job of explaining the concepts of precedence.

The program presented in this chapter prints out all of BASIC's operator tokens along with their token walues and their dual precedence values. Actually, the program provides a visual readout of OPRTAB (Operator PRecedence TABle, at \$AC3F).

In each pair of precedence values listed, the first number is the go-onto-stack value and the second is the come-off-stack value.

```
I\emptyset\emptyset PRINT "A LIST OF OPERATOR TOKENS"
ll\emptyset PRINT " WITH THEIR PRECEDENCE TABLE VALUES"
22Ø SKIP=\emptyset:TOKEN=128:GOSUB 1\varnothingØ\emptyset
lø\emptyset\emptyset ADDR=42979:REM WHERE OP NAMES START
1Øl\emptyset TOKEN=16:REM LOWEST TOKEN VALUE
1\emptyset2\emptyset REM NOW THE MAIN CODE LOOP
ll\emptyset\emptyset IF PEEK(ADDR)=\emptyset THEN STOP
111\emptyset PRINT TOKEN,:PREC=PEEK(44095+TOKEN-16)
112\emptyset PRINT INT(PREC/16);":";PREC-16*INT(PREC/16),
1130 PREC=PEEK (ADDR) : ADDR=ADDR+1
114\varnothing IF PREC<l }28 THEN PRINT CHR$(PREC);:GOTO 113\emptyset
115\emptyset PRINT CHR$(PREC-128):TOKEN=TOKEN+1:GOTO l1ø\emptyset
```

If you closely examined the program in the last chapter, you will note a striking similarity to this program, especially lines 1100 through 1150. Actually, the only thing we have really added is the precedence printout of line 1120.

And note the form of the PEEK in line 1110. Then look at the line of code at address \$AAF1 in the BASIC listing. Given

## Chapter Eight

the limitations of dissimilar languages, the code is identical. This is more evidence that you really can use BASIC as a tool to diagnose itself.

# Using What We Know 

Now that Atari BASIC stands revealed before you, what do you do with it? Many authors have, even without benefit of the listing in this book, either used or fooled BASIC in ways that we who designed it never dreamed of.

For example, consider what happens if you change BASIC's STARP pointer (\$8C) to be equal to its ENDSTAR value (\$8E). Remember, BASIC's SAVE command saves everything from the contents of VNTP to the contents of STARP (as documented in Chapter 10 of Part 1). So changing what is in STARP is tantamount to telling BASIC to SAVE more (or less) than what it normally would. Presto! We can now save the entire array and string space to disk or tape, also.

Is it useful? Here's one program that is, using the concepts we learned in the previous chapters.

```
3øø\emptyset\emptyset PRINT :PRINT "WHAT VARIABLE NUMBER
                DO YOU":PRINT,"WISH TO FIND ";
3øø1\emptyset INPUT QV
3Ø\varnothing2\emptyset QA=PEEK(130)+256*PEEK(131):QN=128
3ø\emptyset3\emptyset IF QN=QV THEN 3Øø6\emptyset
3ø\varnothing4\emptyset IF PEEK(QA)<128 THEN QA=QA+l:GOTO
    3øø4\emptyset
3ø\emptyset5\emptyset QN=QN+1:QA=QA+1:GOTO 3Ø\emptyset 3\emptyset
30ø60 IF PEEK(QA)<128 THEN PRINT CHR$(PE
    EK(QA));:QA=QA+l:GOTO 3ø\varnothing6\emptyset
3ø\emptyset7\emptyset PRINT CHR$(PEEK(QA)-128);" IS THE
        VARIABLE"
3ø1\varnothing\emptyset QA=PEEK(136)+256*PEEK(137)
3\emptysetll\emptyset QN=PEEK(QA)+256*PEEK(QA+1):QL=PEEK
    (QA+2):QSV=QA:QA=QA+3
3ø12\varnothing IF QN>32767 THEN PRINT "--END--":E
    ND
3013\emptyset QS=PEEK(QA):QT=PEEK(QA+1):QA=QA+2:
        IF QT>l AND QT<55 THEN 3Ø15\varnothing
```

| 30140 | $Q A=Q S V+Q L: G O T O$ | 30110 |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 30150 | IF $\operatorname{PEEK}(Q A)=Q V$ | THEN | PRINT | "LINE | "; |
|  | QN:GOTO 3Ø14Ø |  |  |  |  |
| 30160 | IF $\operatorname{PEEK}(Q A)>15$ | THEN | $3020 \emptyset$ |  |  |
| 30170 | $\operatorname{IF} \operatorname{PEEK}(Q A)=14$ $\emptyset 2 \emptyset \emptyset$ | THEN | $\mathrm{QA}=\mathrm{QA}+$ | : GOTO | 3 |
| 30180 | $Q A=Q A+P E E K(Q A+1)$ | ) +1 |  |  |  |
| 30200 | $\mathrm{QA}=\mathrm{QA}+1: \mathrm{IF} \quad \mathrm{QA}<$ | SV+QS | THEN | 30150 |  |
| 30210 | IF QA<QSV+QL TH | EN 3Ø | $13 \varnothing$ |  |  |
| 30220 | GOTO 3Ø11Ø |  |  |  |  |

What does it do? It finds all the places in your program that you used a particular variable. And how do you use it? Type it in, LIST it to disk or cassette, and clear the user memory via NEW. Now type, ENTER, or LOAD the program you wish to investigate (and then SAVE it, if you haven't already done so). Finally, ENTER this program fragment from the disk or cassette where you LISTed it and type GOTO 30000.

Although the program asks you for a variable number (which you can get via the program of Chapter 3), it doesn't really matter if you don't know it. The program will print your chosen variable's name before giving all the references. If you chose wrong, try again.

And how does it work? Somewhat like the program token lister of Chapter 5, except that here we are simply skipping everything but variable name references. First, though, we use a modified Variable Name Table lister (lines 30020 through 30070) to tell you what name you chose.

Then, we start at the beginning of the program (line 30100) and check each user line number (30110 and 30120). Within each line, we loop through, checking all statements (30130), skipping entirely all REMs, DATA lines, and lines with syntax errors (line 30140). If we find ourselves in an expression, we check for a matching variable token reference (line 30150) and print it if found, after which we skip the rest of the line. We also skip over numeric and string constants (lines 30170 and 30180). Finally, we check to see if we are at the end of the statement (30200) or the end of a line (30210 and 30220).

This is a fairly large program fragment, and it will prove most useful in very large programs, where you can't remember, for example, how many places you are using the variable name LOOP. So you might want to try to leave room in memory for this aid; you may be very glad you did.

## Part Three

## Atari BASIC Source Code

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Cupertino, California 95014 (U.S.A.)
Telephone: (408) 446-3099

## Source Code

## Some Miscellaneous Equates

| $=0001$ | PATSIZ | EQU | \＄1 | ；PATCH AREA SIZE |
| :---: | :---: | :---: | :---: | :---: |
| $=\varnothing 020$ | ZICB | EQU | \＄20 | ；zero PageIOCB |
| ＝Øø8ø | ZPG1 | EQU | \＄80 | ；beginning of BASIC＇s zero page |
| $=\varnothing 48 \varnothing$ | MISCR1 | EQU | \＄480 | ；syntax stack，etc． |
| $=\varnothing 5 \emptyset \emptyset$ | MISCRAM | EQU | \＄50ø | ；other RAM usage |
| $=\mathrm{E} 456$ | CIO | EQU | \＄E456 | ；in OS ROMs |
| $=\varnothing 340$ | IOCBORG | EQU | \＄340 | ；where IOCBs start |
| $=\varnothing 300$ | DCBORG | EQU | \＄30ø | ；where DCB（for SIO）is |
| ＝ A ¢øø | ROM | EQU | \＄Aøøø | ；begin code here |
| $=\varnothing \emptyset \mathrm{D} 2$ | ZFP | EQU | \＄D2 | ；begin fltg point work area |
| $=\varnothing 09 \mathrm{~B}$ | CR | EQU | \＄9B | ；ATASCII end of line |
| $=02 \mathrm{E} 7$ | LMADR | EQU | \＄2E7 | ；system lo mem |
| $=\varnothing 2 \mathrm{E} 5$ | HMADR | EQU | \＄2E5 | ；system high mem |
| $=\square 2 \mathrm{E} 5$ | HIMEM | EQU | HMADR |  |
| $=\mathrm{D} 8 \emptyset \emptyset$ | FPORG | EQU | \＄D80ø | ；fltg point in OS ROMs |
| ＝ 0011 | BRKBYT | EQU | \＄11 |  |
| ＝ 0008 | WARMFL | EQU | \＄ø8 | ；warmstart flag |
| $=\mathrm{D} 20 \mathrm{~A}$ | RNDLOC | EQU | \＄D20A | ；get a random byte here |
| ＝BFF9 | CRTGI | EQU | \＄BFFC－3 | ；cartridge init vector |
| $=\varnothing 05 \mathrm{D}$ | EPCHAR | EQU | \＄5D | ；the＂？＂for INPUT statement |
| $=\mathrm{E} 471$ | BYELOC | EQU | \＄E471 | ；where to go for BYE |
| $=\varnothing \varnothing \emptyset \mathrm{A}$ | DOSLOC | EQU | \＄日A | ；via here to exit to DOS |
| ＝Øø55 | SCRX | EQU | \＄55 | ；X AXIS |
| ＝øø54 | SCRY | EQU | \＄54 | ；Y AXIS |
| $=\emptyset 2 \mathrm{C} 4$ | CREGS | EQU | \＄2C4 | ；COLOR REGISTER |
| $=\varnothing 2 \mathrm{FB}$ | SVCOLOR | EQU | \＄2FB | ；SAVE COLOR FOR CIO |
| ＝ D 208 | SREGI | EQU | \＄D208 | ；SOUND REG 1 |
| $=\mathrm{D} 200$ | SREG 2 | EQU | \＄D2øの | ；SOUND REG 2 |
| $=\mathrm{D} 201$ | SREG3 | EQU | \＄D201 | ；SOUND REG 3 |
| $=\mathrm{D} 2 \emptyset \mathrm{~F}$ | SKCTL | EQU | \＄D20F | ；sound control |
| $=\varnothing 270$ | GRFBAS | EQU | \＄270 | ；IST GRAPHICS FUNCTION ADDR |
| $=\varnothing 2 \mathrm{FE}$ | DSPFLG | EQU | \＄2FE | ；ATARI DISPLAY FLAG |
| ＝ØøøE | APHM | EQU | \＄E | ；APPLICATION HIGH MEM |

## Zero Page

## RAM Table Pointers

| øøøø | $=\varnothing 口 8 \varnothing$ | ORG |  | ZPG1 |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Øø8ø |  | LOMEM |  |  | ； | LOW MEMORY POINTER |
| Øø8ø |  | ARGOPS |  |  | ； | ARGUMENT／OPERATOR STACK |
| Øø8ø |  | ARGSTK |  |  |  |  |
| øø8ø | $=\varnothing 002$ | OUTBUFF | DS | 2 | ； | SYNTAX OUTPUT BUFFER |
| 0682 | ＝Øøø2 | VNTP | DS | 2 | ； | VARIABLE NAME TABLE POINTER |
| 0084 | ＝øøø2 | VNTD | DS | 2 | ； | VARIABLE NAME TABLE DUMMY END |
| $\emptyset \emptyset 86$ | ＝Ø0ø2 | VVTP | DS | 2 | ； | VARIABLE VALUE TABLE POINTER |
| 9088 |  | ENDVVT |  |  | ； | END VARIABLE VALUE TABLE |
| øø88 | ＝øøø2 | STMTAB | DS | 2 | ； | STATEMENT TABLE［PROGRAM］； POINTER |
| 0ø8A | ＝Øøø2 | STMCUR | DS | 2 | ； | CURRENT PGM PTR |
| のø8С | ＝Øøø2 | STARP | DS | 2 | ； | STRING／ARRAY TABLE POINTER |
| の08E |  | ENDSTAR |  |  | ； | END STRING／ARRAY SPACE |
| 908E | $=\varnothing \square \square 2$ | RUNSTK | DS | 2 | ； | RUN TIME STACK |
| 0090 |  | TOPRSTK |  |  | ； | END RUN TIME STACK |
| のロ9の | $=\varnothing \emptyset \square 2$ | MEMTOP | DS | 2 | ； | TOP OF USED MEMORY |
| 0092 | ＝Øøロ1 | MEOLFLG | DS | 1 | ； | MODIFIED EOL FLAG |
| 0093 | ＝Øøø1 |  | DS | 1 | ； | ：：SPARE： |

## Source Code

## Miscellaneous Zero Page RAM


SEARCH ADR
ARRAY INDEX 2
EXPAND START ADR
MOVE FROM ADR
MOVE TO ADR
CUR SYNTAX PGM COUNTER
MAX SYNTAX CIX
LINE LENGTH
TEST LINE NO
MOVE LENGTH
MOVE SIZE
DIRECT EXECUTE FLAG
STMT LENGTH BYTE DISPL
SXT STMT DISR
star START CIX
CURR STMT INDEX
SYNTAX STACK LEVEL
INPUT BUFFER INDEX
SEARCH SKIP FACTOR
ARG STACK INDEX
TEOW LENGTH BYTE PTR
SAVED OPERATOR
SAVED OPERATOR PRECEDENCE
SAVE VAR NAME TBL PTR
LIST END LINE \#
TEMP FOR ARRAYS
EARCH TABLE ENTRY NUMBER
LIST SCAN COUNTER
COMMA COUNT FOR EXEXOR
SAVE VAR VALUE EXP SIZE
ASSIGN/DIM FLAG
OVE SRC AR
DISPL INTO LINE OF FOR/GOSUB
LOOP CONTROL FOR OP
SAVE ONT SRC INDEX
NDEX INTO STMT
DATA DISPL
DATA LINO
LINE \# STOPPED AR [FOR CON]
TRAP LINE \# [FOR ERROR]
SAVE CURRENT LINE ADDR
I / O COMMAND
I/O DEVICE
ERROR \# FOR USER
TEMP ADDR CELL
TEMP
PRINT TAB WIDTH
LOAD IN PROGRESS FLAG

## Source Code

## Argument Work Area（AWA）

Floating Point Work Area

| ØøСВ | $=\emptyset \square \mathrm{D} 2$ | ORG |  | ZFP |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ØøD2 |  | TVTYPE |  |  | ； | VARIABLE TYPE |  |
| ØøD2 | ＝ØØØ1 | VTYPE | DS | 1 | ； | VARIABLE TYPE |  |
| ØøD3 |  | TVNUM |  |  | ； | VARIABLE NUMBER |  |
| ØøD3 | ＝Ø0ø1 | VNUM | DS | 1 | ； | VARIABLE NUMBER |  |
|  | $=\varnothing \emptyset 06$ | FPREC | EQU | 6 |  |  |  |
|  | $=\varnothing 005$ | FMPREC | EQU | FPREC－1 | ； | LENGTH OF FLOATING | POINT |
|  |  |  |  |  | ； | MANTISSA |  |
| ØØD4 |  | BININT |  |  | ； | FP REGO |  |
| ØøD4 | $=\varnothing \emptyset 01$ | FRø | DS | 1 | ； | FP REGØ |  |
| ØØD5 | $=\varnothing 005$ | FRØM | DS | FPREC－1 | ； | FP REGØ MANTISSA |  |
| ØØDA | $=\boxed{006}$ | FRE | DS | FPREC | ； | FP REGØ EXP |  |
| ØロEの | $=\varnothing \emptyset \emptyset 1$ | FRI | DS | 1 | ； | FP REG 1 |  |
| の日El | $=0005$ | FRIM | DS | FPREC－1 | ； | FP REGI MANTISSA |  |
| のøE6 | ＝Øøø6 | FR2 | DS | EPREC | ； | FP REG 2 |  |
| ØøEC | ＝ØøØ1 | FRX | DS | 1 | ； | FP SPARE |  |

## RAM for ASCII to Floating Point Conversion

| ØøED | $=\varnothing \emptyset \emptyset 1$ | EEXP | DS | 1 | ； | VALUE OF | E |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ØøEE |  | FRSIGN |  |  | ； | FP SIGN |  |  |
| ØøEE | $=\varnothing 001$ | NSIGN | DS | 1 | ； | SIGN OF | \＃ |  |
| ØøEF |  | SQRCNT |  |  |  |  |  |  |
| ØøEF |  | PLYCNT |  |  |  |  |  |  |
| ØøEF | $=\varnothing 口 \square 1$ | ESIGN | DS | 1 | ； | SIGN OF | EXPONENT |  |
| ØØFの |  | SGNFLG |  |  |  |  |  |  |
| の日Fの | $=\varnothing \varnothing \varnothing 1$ | FCHRFLG | DS | 1 | ； | 1ST CHAR | FLAG |  |
| の日F1 |  | XFMFLG |  |  |  |  |  |  |
| ØØF1 | $=\varnothing 001$ | DIGRT | DS | 1 | ； | \＃OF DIG | ITS RIGHT | DECIMAL |

## Input Buffer Controls



Temps

| ØøF5 | $=$ Øøø2 | ZTEMP1 | DS | 2 | ；LOW LEVEL ZERO PageTEMPS |
| :--- | :--- | :--- | :--- | :--- | :--- |
| ØøF7 | $=$ Øøø2 | ZTEMP4 | DS | 2 |  |

Miscellany

| ØøFB |  | DEGFLG |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ØØFB | $=\varnothing 001$ | RADFLG | DS | 1 | ； | $\emptyset=$ RADIANS，6＝DEGREES |
|  | ＝øøøø | RADON | EQU | $\emptyset$ | ； | INDICATE RADIANS |
|  | ＝øøø6 | DEGON | EQU | 6 | ； | INDICATES DEGREES |
| ØØFC | ＝ 0002 | FLPTR | DS | 2 | ； | POLYNOMIAL POINTERS |
| ØØFE | ＝øøø2 | FPTR2 | DS | 2 |  |  |

## Miscellaneous Non－Zero Page RAM



## Source Code

| ø600 | $=\varnothing 5 \mathrm{E} \varnothing$ | ORG |  | LBUFF＋\＄6 ${ }^{\text {d }}$ |
| :---: | :---: | :---: | :---: | :---: |
| 65Eø | ＝ 9006 | Plyarg | DS | FPREC |
| 65E6 | ＝øøø6 | FPSCR | DS | FPREC |
| ø5EC | ＝øøø6 | FPSCR1 | DS | FPREC |
|  | ＝ 05 E 6 | FSCR | EQU | FPSCR |
|  | $=05 \mathrm{EC}$ | FSCRI | EQU | FPSCR1 |

## IOCB Area

65F2 $=\varnothing 34 \varnothing$

## IOCB－I／O Control Block



## ICCOM Value Equates

| $=\varnothing 001$ | ICOIN | EQU | \＄ø1 | OPEN INPUT |
| :---: | :---: | :---: | :---: | :---: |
| ＝øøø2 | ICOOUT | EQU | \＄ø2 | ；OPEN OUTPUT |
| ＝øøø 3 | ICOIO | EQU | \＄ø3 | ；OPEN UN／OUT |
| $=\varnothing \varnothing \varnothing 4$ | ICGBR | EQU | \＄04 | GET BINARY RECORD |
| ＝øøø5 | ICGTR | EQU | \＄05 | ；GET TEXT RECORDS |
| ＝Øøø6 | ICGBC | EQU | \＄ø6 | ；GET BINARY CHAR |
| ＝ 0007 | ICGTC | EQU | \＄07 | ；GET TEXT CHAR |
| ＝øøø8 | ICPBR | EQU | \＄08 | ；PUT BINARY RECORD |
| ＝Øøø9 | ICPTR | EQU | \＄09 | ；PUT TEXT RECORD |
| ＝ØøøA | ICPBC | EQU | \＄øA | ；PUT BINARY CHAR |
| $=$ øøøВ | ICPTC | EQU | \＄øB | ；PUT TEXT CHAR |
| ＝øøøट | ICCLOSE | EQU | \＄øC | ；CLOSE FILE |
| $=$ の日のD | ICSTAT | EQU | \＄øD | ；GET STATUS |
| ＝øøøE | ICDDC | EQU | \＄øE | ；DEVICE DEPENDENT |
| ＝ØøøE | ICMAX | EQU | \＄øE | ；MAX VALUE |
| $=\varnothing \emptyset F F$ | ICFREE | EQU | \＄FF | ；IOCB FREE INDICATOR |
| ＝Øø1C | ICGR | EQU | \＄1C | ；OPEN GRAPHICS |
| ＝Ø011 | ICDRAW | EQU | \＄11 | DRAW TO |

## ICSTA Value Equates

| $=\varnothing \emptyset \emptyset 1$ | ICSOK | EQU | \＄ø1 | STATUS GOOD，NO ERRORS |
| :---: | :---: | :---: | :---: | :---: |
| $=\boxed{0062}$ | ICSTR | EQU | \＄ø2 | ；TRUNCATED RECORD |
| ＝ 0003 | ICSEOF | EQU | \＄ø3 | ；END OF FILE |
| ＝Øø8ø | ICSBRK | EQU | \＄80 | ；BREAK KEY ABORT |
| $=\varnothing \emptyset 81$ | ICSDNR | EQU | \＄81 | ；DEVICE NOT READY |
| ＝Øø82 | ICSNED | EQU | \＄82 | ；NON－EXISTENT DEVICE |
| $=\varnothing 083$ | ICSDER | EQU | \＄83 | DATA ERROR |
| $=\varnothing \emptyset 84$ | ICSIVC | EQU | \＄84 | ；INVALID COMMAND |
| ＝ 0085 | ICSNOP | EQU | \＄85 | ；DEVICE／FILE NOT OPEN |
| ＝ 0086 | ICSIVN | EQU | \＄86 | ；INVALID IOCB NUMBER |
| $=\varnothing 087$ | ICSWPE | EQU | \＄87 | ；WRITE PROTECTION |

## Equates for Variables

|  | $;$ |  | －IN VARIABLE VALUE TABLE |
| :--- | :--- | :--- | :--- |
|  |  | －ON ARGUMENT STACK |  |

## Equates for Run Stack

| $=0004$ | GFHEAD | EQU | 4 | ； | LENGTH OF | HEADER FOR FOR／GOSUB |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ＝øøøट | FBODY | EQU | 12 | ； | LENGTH OF | BODY OF FOR ELEMENT |
| ＝ 0003 | GFDISP | EQU | 3 | ； | DISP TO SA | AVED LINE DISP |
| ＝Øøø1 | GFLNO | EQU | 1 | ； | DISPL TO L | LINE \＃IN HEADER |
| ＝Øøøø | GFTYPE | EQU | $\emptyset$ | ； | DISPL TO T | TYPE IN HEADER |
| ＝Øøø6 | FSTEP | EQU | 6 | ； | DISPL TO S | STEP IN FOR ELEMENT |
| ＝Øøøø | FLIM | EQU | $\emptyset$ |  | DISPL TO L | LIMIT IN FOR ELEMENT |

## ROM Start

## Cold Start

| Аøøø |  |
| :---: | :---: |
| Аøøø | A5CA |
| Аøø2 | Døø4＾Aøの8 |
| Aøø4 | A508 |
| Аøø6 | Dø45＾Aø4D |
| Аøø8 |  |
| Аøø8 | A2FF |
| AØロA | 9A |
| AøøB | D8 |
| AøロC |  |
| AøøC | AEE762 |
| AØØF | ACE8ø2 |
| AØ12 | 8680 |
| A014 | 8481 |
| AØ16 | A9øø |
| A018 | 8592 |
| AØ1A | 85CA |
| A01C | C8 |
| AØ1D | 8A |
| AØ1E | A282 |
| Aø2ø | 950ø |
| Aø22 | E8 |
| Aø23 | 9400 |
| Aø25 | E8 |
| AØ26 | E092 |
| AØ28 | 90F6＊Aø20 |
| Aø2A | A286 |



Source Code

| Aø2C | Aøø1 |  | LDY | \#1 | ; | FOR END OF VNT |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Aø2E | 207FA8 |  | JSR | EXPLOW | ; | zero byte |
| Aø31 | A28C |  | LDX | \#STARP | ; | EXPAND STMT TBL |
| A.633 | Аøø3 |  | LDY | \#3 | ; | BY 3 BYTES |
| A635 | 2ø7FA8 |  | JSR | EXPLOW | ; | GO DO IT |
| Aø38 | А9øø |  | LDA | \# $\emptyset$ | ; | SET 0 |
| A.03A | A8 |  | TAY |  |  |  |
| Aø3B | 9184 |  | STA | [VNTD], Y | ; | INTO VVTP |
| Aø3D | 918A |  | STA | [STMCUR], Y | ; | INTO STMCUR+ø |
| AØ3F | C8 |  | INY |  |  |  |
| Аø4ø | A980 |  | LDA | \#\$8ø | ; | \$8ø Into |
| Aø42 | 918A |  | STA | [STMCUR], Y | ; | STMCUR+1 |
| Aø44 | C8 |  | INY |  |  |  |
| Aø45 | A903 |  | LDA | \#\$ø3 | ; | \$ø3 Into |
| A047 | 918A |  | STA | [STMCUR], y | ; | STMCUR+2 |
| Aø49 | A90A |  | LDA | \#10 | ; | SET PRINT TAB |
| Aø4B | 85C9 |  | STA | PTABW | ; | WIDTH TO 10 |

## Warm Start



## Syntax

AØ6ø
LOCAL
Editor - Get Lines of Input

| A060 | SYNTAX |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| АØ60 | A 5CA | LDA | LOADFLG | ; IF LOAD IN PROGRESS |
| Aø62 | DØ9C * Aøøø | BNE | COLDSTART | ; GO DO COLDSTART |
| Aø64 | A 2 FF | LDX | \# \$FF | ; RESTORE STACK |
| Aø66 | 9A | TXS |  |  |
| Aø67 | 2ø51DA | JSR | INTLBF | ; GO INT LBUFF |
| Aø6A | A95D | LDA | \#EPCHAR |  |
| AØ6C | 85C2 | STA | PROMPT |  |
| Aø6E | 2092BA | JSR | GLGO | ; |
| Aø71 | 2øF4A9 | JSR | TSTBRK | ; TEST BREAK |
| A074 | DØEA ^AØ6Ø | BNE | SYNTAX | ; BR IF BREAK |
|  |  | ; LDA | \#0 |  |
| Aø78 | A960 85 F 2 | STA | \# ${ }_{\text {CIX }}$ | - INPUT INDEX TO ZERO |
| Aø7A | 859F | STA | MAXCIX |  |
| Aø7C | 8594 | STA | COX | ; OUTPUT INDEX TO ZERO |
| Aø7E | 85A6 | STA | DIRFLG | ; SET DIRECT SMT |
| Аø8ø | 85B3 | STA | SVONTX | ; SET SAVE ONT CIX |
| AØ82 | 85Bø | STA | SVONTC |  |
| Aø84 | 85B1 | STA | SVVVTE | ; VALUE IN CASE |
| Aø86 | A584 | LDA | VNTD | ; OF SYNTAX ERROR |
| Aø88 | 85AD | STA | SVVNTP |  |
| AØ8A | A585 | LDA | VNTD+1 |  |
| Ag8C | 85AE | STA | SVVNTP+1 |  |
|  |  | : |  |  |
| Aø8E | 2øAlDB | JSR | SKBLANK | ; SKIP BLANKS |
| A091 | 209FAl | JSR | : GETLNUM | ; CONVERT AND PUT IN BUFFER |
| Aø94 | 2øC8A2 | JSR | : SETCODE | ; SET DUMMY FOR LINE LENGTH |
| A097 | A5D5 | LDA | BININT+1 |  |
| Aø99 | 1 ¢ø2 ^Aの9D | BPL | : SYNØ |  |
| Aø9B | 85A6 | STA | DIRFLG |  |

## Source Code



Source Code


| A192 | 2ØDØA9 |
| :--- | :--- |
| A195 | 68 |
| A196 | A8 |
| A197 | A28A |
| A199 | 2ØFBA8 |
| Al9C | 4C6ØAØ |


| JSR | GNXTL |  |  |
| :--- | :--- | :--- | :--- |
| PLA |  |  |  |
| TAY |  |  |  |
| LDX | \#STMCUR | ;GET STMCUR DISPL |  |
| JSR | CONTLOW | GO DELETE |  |
| JMP | SYNTAX | GGO FOR NEXT | LINE |

## Get a Line Number



## NEXT



## Source Code

| AlF7 | 9059 ^ 4252 | BCC |  | : POP | ; $\mathrm{CODE}=[2,3, \mathrm{OR} 4] \mathrm{POP}$ UP TO <br> ; NEXT SYNTAX CODE |
| :---: | :---: | :---: | :---: | :---: | :---: |
| AlF9 | 20A9A2 | JSR |  | : TERMTST | ; CODE>5 GO TEST TERMINAL |
| AlFC | 90E4 *AlE2 | BCC |  | : NEXT | ; BR IF SUCCESS |
| AlFE | $4 \mathrm{C6CA} 2$ | JMP |  | : FAIL | ; ELSE GO TO FAIL CODE |
|  |  | ; |  |  |  |
| A201 | 38 | : ERNTV | SEC |  | ; RELATIVE NON TERMINAL |
| A202 | A20ø | LDX |  | \#Ø | ; TOKEN MINUS |
| A204 | E9Cl | SBC |  | \# ${ }^{\text {Cl }} 1$ |  |
| A206 | BØØ2 ^A2ØA | BCS |  | : ERN1 | ; BR IF RESULT PLUS |
| A208 | A2FF | LDX |  | \#SFF | ; ADD A MINUS |
| A20A | 18 | : ERN1 | CLC |  |  |
| A20B | 659D | ADC |  | CPC | ; RESULT PLUS CPC |
| A20D | 48 | PHA |  |  | ; IS NEW CPC-1 |
| A20E | 8A | TXA |  |  |  |
| A20F | 659 E | ADC |  | CPC+1 |  |
| A 211 | 48 | PHA |  |  | ; SAVE NEW PC HIGH |
| A212 | 4C28A2 | JMP |  | : PUSH | ; GO PUSH |
|  | = A 215 | : GETADR | EQU | * | ; GET DOUBLE BYTE ADR [-1] |
| A 215 | 20A1A2 | JSR |  | : NXSC | ; GET NEXT CODE |
| A218 | 48 | PHA |  |  | ; SAVE ON STACK |
| A219 | 20AlA2 | JSR |  | : NXSC | ; GET NEXT CODE |
| A21C | 48 | PHA |  |  | ; SAVE ON STACK |
| A21D | 9009 * 4228 | BCC |  | : PUSH | ; BR IF CODE $=\varnothing$ |
| A21F' | 68 | PLA |  |  | ; EXCHANGE TOP |
| A220 | A8 | TAY |  |  | ; 2 ENTRIES ON |
| A221 | 68 | PLA |  |  | ; CPU STACK |
| A222 | AA | TAX |  |  |  |
| A223 | 98 | TYA |  |  |  |
| A224 | 48 | PHA |  |  |  |
| A225 | 8A | TXA |  |  |  |
| A226 | 48 | PHA |  |  |  |
| A227 | 60 | RTS |  |  | ; ELSE GOTO EXTERNAL SRT VIA RTS |
| PUSH |  | ; |  | PUSH TO NEXT | STACK LEVEL |
|  |  | ; |  |  |  |
|  | = A228 | : PUSH | EQU | * |  |
| A228 | A6A9 | LDX |  | STKLVL | ; GET STACK LEVEL |
| A22A | E8 | INX |  |  | ; PLUS 4 |
| A22B | E8 | INX |  |  |  |
| A 22 C | E8 | INX |  |  |  |
| A22D | E8 | INX |  |  |  |
| A 22 E | *A24F | BEQ |  | : SSTB | ; BR STACK TOO BIG |
| A23 $\varnothing$ |  | STX |  | STKLVL | ; SAVE NEW STACK LEVEL |
|  |  | ; |  |  |  |
| A232 | A5F2 | LDA |  | CIX | ; CIX TO |
| A234 | 9D8øØ4 | STA |  | SIX, X | ; STACK IX |
| A237 | A 594 | LDA |  | COX | ; COX TO |
| A239 | 9D8104 | STA |  | SOX, X | ; STACK OX |
| A23C | A59D | LDA |  | CPC | ; CPC TO |
| A23E | 9D8204 | STA |  | SPC, X | : STACK CPC |
| A241 | A 59 E | LDA |  | CPC+1 |  |
| A243 | $9 \mathrm{D8364}$ | STA |  | SPC+1, X |  |
|  |  | ; |  |  |  |
| A246 | 68 | PLA |  |  | ; MOVE STACKED |
| A247 | 859E | STA |  | CPC+1 | : PC TO CPC |
| A249 | 68 | PLA |  |  |  |
| A24A | 859D | STA |  | CPC |  |
| A24C | 4CE2A1 | JMP |  | : NEXT | ; GO FOR NEXT |
| A.24F | 4C24B9 | : SSTB | JMP | ERLTL |  |
| POP |  | ; |  | LOAD CPC FROM | STACK PC |
|  |  | ; |  | AND DECREMENT | TO PREV STACK LEVEL |
|  |  | : POP | EQU | * |  |
| A252 | A6A9 | LDX |  | STKLVL | ; GET STACK LEVEL |
| A254 | Døø1 * 257 | BNE |  | : POP1 | ; BR NOT TOP OF STACK |

## Source Code



Increment CPC


| NXSC |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | ; | GET | NEXT | SYNTAX | X CODE |  |
|  |  | ; |  |  |  |  |  |
| A2Al |  | : NXSC |  |  |  |  |  |
| A2A1 | 209AA2 | JSR | : INCCPC |  | ; | INC PC |  |
| A2A 4 | A2øø | LDX | \#Ø |  |  |  |  |
| A2A6 | A19D | LDA | [CPC, X] |  |  | GET NEXT | CODE |
| A2A 8 | 60 | RTS |  |  |  | RETURN |  |

Source Code

## TERMTST

|  |  | ; | TEST A TERMINAL CODE |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | ; |  |  |  |  |
| A. 2 A 9 |  | : TERMTST |  |  |  |  |
| A2A9 | C9のF | CMP | \# ¢ $^{\text {g }}$ | ; | TEST COD | $\mathrm{E}=\mathrm{F}$ |
| A. 2 AB | FøøD * A2BA | BEQ | : ECHNG | ; | BR CODE | < $F$ |
| A 2 AD | Bø37 * A2E6 | BCS | : SRCONT | ; | BR CODE | > F |
|  |  | ; |  |  |  |  |
| A. 2 AF | 68 | PLA |  | ; | POP RTN | ADR |
| A2B ${ }^{\text {a }}$ | 68 | PLA |  |  |  |  |
| A2B1 | A90C | LDA | \# : EXP-1\&255 | ; | PUSH EXP | ADR |
| A2B3 | 48 | PHA |  | ; | FOR SPEC | IAL |
| A2B4 | A9A6 | LDA | \#: EXP/256 | ; | EXP ANTV | CALL |
| A2B6 | 48 | PHA |  |  |  |  |
| A2B7 | 4C28A2 | JMP | : PUSH |  | GO PUSH |  |

ECHNG

|  |  | ; | EXTERNAL CODE |  | TO CHANGE COX -1 |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | ; |  |  |  |
| A2BA |  | : ECHNG |  |  |  |
| A2BA | 209AA 2 | JSR | : INCCPC | ; | INC PC TO CODE |
| A2BD | Aøøø | LDY | \# $\varnothing$ |  |  |
| A2BF | B19D | LDA. | [CPC], Y | ; | GET CODE |
|  |  | ; |  |  |  |
| A2Cl | A 494 | LDY | COX | ; | GET COX |
| A2C3 | 88 | DEY |  | ; | MINUS 1 |
| A2C4 | 9180 | STA | [OUTBUFF], Y | ; | SET NEW CODE |
| A2C6 | 18 | CLC |  | ; | SET SUCCESS |
| A2C7 | 60 | RTS |  | ; | RETURN |

## SETCODE



Exits for IF and REM

| A2D 4 | A 2 FF | : EIF | LDX | \# SFF |  | RESET STACK |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A2D6 | 9A | TXS |  |  |  |  |
| A2D7 | A 594 | LDA |  | COX |  | SET STMT LENGTH |
| A 2 D 9 | A4A7 | LDY |  | STMLBD |  |  |
| A2DB | 918ø | STA |  | [OUTBUFF], Y |  |  |
| A2DD | 4 CBIA ¢ | JMP |  | : XIF |  | GO CONTINUE IF |
|  |  | ; |  |  |  |  |
| A2EØ |  | : EREM |  |  |  |  |
| A2EØ |  | : EDATA |  |  |  |  |
| A2EØ | A 2 FF | LDX |  | \#\$FF |  | RESET STACK |
| A2E2 | 9A | TXS |  |  |  |  |
| A2E3 | $4 \mathrm{CFBA} \mathrm{V}^{\text {a }}$ | JMP |  | : XDATA |  | O CONTINUE DATA |
| SRCONT |  | ; |  | SEARCH OP | NAME | E TABLE AND TEST RESULT |
|  |  | , |  |  |  |  |
| A2E6 |  | : SRCONT |  |  |  |  |
| A2E6 | 20A1DB | JSR |  | SKPBLANK |  | SKIP BLANKS |
| A2E9 | A5F2 | LDA |  | CIX | ; | GET CURRENT INPUT INDEX |
| A2EB | C5B3 | CMP |  | SVONTX |  | COMPARE WITH SAVED IX |
| A2ED | F016 ^A305 | BEQ |  | : SONT1 |  | BR IF SAVED IX SAME |
| A2EF | 85B3 | STA |  | SVONTX | S | SAVE NEW IX |
|  |  | ; |  |  |  |  |
| A2Fl | A9A7 | LDA |  | \#OPNTAB/256 |  | SET UP FOR ONT |
| A2F3 | AøE3 | LDY |  | \#OPNTAB\&255 | ; | SEARCH |
| A2F5 | A 2 øø | LDX |  | \# $\varnothing$ |  |  |
| A2F7 | 2ø62A4 | JSR |  | SEARCH |  | GO SEARCH |


| A2FA | Bø28 | *A324 | BCS |  | : SONF | ; | BR NOT FOUND |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A2FC | 86B2 |  | STX |  | SVONTL | ; | SAVE NEW CIX |
| A2FE | 18 |  | CLC |  |  |  |  |
| A 2 FF | A5AF |  | LDA |  | STENUM | ; | ADD \$1Ø TO |
| A301 | 6910 |  | ADC |  | \#\$10 | ; | ENTRY NUMBER TO |
| АЗø3 | 85Bø |  | STA |  | SVONTC | ; | GET OPERATOR CODE |
|  |  |  | ; |  |  |  |  |
| A305 | AøøØ |  | : SONTI | LDY | \#Ø |  |  |
| А 307 | B19D |  | LDA |  | [CPC], Y | ; | GET SYNTAX REQ CODE |
| А309 | C5Bø |  | CMP |  | SVONTC | ; | DOES IT MATCH THE FOUND |
| А 30 В | FØØE | A31B | BEQ |  | : SONT2 | ; | BR IF MATCH |
| A30D | C944 |  | CMP |  | \#CNFNP | ; | WAS REQ NFNP |
| A30F | Døø6 | ^A317 | BNE |  | : SONTF | ; | BR IF NOT |
| A311 | A $5 \mathrm{~B} \varnothing$ |  | LDA |  | SVONTC | ; | GET WHAT WE GOT |
| A313 | C944 |  | CMP |  | \#CNFNP | ; | IS IT NFNA |
| A315 | Bøø2 | ^A319 | BCS |  | : SONTS | ; | BR IF IT IS |
| A317 |  |  | : SONTF |  |  |  |  |
| A317 | 38 |  | SEC |  |  | ; | REPORT FAIL |
| A318 | $6 \square$ |  | RTS |  |  |  |  |
| A 319 | A5Bø |  | : SONTS | LDA | SVONTC | ; | GET REAL CODE |
|  |  |  | ; |  |  |  |  |
| A31B | 2øC8A2 |  | : SONT2 | JSR | : SETCODE | ; | GO SET CODE |
| A31E | A6B2 |  | LDX |  | SVONTL | ; | INC CIX BY |
| A32ø | 86 F 2 |  | STX |  | CIX |  |  |
| A322 | 18 |  | CLC |  |  | ; | REPORT SUCCESS |
| A323 | 60 |  | RTS |  |  | ; | DONE |
| A324 | А9øø |  | : SONF | LDA | \#Ø | ; | SET ZERO AS |
| A326 | 85Bø |  | STA |  | SVONTC | , | SAVED CODE |
| A328 | 38 |  | SEC |  |  |  |  |
| A329 | 60 |  | RTS |  |  | ; | DONE |


| TVAR |  | ; |  | EXTERNAL |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | ROUTINE FOR TNVAR \& TSVAR |  |
| A32A | А900 |  |  | : TNVAR | LDA | \#Ø | ; SET NUMERIC TEST |
| A32C | Føø2 * ${ }^{\text {A330 }}$ | BEQ |  | : TVAR |  |
|  |  | ; |  |  |  |
| A32E | A980 | : TSVAR | LDA | \# \$80 | ; SET STR TEST |
|  |  | ; |  |  |  |
| A330 | 85D2 | : TVAR | STA | TVTYPE | ; SAVE TEST TYPE |
| A332 | 20AlDB | JSR |  | SKPBLANK | ; SKIP LEADING BLANKS |
| A335 | A5F2 | LDA |  | CIX | ; GET INDEX |
| A337 | 85AC | STA |  | TVSCIX | ; FOR SAVING |
|  |  | ; |  |  |  |
| A339 | 20F3A3 | JSR |  | : TSTALPH | ; GO TEST FIRST CHAR |
| A33C | Bø25 ^A363 | BCS |  | : TVFAIL | ; BR NOT ALPHA |
| A33E | 20E6A2 | JSR |  | : SRCONT | ; IF THIS IS A |
| A341 | A5B6 | LDA |  | SVONTC | ; RESVD NAME |
| A343 | Føø8 *A34D | BEQ |  | :TV1 | ; BR NOT RSVDNAME |
| A 345 | A4B2 | LDY |  | SVONTL | ; IF NEXT CHAR AFTER |
| A347 | B1F3 | LDA |  | [INBUFF], Y | ; RESERVED NAME |
| A 349 | C930 | CMP |  | \# ${ }^{\text {3 }}$ ¢ | ; NOT ALARM NUMERIC |
| A 34 B | 9016 ^A363 | BCC |  | : TVFAIL | ; THEN ERROR |
|  |  | , |  |  |  |
| A34D | E6F2 | :TV1 | INC | CIX | ; INC TO NEXT CHAR |
| A 34 F | 20F3A3 | JSR |  | : TSTALPH | ; TEST ALPHA |
| A352 | 9ØF9 ^A34D | BCC |  | : TV1 | ; BR IF ALPHA |
| A 354 | 20AFDB | JSR |  | TSTNUM | ; TRY NUMBER |
| A 357 | 9ØF4 ^A34D | BCC |  | :TV1 | ; BR IF NUMBER |
| A359 |  | ; |  | NBUFF] |  |
| A359 | B1F3 | LDA |  | [INBUFF], Y | ; GET OFFENDING CHAR |
| A35B | C924 | CMP |  | \#'\$' | ; IS IT \$ |
| A35D | Fの06 ~A365 | BEQ |  | : TVSTR | ; BR IF \$ [STRING] |
| A 35 F | 24D2 | BIT |  | TVTYPE | ; THIS A NVAR SEARCH |
| A361 | 1009 *A36C | BPL |  | : TVOK | ; BR 'IF NVAR |
|  |  | ; |  |  |  |
| A363 | 38 | : TVFAIL | SEC |  | ; SET FAIL CODE |
| A364 | $6 \emptyset$ | RTS |  |  | ; DONE |
| A365 | 24D2 | ; TVSTR | BIT | TVTYPE | ; TEST SVAR SEARCH |
| A367 | 10FA *A363 | BPL |  | : TVFAIL | ; BR IF SVAR |

Source Code

| A369 | C8 | INY |  |  | ; INC OVER \$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| A36A | DøøD ^A379 | BNE |  | :TVOK2 | ; BR ALWAYS |
| A36C | BlF3 | :TVOK | LDA | [INBUFF], Y | ; GET NEXT CHAR |
| A36E | C928 | CMP |  | \#' ${ }^{\prime}$ | ; IS It PAREN |
| A370 | Dø07 ^A379 | BNE |  | :TVOK2 | ; BR NOT PAREN |
| A372 | C8 | INY |  |  | ; INC OVER PAREN |
| A373 | A940 | LDA |  | \#\$4ø | ; OR IN ARRAY |
| A375 | 65D2 | ORA |  | TVTYPE | ; CODE TO TVTYPE |
| A377 | 85D2 | STA |  | TVTYPE |  |
| A379 | A5AC | :TVOK2 | LDA | TVSCIX | ; GET SAVED CIX |
| A37B | 85F2 | STA |  | CIX | ; PUT BACK |
| A37D | 84AC | STY |  | TVSCIX | ; SAVE NEW CIX |
| A37F | A583 | LDA |  | VNTP+1 | ; SEARCH VNT |
| A381 | A482 | LDY |  | VNTP | ; FOR THIS GUY |
| A383 | A2øø | LDX |  | \# $\varnothing$ |  |
| A385 | 2062A4 | JSR |  | SEARCH |  |
| A388 |  | : TVRS |  |  |  |
| A388 | BøøA ^A394 | BCS |  | :TVSø | ; BR NOT FOUND |
| A38A | E4AC | CPX |  | TVSCIX | ; FOUND RIGHT ONE |
| A38C | F04D ^A3DB | BEQ |  | :TVSUC | ; BR IF YES |
| A38E | 2090A4 | JSR |  | SRCNXT | ; GO SEARCH MORE |
| A391 | 4C88A3 | JMP |  | :TVRS | ; TEST THIS RESULT |
| A394 |  | ; ${ }^{\text {; TVSø }}$ |  |  |  |
| A394 | 38 | SEC |  |  | ; SIGH: |
| A395 | A5AC | LDA |  | TVSCIX | ; VAR LENGTH IS |
| A397 | E5F2 | SBC |  | CIX | ; NEW CIX-OLD CIX |
| A399 | 85F2 | STA |  | CIX |  |
|  |  | ; |  |  |  |
| A39B | A8 | TAY |  |  | ; GO EXPAND VNT |
| A39C | A284 | LDX |  | \#VNTD | ; BY VAR LENGTH |
| A39E | 207FA8 | JSR |  | EXPLOW |  |
| A3Al | A5AF | LDA |  | STENUM | ; SET VARIABLE NUMBER |
| A3A3 | 85D3 | STA |  | TVNUM |  |
| A3A5 | A4F2 | LDY |  | CIX | AND |
| A3A 7 | 88 | DEY |  |  |  |
| A3A8 | A6AC | LDX |  | TVSCIX | ; GET DISPL TO EQU+1 |
| A3AA | CA | DEX |  |  |  |
| А 3 AB | BD8øø5 | :TVS 1 | LDA | LBUFF, X | MOVE VAR TO |
| A3AE | 9197 | STA |  | [SVESA], Y |  |
| А3Bø | CA | DEX |  |  |  |
| A3B1 | 88 | DEY |  |  |  |
| A3B2 | 10F7 ^A3AB | BPL |  | :TVS1 |  |
| A3B4 | A4F2 | LDY |  | CIX | ; TURN ON MSB |
| A3B6 | 88 | DEY |  |  | ; OF LAST CHAR |
| А 3 B 7 | B197 | LDA |  | [SVESA], Y | ; IN VTVT ENTRY |
| А389 | ø980 | ORA |  | \#\$8® |  |
| A3BB | 9197 | STA |  | [SVESA], Y |  |
| A3BD | Аøø8 | LDY |  | \#8 | ; THEN EXPAND |
| A3BF | A288 | LDX |  | \#STMTAB | ; VVT BY 8 |
| A3C1 | 2ø7FA8 | JSR |  | EXPLOW |  |
| А 3 C 4 | E6Bl | INC |  | SVVVTE | ; INC VVT EXP SIZE |
| A3C6 | Aøø2 | LDY |  | \#2 | ; Clear value |
| A3C8 | A9øø | LDA |  | \#ø | ; PART OF |
| A3CA | 99D2øб | :TVS1A | STA | TVTYPE, Y | ; ENTRY |
| A3CD | C8 | INY |  |  |  |
| A3CE | сøø8 | CPY |  | \#8 |  |
| A3Dø | $90 \mathrm{F8}{ }^{\text {^ }} \mathrm{A} 3 \mathrm{CA}$ | BCC |  | :TVSlA |  |
| A3D2 | 88 | DEY |  |  | ; AND THEN |
| A3D3 | B9D2øø | :TVS2 | LDA | TVTYPE, Y | ; PUT IN VAR table |
| A3D6 | 9197 | STA |  | [SVESA], Y | ; ENTRY |
| A3D8 | 88 | DEY |  |  |  |
| A3D9 | 10 F 8 ^A3D3 | BPL |  | :TVS2 |  |

## Source Code

| A3DB | 24D2 | : TVSUC | BIT | TVTYPE | ; | WAS THERE A PAREN |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A 3DD | 50ø2 ^A3E1 | BVC |  | : TVNP | ; | BR IF NOT |
| A3DF | C6AC | DEC |  | TVSCIX | ; | LET SYNTAX SEE PAREN |
|  |  | ; |  |  |  |  |
| A3E1 | A5AC | : TVNP | LDA | TVSCIX | ; | GET NEW CIX |
| A3E3 | 85F2 | STA |  | CIX | ; | TO CIX |
|  |  | ; |  |  |  |  |
| A3E5 | ${ }^{\wedge} \mathrm{A} 3 \mathrm{~F} \varnothing$ | LDA |  | STENUM | ; | GET TABLE ENTRY NO |
| A 3 E 7 |  | BMI |  | :TVFULL | ; | BR IF > \$7F |
| A3E9 | 0980 | ORA |  | \# \$80 | ; | MAKE IT > \$7F |
| A3EB | 2ØC8A2 | JSR |  | : SETCODE | ; | SET CODE TO OUTPUT BUFFER |
| A3EE | 18 | CLC |  |  | ; | SET SUCCESS CODE |
| A3EF | 60 | RTS |  |  | ; | RETURN |
|  |  | : TVFULL | JMP | ERRVSF |  |  |
| A3FØ | 4C38B9 |  |  |  | ; GO TO ERROR RTN |  |

TSTALPH

|  |  |  | ; | TEST |
| :---: | :---: | :---: | :---: | :---: |
|  |  |  | ; |  |
| A3F 3 |  |  | : TSTALPH |  |
| A3F3 | A 4 F 2 |  | LDY | CIX |
| A3F5 | BlF3 |  | LDA | [INBUFF], Y |
| A3F7 |  |  | TSTALPH |  |
| A3F7 | C941 |  | CMP | \#'A |
| A3F9 | 9 Ø03 | * A 3FE | BCC | : TAFAIL |
| A3FB | C95B |  | CMP | \# \$5B |
| A3FD | 60 |  | RTS |  |
|  |  |  | ; |  |
| A3FE | 38 |  | :TAFAIL |  |
| A 3FF | 60 |  | RTS |  |

TNCON

| A4øø |  |
| :---: | :---: |
| А4øø | 20A1DB |
| A403 | A5F2 |
| A405 | 85AC |
| A407 | 2øøøD8 |
| A 40 A | 9005 *A411 |
| A40C | A5AC |
| A40E | 85F2 |
| A410 | 60 |


| ; | EXTERNAL SUBROUTINE TO CHECK FOR |  |  |
| :---: | :---: | :---: | :---: |
| ; |  |  |  |
| : TNCON |  |  |  |
| JSR |  | SKBLANK |  |
| LDA |  | CIX |  |
| STA |  | TVSCIX |  |
| JSR |  | CVAFP | ; GO TEST AND CONV |
| BCC |  | :TNC1 | ; BR IF NUMBER |
| LDA |  | TVSCIX |  |
| STA |  | CIX |  |
| RTS |  |  | ; RETURN FAIL |
|  |  |  |  |
| : TNC1 | LDA | \# ${ }^{\text {¢ }}$ ¢E | ; SET NUMERIC CONST |
| JSR |  | : SETCODE |  |
| ; |  |  |  |
| LDY |  | Cox |  |
| LDX |  | \# $\emptyset$ |  |
| : TNC 2 | LDA | FRø, X | ; MOVE CONST TO STMT |
| STA |  | [OUTBUFF], Y |  |
| INY |  |  |  |
| INX |  |  |  |
| CPX |  | \# 6 |  |
| BCC |  | : TNC2 |  |
| STY |  | COX |  |
| CLC |  |  |  |
| RTS |  |  |  |
|  |  |  |  |
| ; |  | EXT SRT TO | CHECK FOR STR CONST |
| : TSCON |  |  |  |
|  |  |  |  |
| JSR |  | SKBLANK |  |
| LDY |  | CIX | ; GET INDEX |
| LDA |  | [INBUFF], Y | ; GET CHAR |
| CMP |  | \#\$22 | ; IS IT DQUOTE |
| BEQ |  | : TSC1 | ; BR IF DQ |
| SEC |  |  | ; SET FAIL |
| RTS |  |  | ; RETURN |


| TSCON |  |  |
| :--- | :--- | :--- |
|  |  |  |
| A428 |  |  |
| A428 | 2ØA1DB |  |
| A42B | A4F2 |  |
| A42D | B1F3 |  |
| A42F | C922 |  |
| A431 | Føø2 | A435 |
| A433 | 38 |  |
| A434 | $6 \varnothing$ |  |

Source Code

| A 435 | A90F | :TSCl | LDA | \# ¢ $^{\text {F }}$ | ; | SET SCON CODE |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A437 | 2øC8A2 | JSR |  | : SETCODE |  |  |
| A43A | A594 | LDA |  | COX | ; | SET COX |
| A43C | 85AB | STA |  | TSCOX | ; | SAVE FOR LENGTH |
| A43E | 2øС8A 2 | JSR |  | : SETCODE | ; | SET DUMMY FOR NOW |
| A441 | E6F2 | :TSC2 | INC | CIX | ; | NEXT INPUT CHAR |
| A443 | A 4 F 2 | LDY |  | CIX |  |  |
| A445 | BlF3 | LDA |  | [INBUFF], Y |  |  |
| A447 | C99B | CMP |  | \#CR | ; | IS IT CR |
| A449 | FøøC ^A457 | BEQ |  | : TSC4 | ; | BR IF CR |
| A44B | C922 | CMP |  | \#\$22 | ; | IS IT DQ |
| A44D | FøØ6 ^A455 | BEQ |  | : TSC3 | ; | BR IF DQ |
| A44F | 2øC8A2 | JSR |  | : SETCODE | ; | OUTPUT IT |
| A452 | 4C41A4 | JMP |  | : TSC2 | ; | NEXT |
| A 455 | E6F2 | ; TSC3 | INC | CIX | ; | INC CIX OVER DQ |
| A457 | 18 | :TSC4 | CLC |  |  |  |
| A458 | A 594 | LDA |  | COX | ; | LENGTH IS COX MINUS |
| A45A | E5AB | SBC |  | TSCOX | ; | LENGTH BYTE COX |
| A 45 C | A 4 AB | LDY |  | TSCOX |  |  |
| A 45 E | 9180 | STA |  | [OUTBUFF], Y | ; | SET LENGTH |
|  |  | ; |  |  |  |  |
| A460 | 18 | CLC |  |  | ; | SET SUCCESS |
| A461 | $6 \emptyset$ | RTS |  |  | ; | DONE |

## Search a Table

; TABLE FORMAT:

```
```

;

```
```

;
GARBAGE TO SKIP [N]
GARBAGE TO SKIP [N]
GARBAGE TO SKIP [N]
GARBAGE TO SKIP [N]
WITH LEAST SIGNIFICANT BYTE HAVING
WITH LEAST SIGNIFICANT BYTE HAVING
MOST SIGNIFICANT BIT ON
MOST SIGNIFICANT BIT ON
LAST TABLE ENTRY MUST HAVE FIRST ASCII
LAST TABLE ENTRY MUST HAVE FIRST ASCII
CHAR = Ø
CHAR = Ø
ENTRY PARMS:
ENTRY PARMS:
Y. = SKIP LENGTH
Y. = SKIP LENGTH
A,Y = TABLE ADR [HIGH LOW]
A,Y = TABLE ADR [HIGH LOW]
ARGUMENT = INBUFF + CIX
ARGUMENT = INBUFF + CIX
EXIT PARMS:
EXIT PARMS:
CARRY = CLEAR IF FOUND
CARRY = CLEAR IF FOUND
X = FOUND ARGUMENT END CIX+1
X = FOUND ARGUMENT END CIX+1
SRCADR = TABLE ENTRY ADR
SRCADR = TABLE ENTRY ADR
STENUM = TABLE ENTRY NUMBER
STENUM = TABLE ENTRY NUMBER
SEARCH
SEARCH
STX
STX
; LDX
; LDX
STX
STX
;:SRCl STA SRCADR+1 ; SET SEARCH ADR

```
;:SRCl STA SRCADR+1 ; SET SEARCH ADR
```

```
STY SRCADR ; INC ENTRY NUMBER
```

STY SRCADR ; INC ENTRY NUMBER
INC STENUM CIX GDX ARG DISPL
INC STENUM CIX GDX ARG DISPL
LDX SRCSKP ; GET SKIP LENGTH
LDX SRCSKP ; GET SKIP LENGTH
LDA [SRCADR],Y ; GET FIRST CHAR
LDA [SRCADR],Y ; GET FIRST CHAR
BEQ :SRCNF ; BR IF EOT
BEQ :SRCNF ; BR IF EOT
LDA \# | SET STATUS = EQ
LDA \# | SET STATUS = EQ
PHP ; AND PUSH IT
PHP ; AND PUSH IT
;:SRC2 LDA LBUFF,X ; GET INPUT CHAR
;:SRC2 LDA LBUFF,X ; GET INPUT CHAR
AND \#\$7F ; TURN OFF MSB
AND \#\$7F ; TURN OFF MSB
\#\$7F ; I IF WILD CARD
\#\$7F ; I IF WILD CARD
:SRC5 ; THEN BR
:SRC5 ; THEN BR
[SRCADR],Y ; EX-OR WITH TABLE CHAR
[SRCADR],Y ; EX-OR WITH TABLE CHAR
; SHIFT MSB TO CARRY
; SHIFT MSB TO CARRY
A
A
:SRC3 ; BR IF [ARG=TAB] CHAR

```
        :SRC3 ; BR IF [ARG=TAB] CHAR
```

$\begin{array}{lll}\text { A462 } & & \\ \text { A462 } & \text { 86AA } \\ & & \\ \text { A464 } & \text { A2FF } & \\ \text { A466 } & 86 A F & \\ & & \\ \text { A468 } & 8596 & \\ \text { A46A } & 8495 & \\ \text { A46C } & \text { E6AF } & \\ \text { A46E } & \text { A6F2 } & \\ \text { A47Ø } & \text { A4AA } & \\ \text { A472 } & \text { B195 } & \\ \text { A474 } & \text { FØ27 } & \text { A49D } \\ \text { A476 } & \text { A9Øø } & \\ \text { A478 } & \text { Ø8 } & \\ & & \\ \text { A479 } & \text { BD8øØ5 } \\ \text { A47C } & 297 F & \\ \text { A47E } & \text { C92E } & \\ \text { A48Ø } & \text { FØ1D } & \text { A49F } \\ \text { A482 } & & \\ \text { A482 } & 5195 & \\ \text { A484 } & & \\ \text { A484 } & \text { +ØA } & \\ \text { A485 } & \text { FØØ2 } & \text { A489 }\end{array}$
$\begin{array}{lll}\text { A462 } & & \\ \text { A462 } & \text { 86AA } \\ & & \\ \text { A464 } & \text { A2FF } & \\ \text { A466 } & 86 A F & \\ & & \\ \text { A468 } & 8596 & \\ \text { A46A } & 8495 & \\ \text { A46C } & \text { E6AF } & \\ \text { A46E } & \text { A6F2 } & \\ \text { A47Ø } & \text { A4AA } & \\ \text { A472 } & \text { B195 } & \\ \text { A474 } & \text { FØ27 } & \text { A49D } \\ \text { A476 } & \text { A9Øø } & \\ \text { A478 } & \text { Ø8 } & \\ & & \\ \text { A479 } & \text { BD8øØ5 } \\ \text { A47C } & 297 F & \\ \text { A47E } & \text { C92E } & \\ \text { A48Ø } & \text { FØ1D } & \text { A49F } \\ \text { A482 } & & \\ \text { A482 } & 5195 & \\ \text { A484 } & & \\ \text { A484 } & \text { +ØA } & \\ \text { A485 } & \text { FØØ2 } & \text { A489 }\end{array}$
$\begin{array}{lll}\text { A462 } & & \\ \text { A462 } & \text { 86AA } \\ & & \\ \text { A464 } & \text { A2FF } & \\ \text { A466 } & 86 A F & \\ & & \\ \text { A468 } & 8596 & \\ \text { A46A } & 8495 & \\ \text { A46C } & \text { E6AF } & \\ \text { A46E } & \text { A6F2 } & \\ \text { A47Ø } & \text { A4AA } & \\ \text { A472 } & \text { B195 } & \\ \text { A474 } & \text { FØ27 } & \text { A49D } \\ \text { A476 } & \text { A9Øø } & \\ \text { A478 } & \text { Ø8 } & \\ & & \\ \text { A479 } & \text { BD8øØ5 } \\ \text { A47C } & 297 F & \\ \text { A47E } & \text { C92E } & \\ \text { A48Ø } & \text { FØ1D } & \text { A49F } \\ \text { A482 } & & \\ \text { A482 } & 5195 & \\ \text { A484 } & & \\ \text { A484 } & \text { +ØA } & \\ \text { A485 } & \text { FØØ2 } & \text { A489 }\end{array}$
$\begin{array}{lll}\text { A462 } & & \\ \text { A462 } & \text { 86AA } \\ & & \\ \text { A464 } & \text { A2FF } & \\ \text { A466 } & 86 A F & \\ & & \\ \text { A468 } & 8596 & \\ \text { A46A } & 8495 & \\ \text { A46C } & \text { E6AF } & \\ \text { A46E } & \text { A6F2 } & \\ \text { A47Ø } & \text { A4AA } & \\ \text { A472 } & \text { B195 } & \\ \text { A474 } & \text { FØ27 } & \text { A49D } \\ \text { A476 } & \text { A9Øø } & \\ \text { A478 } & \text { Ø8 } & \\ & & \\ \text { A479 } & \text { BD8øØ5 } \\ \text { A47C } & 297 F & \\ \text { A47E } & \text { C92E } & \\ \text { A48Ø } & \text { FØ1D } & \text { A49F } \\ \text { A482 } & & \\ \text { A482 } & 5195 & \\ \text { A484 } & & \\ \text { A484 } & \text { +ØA } & \\ \text { A485 } & \text { FØØ2 } & \text { A489 }\end{array}$
$\begin{array}{lll}\text { A462 } & & \\ \text { A462 } & \text { 86AA } \\ & & \\ \text { A464 } & \text { A2FF } & \\ \text { A466 } & 86 A F & \\ & & \\ \text { A468 } & 8596 & \\ \text { A46A } & 8495 & \\ \text { A46C } & \text { E6AF } & \\ \text { A46E } & \text { A6F2 } & \\ \text { A47Ø } & \text { A4AA } & \\ \text { A472 } & \text { B195 } & \\ \text { A474 } & \text { FØ27 } & \text { A49D } \\ \text { A476 } & \text { A9Øø } & \\ \text { A478 } & \text { Ø8 } & \\ & & \\ \text { A479 } & \text { BD8øØ5 } \\ \text { A47C } & 297 F & \\ \text { A47E } & \text { C92E } & \\ \text { A48Ø } & \text { FØ1D } & \text { A49F } \\ \text { A482 } & & \\ \text { A482 } & 5195 & \\ \text { A484 } & & \\ \text { A484 } & \text { +ØA } & \\ \text { A485 } & \text { FØØ2 } & \text { A489 }\end{array}$
$\begin{array}{lll}\text { A462 } & & \\ \text { A462 } & 86 A A & \\ & & \\ \text { A464 } & \text { A2FF } & \\ \text { A466 } & 86 A F & \\ & & \\ \text { A468 } & 8596 & \\ \text { A46A } & 8495 & \\ \text { A46C } & \text { E6AF } & \\ \text { A46E } & \text { A6F2 } & \\ \text { A47Ø } & \text { A4AA } & \\ \text { A472 } & \text { B195 } & \\ \text { A474 } & \text { FØ27 } & \text { A49D } \\ \text { A476 } & \text { A9ØØ } & \\ \text { A478 } & \text { Ø8 } & \\ & & \\ \text { A479 } & \text { BD8øø5 } \\ \text { A47C } & 297 F & \\ \text { A47E } & \text { C92E } & \\ \text { A48Ø } & \text { FØ1D } & \text { A49F } \\ \text { A482 } & & \\ \text { A482 } & 5195 & \\ \text { A484 } & & \\ \text { A484 } & \text { +ØA } & \\ \text { A485 } & \text { FØØ2 } & \text { A489 }\end{array}$
$\begin{array}{lll}\text { A462 } & & \\ \text { A462 } & 86 A A & \\ & & \\ \text { A464 } & \text { A2FF } & \\ \text { A466 } & 86 A F & \\ & & \\ \text { A468 } & 8596 & \\ \text { A46A } & 8495 & \\ \text { A46C } & \text { E6AF } & \\ \text { A46E } & \text { A6F2 } & \\ \text { A47Ø } & \text { A4AA } & \\ \text { A472 } & \text { B195 } & \\ \text { A474 } & \text { FØ27 } & \text { A49D } \\ \text { A476 } & \text { A9ØØ } & \\ \text { A478 } & \text { Ø8 } & \\ & & \\ \text { A479 } & \text { BD8øø5 } \\ \text { A47C } & 297 F & \\ \text { A47E } & \text { C92E } & \\ \text { A48Ø } & \text { FØ1D } & \text { A49F } \\ \text { A482 } & & \\ \text { A482 } & 5195 & \\ \text { A484 } & & \\ \text { A484 } & \text { +ØA } & \\ \text { A485 } & \text { FØØ2 } & \text { A489 }\end{array}$
$\begin{array}{lll}\text { A462 } & & \\ \text { A462 } & \text { 86AA } \\ & & \\ \text { A464 } & \text { A2FF } & \\ \text { A466 } & 86 A F & \\ & & \\ \text { A468 } & 8596 & \\ \text { A46A } & 8495 & \\ \text { A46C } & \text { E6AF } & \\ \text { A46E } & \text { A6F2 } & \\ \text { A47Ø } & \text { A4AA } & \\ \text { A472 } & \text { B195 } & \\ \text { A474 } & \text { FØ27 } & \text { A49D } \\ \text { A476 } & \text { A9Øø } & \\ \text { A478 } & \text { Ø8 } & \\ & & \\ \text { A479 } & \text { BD8øØ5 } \\ \text { A47C } & 297 F & \\ \text { A47E } & \text { C92E } & \\ \text { A48Ø } & \text { FØ1D } & \text { A49F } \\ \text { A482 } & & \\ \text { A482 } & 5195 & \\ \text { A484 } & & \\ \text { A484 } & \text { +ØA } & \\ \text { A485 } & \text { FØØ2 } & \text { A489 }\end{array}$

## Source Code



## Statement Name Table



Source Code

| A4E5 | 4E4558D4 | DC | 'NEXT ' |
| :---: | :---: | :---: | :---: |
| A4E9 | BCA6 | DW | : SGOTO-1 |
| A4EB | 474F54CF | DC | 'GOTO' |
| A 4 EF | BCA6 | ; DW | - SGOTO-1 |
| A4F1 | 474F2054CF | DC | 'GO TO' |
|  |  | ; |  |
| A4F6 | BCA6 | DW | : SGOSUB-1 |
| A4F8 | 474F5355C2 | DC | 'GOSUB' |
| A4FD | BCA6 | DW | : STRAP-1 |
| A4FF | 545241Dø | DC | 'TRAP' |
|  |  | ; |  |
|  |  | ; |  |
| A503 | BDA6 | DW | : SBYE-1 |
| A505 | 4259C5 | DC | 'BYE' |
| A508 | BDA6 | DW | : SCONT-1 |
| A50A | 434F4ED4 | DC | 'CONT' |
| A50E | 5FA7 | DW | : SCOM-1 |
| A510 | 434 FCD | DC | ' $\mathrm{COM}^{\prime}$ |
|  |  | ; |  |
|  |  | ; |  |
| A513 | 20A7 | DW | : SCLOSE-1 |
| A515 | 434C4F53C5 | DC | 'Close' |
| A51A | BDA6 | DW | : SCLR-1 |
| A51C | 434CD2 | DC | 'CLR' |
| A51F | BDA6 | DW | :SDEG-1 |
| A521 | 4445C7 | DC | 'DEG ' |
|  | 5FA7 | ; |  |
|  |  | DW | :SDIM-1 |
| A526 | 4449 CD | DC | 'DIM ${ }^{\prime}$ |
| A529 | BDA6 | DW | :SEND-1 |
| A52B | 454EC4 | DC | 'END' |
|  |  | ; |  |
| A52E | BDA6 | DW | :SNEW-1 |
| A530 | 4E45D7 | DC | 'NEW' |
| A535 | 4F5®45CE | DC | 'OPEN ${ }^{\text {' }}$ |
| A539 | 23A7 | DW | :SLOAD-1 |
| A53B | 4C4F41C4 | DC | 'LOAD' |
| A53F | 23A7 | DW | :SSAVE-1 |
| A541 | 534156C5 | DC | 'SAVE' |
| A545 | 40A7 | DW | :SSTATUS-1 |
| A547 | 5354415455 | DC | 'Status' |
| A54D | 49A7 | DW | :SNOTE-1 |
| A54F | 4E4F54C5 | DC | 'note' |
| A553 | 49A7 | DW | : SPOINT-1 |
| A555 | 504F494ED4 | DC | 'POINT' |
| A55A | 17A7 | DW | : SxIO-1 |
| A55C | 5849CF | DC | 'XIO' |
| A55F | 62A7 | DW | : SON-1 |
| A561 | 4FCE | DC | 'ON' |
| A563 | $5 \mathrm{CA7}$ | DW | -SPOKE-1 |
| A565 | 5ø4F4BC5 | DC | 'POKE ' |
|  |  | ; ${ }^{\text {W }}$ |  |
| A569 | FBA6 | DW | :SPRINT-1 |
| A56B | 5ø52494ED4 | DC | 'PRINT ' |
|  |  | ; |  |
| A57ø | BDA6 | DW | : SRAD-1 |
| A572 | 5241C4 | DC | 'RAD' |
| A575 | F4A6 | DW | :SREAD-1 |



## Syntax Tables

Syntax Table OP Codes

| $=\varnothing \varnothing \square \varnothing$ | : ANTV | EQU | \$øØ | ; | ABSOLUTE NON TERMINAL VECTOR FOLLOWED BY 2 BYTE ADR -1 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| = øøø1 | : ESRT | EQU | \$Ø1 | ; | EXTERNAL SUBROUTINE CALL |
|  | ; |  |  |  | FOLLOWED BY 2 BYTE ADR -1 |
| = øøø2 | : OR | EQU | \$ø2 | ; | ALTERNATIVE, BNF OR (]) |
| = Фøø3 | : RTN | EQU | \$03 | ; | RETURN, (\#) |
| $=\emptyset \emptyset 04$ | : NULL | EQU | \$ø4 | ; | ACCEPT TO THIS POINT ( \& ) |
| $=\varnothing \varnothing \square \mathrm{E}$ | : VEXP | EQU | \$ $\varnothing \mathrm{E}$ | ; | SPECIAL NTV FOR EXP (<EXP>) |
| = øøøF | : CHNG | EQU | \$øF | ; | CHANGE LAST OUTPUT TOKEN |

$<\operatorname{EXP}>=(<\operatorname{EXP}>)<$ NOP $>|<U N A R Y><E X P>|<N V><$ NOP $>\#$

| A60D |  | : EXP SYN | CLPRN |
| :---: | :---: | :---: | :---: |
| A60D | $+2 \mathrm{~B}$ | DB | CLPRN |
| A6øE |  | SYN | JS, : EXP |
| A60E | +BF | DB | \$8Ø+( ( $:$ EXP-*)\&\$7F) XOR \$40 ) |
| A60F |  | SYN | CRPRN |
| A60F | $+2 \mathrm{C}$ | DB | CRPRN |
| A610 |  | SYN | JS, : NOP |
| A610 | $+D E$ | DB | \$8Ø+( ( (:NOP-*) \& 7 F ) XOR \$4Ø ) |
| A611 |  | SYN | : OR |
| A611 | +ø2 | DB | : OR |
| A612 |  | SYN | JS, : UNARY |
| A612 | +C6 | DB | \$8Ø+(():UNARY-*)\&\$7F) XOR \$4Ø.) |
| A613 |  | SYN | JS, : EXP |
| A613 | +BA | DB | \$8Ø+( ( $:$ EXP-*)\&\$7F) XOR \$4Ø) |
| A614 |  | SYN | : OR |
| A614 | $+\varnothing 2$ | DB | : OR |
| A615 |  | SYN | JS, : NV |
| A615 | +CD | DB | \$80+(():NV-*)\&\$7F) XOR \$40 ) |
| A616 |  | SYN | JS, : NOP |
| A616 | +D8 | DB | \$8Ø+( ( (:NOP-*) \& \$7F) XOR \$4Ø ) |
| A617 |  | SYN | : RTN |
| A617 | $+\emptyset 3$ | DB | : RTN |

$<$ UNARY $>=+|\cdot|$ NOT\#

| A618 |  | : UNARY SYN | CPLUS |  |
| :---: | :---: | :---: | :---: | :---: |
| A618 | +25 | DB | CPLUS |  |
| A619 |  | SYN | : CHNG, CUPLUS |  |
| A619 | $+\emptyset F$ | DB | : CHNG |  |
| A61A | +35 | DB | CUPLUS |  |
| A61B |  | SYN | : OR |  |
| A61B | $+\square 2$ | DB | : OR |  |
| A61C |  | SYN | CMINUS |  |
| A61C | $+26$ | DB | CMINUS |  |
| A61D |  | SYN | : CHNG, CUMINUS |  |
| A61D | $+\emptyset F$ | DB | : CHNG |  |
| A61E | $+36$ | DB | CUMINUS |  |
| A61F |  | SYN | : OR |  |
| A61F | $+\varnothing 2$ | DB | : OR |  |
| A620 |  | SYN | CNOT |  |
| A62ø | $+28$ | DB | CNOT |  |
| A621 |  | SYN | : RTN |  |
| A621 | $+\square 3$ | DB | : RTN |  |
| $<\mathrm{NV}>=<\mathrm{NFUN}>$ |  | $<$ NVAR $>$ | <NCON> $\mid<$ STCOMP $>\#$ |  |
| A622 |  | : NV SYN | JS, : NFUN, : OR |  |
| A622 | +FD | DB | \$8ø+( ( $:$ NFUN-*) \&\$7F) XOR | \$40) |
| A623 | $+\varnothing 2$ | DB | : OR |  |
| A624 |  | SYN | JS, : NVAR, : OR |  |
| A624 | +E8 | DB | \$8ø+(():NVAR-*) 8 \$ 7 F ) XOR | \$40) |
| A625 | $+\varnothing 2$ | DB | : OR |  |
| A626 |  | SYN | : ESRT, AD, :TNCON-1, : OR |  |
| A626 | $+\emptyset 1$ | DB | : ESRT |  |


| A627 | +FFA3 | DW | (: TNCON-1) |
| :---: | :---: | :---: | :---: |
| A629 | $+\emptyset 2$ | DB | : OR |
| A62A |  | SYN | :ANTV, AD, : STCOMP-1 |
| A62A | $+\emptyset \emptyset$ | DB | : ANTV |
| A62B | +7DA6 | DW | (: STCOMP-1) |
| A62D |  | SYN | : RTN |
| A62D | +Ø3 | DB | : RTN |
| $<$ NOP $>=<\mathbf{O P}><\mathbf{E X P}>\mid \& \#$ |  |  |  |
| A62E |  | :NOP SYN | JS, OP |
| A62E | +C4 | DB | \$8Ø+( ( $:$ OP-*) \& $\left.{ }^{\text {P }} 7 \mathrm{~F}\right)$ XOR \$4Ø) |
| A62F |  | SYN | JS , : EXP |
| A62F | $+9 \mathrm{E}$ | DB |  |
| A630 |  | SYN | : OR |
| A630 | $+\emptyset 2$ | DB | : OR |
| A631 |  | SYN | : RTN |
| A631 | +ø3 | DB | : RTN |
| $<\mathbf{O P}>=\left.{ }^{*}\right\|^{*}\|/\|<=\|\mathbf{S}=\|<>\|<\|>\|=\|$ AND $\|$ OR\# |  |  |  |
| A632 |  | : OP SYN | CEXP, : OR |
| A632 | +23 | DB | CEXP |
| A633 | +Ø2 | DB | : OR |
| A634 |  | SYN | CPLUS, : OR |
| A634 | $+25$ | DB | CPLUS |
| A635 | $+\emptyset 2$ | DB | : OR |
| A636 |  | SYN | CMINUS, : OR |
| A636 | +26 | DB | CMINUS |
| A637 | +ø2 | DB | : OR |
| A638 |  | SYN | CMUL, : OR |
| A638 | +24 | DB | CMUL |
| A639 | $+\emptyset 2$ | DB | : OR |
| A63A |  | SYN | CDIV, : OR |
| A63A | $+27$ | DB | CDIV |
| A63B | $+\varnothing 2$ | DB | : OR |
| A63C |  | SYN | CLE, OR |
| A63C | +1D | DB | CLE |
| A63D | $+\varnothing 2$ | DB | : OR |
| A63E |  | SYN | CGE, : OR |
| A63E | +1F | DB | CGE |
| A63F | $+\varnothing 2$ | DB | : OR |
| A640 |  | SYN | CNE, : OR |
| A640 | +1E | DB | CNE |
| A641 | $+\varnothing 2$ | DB | : OR |
| A642 |  | SYN | CLT, : OR |
| A642 | $+2 \emptyset$ | DB | CLT |
| A643 | $+\varnothing 2$ | DB | : OR |
| A644 |  | SYN | CGT, : OR |
| A644 | +21 | DB | CGT |
| A645 | $+\emptyset 2$ | DB | : OR |
| A646 |  | SYN | CEQ, : OR |
| A646 | +22 | DB | CEQ |
| A647 | $+\varnothing 2$ | DB | : OR |
| A648 |  | SYN | CAND, : OR |
| A648 | $+2 \mathrm{~A}$ | DB | CAND |
| A649 | $+\square 2$ | DB | : OR |
| A64A |  | SYN | COR |
| A64A | $+29$ | DB | COR |
| A64B |  | SYN | : RTN |
| A64B | +ø3 | DB | : RTN |



Source Code


| A675 | $+\emptyset E$ | DB | : VEXP |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A676 |  | SYN | CRPRN |  |  |  |
| A676 | $+2 \mathrm{C}$ | DB | CRPRN |  |  |  |
| A677 |  | SYN | : RTN |  |  |  |
| A677 | +ø3 | DB | : RTN |  |  |  |
| $<$ SFP $>=<$ STR $>$ ) \# |  |  |  |  |  |  |
| A678 |  | : SFP SYN | CLPRN, : CHNG, CFL | PRN |  |  |
| A678 | +2B | DB | CLPRN |  |  |  |
| A679 | $+\emptyset F$ | DB | : CHNG |  |  |  |
| A67A | $+3 \mathrm{~A}$ | DB | CFLPRN |  |  |  |
| A67B |  | SYN | JS, : STR |  |  |  |
| A67B | +C7 | DB | \$8Ø+(( $:$ STR-*) \& 7 F ) | XOR | \$40 | ) |
| A67C |  | SYN | CRPRN |  |  |  |
| A67C | $+2 \mathrm{C}$ | DB | CRPRN |  |  |  |
| A67D |  | SYN | : RTN |  |  |  |
| A67D | +ø3 | DB | : RTN |  |  |  |
| $<\mathrm{STCOMP}>=<\mathrm{STR}><\mathrm{SOP}><\mathrm{STR}>\#$ |  |  |  |  |  |  |
| A67E |  | : STCOMP SYN | JS, : STR |  |  |  |
| A67E | +C4 | DB | \$80+( ( $:$ STR-*) \& ${ }^{\text {S }} 7 \mathrm{~F}$ ) | XOR | \$40 | ) |
| A67F |  | SYN | JS, : SOP |  |  |  |
| A67F | +E3 | DB | \$80+( ( (:SOP-*) \& ${ }^{\text {S }}$ (F) | XOR | \$40 | ) |
| A680 |  | SYN | JS, : STR |  |  |  |
| A680 | +C2 | DB | \$8Ø+( ( (:STR-*) \& ${ }^{\text {S }}$ (F) | XOR | \$4ø | ) |
| A681 |  | SYN | : RTN |  |  |  |
| A681 | +ø3 | DB | : RTN |  |  |  |
| $<$ STR $>=<$ SFUN $>\mid<$ SVAR $>\mid<S C O N>\#$ |  |  |  |  |  |  |
| A682 |  | : STR SYN | JS, : SFUN |  |  |  |
| A682 | +C8 | DB | \$8Ø+( ( $:$ SFUN-*) \& ${ }^{\text {P }} 7 \mathrm{~F}$ ) | XOR | \$4ø | ) |
| A683 |  | SYN | : OR |  |  |  |
| A683 | $+\emptyset 2$ | DB | : OR |  |  |  |
| A684 |  | SYN | JS, : SVAR |  |  |  |
| A684 | +CB | DB | \$8Ø+( ( $:$ SVAR-*) \& \$ 7F ) | XOR | \$4ø | ) |
| A685 |  | SYN | : OR |  |  |  |
| A685 | $+\emptyset 2$ | DB | : OR |  |  |  |
| A686 |  | SYN | : ESRT, AD, : TSCON-1 |  |  |  |
| A686 | $+\varnothing 1$ | DB | : ESRT |  |  |  |
| A687 | $+27 \mathrm{~A} 4$ | DW | (:TSCON-1) |  |  |  |
| A689 |  | SYN | : RTN |  |  |  |
| A689 | +ø3 | DB | : RTN |  |  |  |
| $<$ SFUN $>=$ SFNP < NFP $>$ \# |  |  |  |  |  |  |
| A68A |  | : SFUN SYN | : ANTV, AD, : SFNP- |  |  |  |
| A68A | $+\emptyset \emptyset$ | DB | : ANTV |  |  |  |
| A68B | +D5A7 | DW | (:SFNP-1) |  |  |  |
| A68D |  | SYN | JS, : NFP |  |  |  |
| A68D | +A5 | DB | \$8Ø+( ( : NFP-*)\&\$7F) | XOR | \$40 | ) |
| A68E |  | SYN | : RTN |  |  |  |
| A68E | $+\varnothing 3$ | DB | : RTN |  |  |  |
| $<$ SVAR $>=<$ TSVAR $><$ SMAT $>\#$ |  |  |  |  |  |  |
| A68F |  | : SVAR SYN | : ESRT, AD, :TSVAR | R-1 |  |  |
| A68F | $+\varnothing 1$ | DB | : ESRT |  |  |  |
| A690 | +2DA3 | DW | (: TSVAR-1) |  |  |  |
| A692 |  | SYN | JS, : SMAT |  |  |  |
| A692 | +C2 | DB | \$80+( ( $:$ SMAT-*) \& ${ }^{\text {P }} 7 \mathrm{~F}$ ) | ) XOR | \$4ø | ) |
| A693 |  | SYN | : RTN |  |  |  |
| A693 | +ø3 | DB | : RTN |  |  |  |
| $<\mathrm{SMAT}>=(<$ EXP $><$ SMAT2 $>) \mid \& \#$ |  |  |  |  |  |  |
| A694 |  | ; SMAT SYN | CLPRN, : CHNG, CSLPRN |  |  |  |
| A694 | $+2 \mathrm{~B}$ | DB | CLPRN |  |  |  |
| A695 | $+\emptyset F$ | DB | : CHNG |  |  |  |
| A696 | $+37$ | DB | CSLPRN |  |  |  |

## Source Code

| A697 |  | SYN | : VExP |  |
| :---: | :---: | :---: | :---: | :---: |
| A697 | $+\varnothing \mathrm{E}$ | DB | :VEXP |  |
| A698 |  | SYN | JS, : SMAT2 |  |
| A698 | +C4 | DB | \$8ø+( ( $:$ SMAT2-*) $\&$ \$ 7 F$)$ | XOR \$4ø |
| A699 |  | SYN | CRPRN |  |
| A699 | +2C | DB | CRPRN |  |
| A69A |  | SYN | : OR |  |
| A69A | $+\varnothing 2$ | DB | : OR |  |
| A69B |  | SYN | : RTN |  |
| A69B | +ø3 | DB | : RTN |  |
| $<$ SMAT2> $=$, <EXP $>\mid \& \#$ |  |  |  |  |
| A69C |  | :SMAT2 SYN | CCOM, : CHNG, CACOM |  |
| A69C | +12 | DB | ССом |  |
| A69D | $+\emptyset \mathrm{F}$ | DB | : CHNG |  |
| A69E | $+3 \mathrm{C}$ | DB | CACOM |  |
| A69F |  | SYN | : VEXP |  |
| A69F | $+\emptyset E$ | DB | : VEXP |  |
| A6Aø |  | SYN | : OR |  |
| А6Аø | +ø2 | DB | : OR |  |
| A6Al |  | SYN | : RTN |  |
| A6Al | +ø3 | DB | : RTN |  |
| $<$ SOP $>=<><$ \# |  |  |  |  |
| A6A2 |  | : SOP |  |  |
| A6A2 |  | SYN | CLE, : CHNG, CSLE, : OR |  |
| A6A2 | +1D | DB | CLE |  |
| A6A3 | $+\emptyset \mathrm{F}$ | DB | : CHNG |  |
| A6A4 | +2F | DB | CSLE |  |
| A6A5 | +ø2 | DB | : OR |  |
| A6A6 |  | SYN | CNE, : CHNG, CSNE, : OR |  |
| A6A6 | +1E | DB | CNE |  |
| A6A7 | $+\varnothing \mathrm{F}$ | DB | : CHNG |  |
| A6A8 | $+3 \varnothing$ | DB | CSNE |  |
| A6A9 | $+\varnothing 2$ | DB | : OR |  |
| A6AA |  | SYN | CGE, : CHNG, CSGE, : OR |  |
| A6AA | +1F | DB | CGE |  |
| A6AB | $+\emptyset F$ | DB | : CHNG |  |
| A6AC | +31 | DB | CSGE |  |
| A6AD | $+\emptyset 2$ | DB | : OR |  |
| A6AE |  | SYN | CLT, : CHNG, CSLT, : OR |  |
| A6AE | $+2 \emptyset$ | DB | CLT |  |
| A6AF | $+\square \mathrm{F}$ | DB | : CHNG |  |
| A6Bø | +32 | DB | CSLT |  |
| A6Bl | +ø2 | DB | : OR |  |
| A6B2 |  | SYN | CGT, : CHNG, CSGT, : OR |  |
| A6B2 | +21 | DB | CGT |  |
| A6B3 | $+\square \mathrm{F}$ | DB | : CHNG |  |
| A6B4 | +33 | DB | CSGT |  |
| A6B5 | +ø2 | DB | : OR |  |
| A6B6 |  | SYN | CEQ, : CHNG, CSEQ |  |
| A6B6 | +22 | DB | CEQ |  |
| A6B7 | $+\emptyset \mathrm{F}$ | DB | : CHNG |  |
| A6B8 | +34 | DB | CSEQ |  |
| A6B9 |  | SYN | :RTN |  |
| A6B9 | $+\varnothing 3$ | DB | : RTN |  |
| <PUT $>=<$ D1 $>$, <EXP $><$ EOS $>\#$ |  |  |  |  |
| A6BA |  | : SPUT |  |  |
| A6BA |  | SYN | CPND, : VEXP |  |
| A6BA | +1C | DB | CPND |  |
| A6BB | $+\emptyset \mathrm{E}$ | DB | : VEXP |  |
| A6BC |  | SYN | ССОМ |  |
| A6BC | +12 | DB | ССом |  |

```
< > =<EXP><EOS > #
llol
< > =<EOS >#
A6BE :SCSAVE
A6BE :SDOS
A6BE :SCLR
A6BE :SRET
A6BE :SEND
A6BE :SSTOP
A6BE :SPOP
A6BE :SNEW
A6BE :SBYE
A6BE :SCONT
A6BE :SDEG
A6BE :SRAD SYN JS,:EOS
```



```
A6BF +ø3 DB :RTN
<LET> = <NVAR > = <EXP><EOS> | <SVAR > = <STR>< <EOS>#
A6CØ 
\@CO
A6CØ +\emptyset\emptyset
A6Cl +4BA6
SYN :ANTV,AD,:NVAR-1
                    DB :ANTV
                DW (:NVAR-1)
        SYN CEQ,:CHNG,CAASN
A6C3 SYN CE CEQ
A6C4 +ØF DB : CHNG
A6C5 +2D DB CAASN
A6C6 SYN :VEXP
A6C6 +ØE DB :VEXP
A6C7 SYN JS,:EOS
A6C7 +F1 
A6C8 +Ø2 DB :OR
A6C9 ; SYN JS,:SVAR
A6C9 +86 DB $8\emptyset+(((:SVAR-*)&$7F) XOR $4\emptyset )
A6CA SYN CEQ,:CHNG,CSASN
A6CA +22 DB CEQ
A6CB +\emptysetF DB :CHNG
A6CC +2E DB CSASN
A6CD SYN :ANTV,AD,:STR-1
A6CD +ø\emptyset DB :ANTV
A6CE +81A6 DW (:STR-1)
A6DØ SYN JS,:EOS
A6D\emptyset +E8 DB $ $8\emptyset+(((:EOS-*)&$7F) XOR $4\emptyset)
A6D1 SYN :RTN
A6D1 +ø3 DB :RTN
<FOR> = <TNVAR > = <EXP> TO <EXP><FSTEP><EOS >#
A6D2 :SFOR SYN :ESRT,AD,:TNVAR-1
A6D2 +Ø1 
A6D3 +29A3 (%)
A6D6 +ØF
DB CEQ
DB : CHNG
A6D7 +2D DB CAASN
```


## Source Code



| $<$ LOCATE $>=<$ EXP $>,<$ EXP $>,<$ TNVAR $><$ EOL $>\#$ |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| A6E 2 |  | : SLOCATE |  |  |  |
| A6E2 |  | SYN | : VEXP |  |  |
| A6E2 | $+\emptyset E$ | DB | : VEXP |  |  |
| A6E3 |  | SYN | CCOM |  |  |
| A6E3 | +12 | DB | CCOM |  |  |
| A6E4 |  | SYN | : VEXP |  |  |
| A6E4 | $+\emptyset E$ | DB | : VEXP |  |  |
| A6E5 |  | SYN | CCOM |  |  |
| A6E5 | +12 | DB | CCOM |  |  |
| A6E6 |  | SYN | JS, : SNEXT |  |  |
| A6E6 | +C4 | DB | \$8Ø+( ( (:SNEXT-*) \& \$ 7F ) | XOR | \$4Ø |
| A6E7 |  | SYN | : RTN |  |  |
| A6E7 | $+\varnothing 3$ | DB | : RTN |  |  |


$<$ NEXT $>=<$ TNVAR $><$ EOS $>\#$



## Source Code



| A72A |  | SYN | JS, : EOS |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A72A | $+8 \mathrm{E}$ | DB | \$8ø+( ( $:$ EOS-*) \& ${ }^{\text {P }} 7 \mathrm{~F}$ ) | XOR | \$4ø | ) |
| A72B |  | SYN | :RTN |  |  |  |
| A72B | +ø3 | DB | : RTN |  |  |  |
| $<\mathrm{OPD}>=<\mathrm{D} 1>$, \| \# |  |  |  |  |  |  |
| A72C |  | : OPD |  |  |  |  |
| A72C |  | SYN | JS, : Dl |  |  |  |
| A72C | +99 | DB | \$8Ø+(():Dl-*)\&\$7F) X | XOR \$ | \$40) |  |
| A72D |  | : OPDX SYN | CCOM |  |  |  |
| A72D | +12 | DB |  |  |  |  |
| A72E |  | SYN | : OR |  |  |  |
| A72E | $+\varnothing 2$ | DB | : OR |  |  |  |
| A72F |  | SYN | JS, : D1 |  |  |  |
| A72F | +96 | DB | $\underset{\text { CSC }}{\text { \$8®+( }}$ ( $(:$ Dl-*) $¢ \$ 7 \mathrm{~F})$ XOR \$4Ø) |  |  |  |
| A730 |  | SYN |  |  |  |  |
| A730 | +15 | DB | CsC |  |  |  |
| A731 |  | SYN | : OR |  |  |  |
| A731 | $+\varnothing 2$ | DB ${ }^{\text {- }}$ | : OR |  |  |  |
| A732 |  | SYN | : RTN |  |  |  |
| A732 | +ø3 | DB | :RTN |  |  |  |
| $<$ LIST $>=<$ FS $>$; <L2 $>\mid<$ L2 $>$ \# |  |  |  |  |  |  |
| A733 |  | : SLIST |  |  |  |  |
| A733 |  | SYN | JS,:FS |  |  |  |
| A733 | +DE | DB | \$80+(():FS-*)\&\$7F) XOR \$4ø |  |  |  |
| A734 |  | SYN | JS, : EOS |  |  |  |
| A734 | +84 | DB | \$8Ø+(():EOS-*)\&\$7F) XOR \$4ø ) |  |  |  |
| A735 |  | SYN | : OR |  |  |  |
| A735 | +ø2 | DB | : OR ${ }^{\text {JS, }}$, FS |  |  |  |
| A736 |  | SYN | JS, :FS |  |  |  |
| A736 | +DB | DB | \$8ø+(():FS-*)\&\$7F) XOR \$4ø ) |  |  |  |
| A737 |  | SYN | CCOM |  |  |  |
| A737 | +12 | DB | ССОМ |  |  |  |
| A738 |  | SYN | JS, : LIS |  |  |  |
| A738 | +C4 | DB | $\$ 8 \varnothing+(((: L I S-*) \& \$ 7 F)$ XOR $\$ 4 \emptyset)$ : OR |  |  |  |
| A739 |  | SYN |  |  |  |  |
| A739 | $+\varnothing 2$ | DB | : OR |  |  |  |
| A73A |  | SYN |  |  |  |  |
| A73A | +C2 | DB |  |  |  |  |
| A73B |  | SYN | : RTN |  |  |  |
| A73B | +ø3 | DB | : RTN |  |  |  |
| $<$ LIS $>=<$ L1 $><$ EOS2 $>$ \# |  |  |  |  |  |  |
| A73C |  | :LIS |  |  |  |  |
| A73C |  | SYN | : ANTV, AD, : L1-1 |  |  |  |
| A73C | $+\varnothing \varnothing$ | DB | : ANTV$(: \mathrm{Ll}-1)$ |  |  |  |
| A73D | +BFA 7 | DW |  |  |  |  |
| A73F |  | SYN | JS, : EOS2 |  |  |  |
| A73F | +F4 | DB | $\$ 8 \varnothing+((): \operatorname{RTN}$ <br> $: \operatorname{RTN} 2-*) \& \$ 7 F)$ XOR $\$ 4 \varnothing$ |  |  |  |
| A74ø |  | SYN |  |  |  |  |
| A740 | $+\varnothing 3$ | DB | :RTN |  |  |  |
| $<$ STATUS $>=<$ STAT $><$ EOS2 $>$ \# |  |  |  |  |  |  |
| A741 |  | : SSTATUS | JS, : STAT |  |  |  |
| A741 |  | SYN |  |  |  |  |
| A741 | +C3 | DB | \$8ø+(():STAT-*)\&\$7F) XOR \$40 |  |  |  |
| A742 |  | SYN | JS, : EOS2 |  |  |  |
| A742 | +F1 | DB |  |  |  | ) |
| A743 |  | SYN |  |  |  |  |
| A743 | $+\square 3$ | DB | : RTN |  |  |  |
| $<$ STAT $>=<$ D1 $>$, <NVAR $>$ \# |  |  |  |  |  |  |
| A744 |  | : STAT |  |  |  |  |
| A744 |  | SYN | JS, ${ }^{\text {Dl }}$ |  |  |  |
| A744 | +81 | DB | \$8ø+(():D1-*)\&\$7F) X | XOR \$ | \$40 ) |  |

Source Code

| A745 |  | SYN | ССОм |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A745 | +12 | DB | CCOM |  |  |  |
| A746 |  | SYN | : ANTV, AD, : NVAR-1 |  |  |  |
| A746 | +øø | DB | : ANTV |  |  |  |
| A747 | +4BA6 | DW | (:NVAR-1) |  |  |  |
| A749 |  | SYN | : RTN |  |  |  |
| A749 | $+\varnothing 3$ | DB | : RTN |  |  |  |
| $<>=<$ STAT $>,<$ NVAR $><$ EOS2 $>\#$ |  |  |  |  |  |  |
| A74A |  | : SNOTE |  |  |  |  |
| A74A |  | :SPOINT |  |  |  |  |
| A74A |  | SYN | JS, : STAT |  |  |  |
| A74A | +BA | DB | \$8Ø+( ( : STAT-*) \&\$7F) | XOR | \$40 |  |
| A74B |  | SYN | CCOM |  |  |  |
| A74B | +12 | DB | ССОМ |  |  |  |
| A74C |  | SYN | : ANTV, AD, : NVAR-1 |  |  |  |
| A74C | $+\varnothing \varnothing$ | DB | : AnTV |  |  |  |
| A74D | +4BA6 | DW | (:NVAR-1) |  |  |  |
| A74F |  | SYN | JS, : EOS2 |  |  |  |
| A74F | +E4 | DB | \$8Ø+(():EOS2-*)\&\$7F) | XOR | \$40 |  |
| A750 |  | SYN | : RTN |  |  |  |
| A750 | $+\varnothing 3$ | DB | : RTN |  |  |  |
| $<\mathrm{FS}>=<$ STR $>$ |  |  |  |  |  |  |
| A751 |  | : FS |  |  |  |  |
| A751 |  | SYN | : ANTV, AD, : STR-1 |  |  |  |
| A751 | +øø | DB | : ANTV |  |  |  |
| A752 | +81A6 | DW | (:STR-1) |  |  |  |
| A754 |  | SYN | : RTN |  |  |  |
| A754 | $+\varnothing 3$ | DB | : RTN |  |  |  |
| $<$ TEXP $>=<$ EXP $>,<$ EXP $>\#$ |  |  |  |  |  |  |
|  |  | ; |  |  |  |  |
| A755 |  | :TEXP |  |  |  |  |
| A755 |  | SYN | : VEXP |  |  |  |
| A755 | $+\varnothing \mathrm{E}$ | DB | : VEXP |  |  |  |
| A756 |  | SYN | CCOM |  |  |  |
| A756 | +12 | DB | CCOM |  |  |  |
| A757 |  | SYN | : VEXP |  |  |  |
| A757 | $+\emptyset E$ | DB | : VEXP |  |  |  |
| A758 |  | SYN | :RTN |  |  |  |
| A758 | $+03$ | DB | : RTN |  |  |  |
| $<$ SOUND $>=<$ EXP $>,<$ EXP $>,<$ EXP $>,<$ EXP $><$ EOS $>\#$ |  |  |  |  |  |  |
| A759 |  | :SSOUND |  |  |  |  |
| A759 |  | SYN | : VEXP |  |  |  |
| A759 | $+\varnothing$ E | DB | : VEXP |  |  |  |
| A75A |  | SYN | ССОМ |  |  |  |
| A75A | +12 | DB | CCOM |  |  |  |
| A75B |  | : SSETCOLOR |  |  |  |  |
| A75B |  | SYN | : VEXP |  |  |  |
| A75B | $+\varnothing$ E | DB | : VEXP |  |  |  |
| A75C |  | SYN | CCOM |  |  |  |
| A75C | +12 | DB | CCOM |  |  |  |
| $<>=<E X P>,<E X P><E O S>\#$ |  |  |  |  |  |  |
| A75D |  | : SPOKE |  |  |  |  |
| A75D |  | : SPLOT |  |  |  |  |
| A75D |  |  |  |  |  |  |
| A75D |  | : SDRAWTO |  |  |  |  |
| A75D |  | SYN | JS, : TEXP |  |  |  |
| A75D | +B8 | DB | \$8ø+( ( (:TEXP-*) $¢ \$ 7 \mathrm{~F})$ | XOR | \$4ø | ) |
| A75E |  | SYN | JS, : EOS2 |  |  |  |
| A75E | +D5 | DB | \$8ø+( ( $(: E O S 2-*) \& \$ 7 F)$ | XOR | \$40 |  |
| A75F |  | SYN | :RTN |  |  |  |
| A 75 F | +ø3 | DB | : RTN |  |  |  |



## Source Code

| A778 | +29A3 | DW | (:TNVAR-1) |
| :---: | :---: | :---: | :---: |
| A77A |  | SYN | CLPRN, : CHNG, CDLPRN |
| A77A | $+2 \mathrm{~B}$ | DB | CLPRN |
| A77B | $+\emptyset \mathrm{F}$ | DB | : CHNG |
| A 77 C | +39 | DB | CDLPRN |
| A77D |  | SYN | : VEXP |
| A77D | $+\emptyset E$ | DB | : VEXP |
| A77E |  | SYN | : ANTV, AD, : NMAT2-1 |
| A77E | $+\emptyset \emptyset$ | DB | : ANTV |
| A 77 F | +58A6 | DW | (: NMAT2-1) |
| A781 |  | SYN | CRPRN |
| A781 | $+2 \mathrm{C}$ | DB | CRPRN |
| A782 |  | SYN | : OR |
| A782 | +62 | DB | : OR |
| A783 |  | SYN | : ESRT, AD, : TSVAR-1 |
| A 783 | $+\varnothing 1$ | DB | : ESRT |
| A784 | +2 DA3 | DW | (:TSVAR-1) |
| A786 |  | SYN | CLPRN, : CHNG, CDSLPR |
| A786 | +2B | DB | CLPRN |
| A 787 | $+\emptyset F$ | DB | : CHNG |
| A.788 | +3B | DB | CDSLPR |
| A789 |  | SYN | : VEXP |
| A789 | $+\emptyset E$ | DB | : VEXP |
| A78A |  | SYN | CRPRN |
| A78A | $+2 \mathrm{C}$ | DB | CRPRN |
| A78B |  | SYN | : RTN |
| A.78B | $+\varnothing 3$ | DB | : RTN |
| $<$ NSML $>=<$ NSMAT $><$ NSML2> $\mid \& \#$ |  |  |  |
| A78C |  | :NSML SYN | JS, : NSMAT |
| A78C | $+\mathrm{AB}$ | DB | \$8Ø+(():NSMAT-*)\&\$7F) XOR \$40) |
| A78D |  | SYN | JS, : NSML2 |
| A78D | +C3 | DB |  |
| A 78 E |  | SYN | :OR, : RTN |
| A 78 E | $+\emptyset 2$ | DB | : OR |
| A78F | +ø3 | DB | : RTN |
| <NSML2> $=,<$ NSML> $\mid$ \&\# |  |  |  |
| A790 |  | : NSML2 SYN | CCOM |
| A790 | +12 | DB | CCOM |
| A791 |  | SYN | JS, : NSML |
| A791 | +BB | DB |  |
| A792 |  | SYN | : OR, : RTN |
| A 792 | $+\varnothing 2$ | DB | : OR |
| A793 | $+\varnothing 3$ | DB | : RTN |
| $<\boldsymbol{\\| F}>=<$ EXP $>$ THEN $<$ IFA $><$ EOS $>\#$ |  |  |  |
| A794 |  | :SIF SYN | : VEXP |
| A794 | $+\emptyset E$ | DB | :VEXP |
| A795 |  | SYN | CTHEN |
| A795 | +1B | DB | CTHEN |
| A796 |  | SYN | JS, : IFA |
| A796 | $+\mathrm{C} 3$ | DB |  |
| A797 |  | SYN | JS, : EOS 2 |
| A797 | $+9 \mathrm{C}$ | DB | \$80+( ( $:$ EOS $2-*) \& \$ 7 \mathrm{~F})$ XOR \$40) |
| A798 |  | SYN | : RTN |
| A798 | +ø3 | DB | : RTN |
| $<\mathrm{IFA}>=<$ TNCON $>1<$ EIF $>$ |  |  |  |
| A799 |  | :IFA SYN | : ESRT, AD, :TNCON-1 |
| A799 | $+\emptyset 1$ | DB | : ESRT |
| A79A | +FFA3 | DW | (:TNCON-1) |
| A79C |  | SYN | : OR |
| A 79 C | $+\varnothing 2$ | DB | : OR |
| A79D |  | SYN | : ESRT, AD, : EIF-1 |
| A79D | $+\varnothing 1$ | DB | : ESRT |
| A79E | +D3A2 | DW | (:EIF-1) |

## Source Code

| $<\mathbf{P R 1}$＞$=$＜PEL＞｜ | ＜PSL $>$＜PR2＞ | ＞\＆\＃ |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| A7Aの | ：PR1 |  |  |  |  |
| A7Aの | SYN | JS，：PEL，OR |  |  |  |
| A7Aロ＋C9 | DB \＄ | \＄8Ø＋（（ $(:$ PEL－＊）$\%$ \＄ 7 F$)$ | XOR | \＄4ø | ） |
| A7Al $+\square 2$ | DB ： | ：OR |  |  |  |
| A7A2 | SYN | JS，：PSL |  |  |  |
| A7A2＋D4 | DB \＄ | \＄8Ø＋（（）：PSL－＊）\＆\＄7F） | XOR | \＄4ø | ） |
| A7A3 | SYN | JS，：PR2 |  |  |  |
| A7A3 +C 3 | DB \＄ | \＄8Ø＋（（（：PR2－＊）\＆\＄7F） | XOR | \＄40 | ） |
| A7A4 | SYN | ：OR |  |  |  |
| A7A $4+\emptyset 2$ | DB ： | ：OR |  |  |  |
| A7A5 | SYN | ：RTN |  |  |  |
| A7A5＋ 63 | DB ： | ：RTN |  |  |  |
|  | ； |  |  |  |  |
| $<\mathbf{P R 2}>=<$ PEL $>\mid \& \#$ |  |  |  |  |  |
| A7A6 | ：PR2 SYN | JS，PEL |  |  |  |
| A7A6＋C3 | DB | \＄80＋（（ $(:$ PEL－＊） 8 \＄ 7 F$)$ | XOR | \＄40 | ） |
| A7A7 | SYN | ：OR |  |  |  |
| A7A7 $+\emptyset 2$ | DB ： | ：OR |  |  |  |
| A7A8 | SYN | ：RTN |  |  |  |
| A7A8 $+\varnothing 3$ | DB ： | ：RTN |  |  |  |
| ＜PEL＞$=<$ PES $><$ PELA $>$ \＃ |  |  |  |  |  |
| A7A9 | ：PEL SYN | JS，PES |  |  |  |
| A7A9＋C3 | DB \＄ | \＄8ø＋（（ $(:$ PES $-*) \& \$ 7 \mathrm{~F})$ | XOR | \＄4ø | ） |
| A7AA | SYN | Js，：PELA |  |  |  |
| A7AA +C 8 | DB \＄ | \＄8ø＋（（ $(:$ PELA－＊）\＆$\$ 7 \mathrm{~F})$ | ）XOR | R \＄4ø | ） |
| A7AB | SYN | ：RTN |  |  |  |
| $A 7 A B+\varnothing 3$ | DB ： | ：RTN |  |  |  |
| $<$ PES＞$=$＜EXP $>$｜＜STR $>$ |  |  |  |  |  |
| A7AC | ：PES SYN | ：VEXP |  |  |  |
| A7AC $+\emptyset E$ | DB | ：VEXP |  |  |  |
| A7AD | SYN | ：OR |  |  |  |
| A7AD $+\emptyset 2$ | DB | ：OR |  |  |  |
| A7AE | SYN | ：ANTV，AD，：STR－1 |  |  |  |
| A7AE＋øø | DB | ：ANTV |  |  |  |
| A7AF＋81A6 | DW | （：STR－1） |  |  |  |
| A7B1 | SYN | ：RTN |  |  |  |
| A7B1 $+\square 3$ | DB | ：RTN |  |  |  |
| ＜PELA $>=<$ PSL $><$ PEL $>\mid$ \＆\＃ |  |  |  |  |  |
| A7B2 | ：PELA SYN | JS，：PSL |  |  |  |
| A7B2 +C 4 | DB \＄ | \＄8ø＋（（ $:$ PSL－＊）\＆${ }^{\text {P }} 7 \mathrm{~F}$ ） | XOR | \＄4ø | ） |
| A7B3 | SYN | JS，：PR2 |  |  |  |
| A7B3 +B 3 | DB \＄ | \＄8ø＋（（）：PR2－＊）\＆\＄7F） | XOR | \＄4ø | ） |
| A7B4 | SYN | ：OR |  |  |  |
| A7B4＋ 62 | DB | ：OR |  |  |  |
| A7B5 | SYN | ：RTN |  |  |  |
| A7B5 + Ø3 | DB | ：RTN |  |  |  |
| $<\mathbf{P S L}>=<\mathbf{P S}><\mathbf{P S L A}>\#$ |  |  |  |  |  |
| A7B6 | ：PSL SYN | JS，：PS |  |  |  |
| A7B6＋C6 | DB \＄ | \＄80＋（（）：PS－＊）\＆\＄7F）X | XOR | \＄40） |  |
| A7B7 | SYN | JS，：PSLA |  |  |  |
| А $7 \mathrm{~B} 7+\mathrm{C} 2$ | DB | \＄8ø＋（（）：PSLA－）\＆\＄7F） | ）XOR | \＄4ø | ） |
| A7B8 | SYN | ：RTN |  |  |  |
| A7B8 $+\varnothing 3$ | DB ： | ：RTN |  |  |  |
| ＜PSLA $>=<\mathbf{P S L}>1$ \＆\＃ |  |  |  |  |  |
| A7B9 | ：PSLA SYN | JS，：PSL |  |  |  |
| A7B9＋BD | DB \＄ | \＄80＋（（ $:$ PSL－＊）$¢ \$ 7 \mathrm{~F})$ | XOR | \＄4ø | ） |
| A7BA | SYN | ：OR |  |  |  |

## Source Code



```
\(<\) SFNP \(>=\) STR \(\mid\) CHR\#
```



```
\(<\) PUSR \(>=<\) EXP \(><\) PUSR1 \(>\) \#
\begin{tabular}{|c|c|c|c|c|}
\hline A7DA & & :PUSR SYN & : VEXP & \\
\hline A7DA & \(+\varnothing \mathrm{E}\) & DB & : VEXP & \\
\hline A7DB & & SYN & JS, : PUSR1 & \\
\hline A7DB & +C2 & DB & \$8ø+(():PUSR1-*)\&\$7F) & XOR \$4ø \\
\hline A7DC & & SYN & :RTN & \\
\hline A7DC & \(+ø 3\) & DB & :RTN & \\
\hline
\end{tabular}
\(<\) PUSR1 \(>=,<\) PUSR \(>\mid \& \#\)
```



```
            OPNTAB - Operator Name Table
\begin{tabular}{|c|c|c|c|c|c|}
\hline \multirow[t]{4}{*}{A7E3} & & opntab & & & \\
\hline & \(=\varnothing \varnothing \emptyset F\) & & SET & \$ \(\square_{\mathrm{F}}\) & \multirow[t]{3}{*}{;FIRST ENTRY VALUE=\$10} \\
\hline & = 9010 & & & & \\
\hline & = øø1ø & CDQ & \[
\begin{aligned}
& \text { SET } \\
& \text { EQU }
\end{aligned}
\] & \({ }_{\text {c }}+\) & \\
\hline \multirow[t]{4}{*}{A7E3} & 82 & DB & \$82 & & \multirow[t]{4}{*}{; DOUBLE QUOTE} \\
\hline & \(=9011\) & ; & & & \\
\hline & = & C & SET & c+1 & \\
\hline & = 9011 & CSOE & EQU & C & \\
\hline \multirow[t]{3}{*}{A 7 E 4} & \(8 \varnothing\) & DB & \$8ø & & \multirow[t]{3}{*}{; DUMMY FOR SOE} \\
\hline & = 0012 & ' & SET & C+1 & \\
\hline & \(=\varnothing 012\) & ССом & EQU & C & \\
\hline \multirow[t]{4}{*}{A7E5} & AC & DC & ', ' & & \\
\hline & & ; & & & \\
\hline & = \(\square 013\) & C & SET & C+1 & \\
\hline & = 9013 & CDOL & EQU & C & \\
\hline \multirow[t]{3}{*}{A7E6} & A4 & DC & ' \({ }^{\prime}\) & & \\
\hline & = øø14 & ; & SET & C+1 & \\
\hline & = ø014 & ceos & EQU & c & \\
\hline \multirow[t]{3}{*}{A7E7} & BA & DC & ': ' & & \\
\hline & \(=\varnothing 015\) & ; & SET & C+1 & \\
\hline & \(=\varnothing 015\) & csc & EQU & C & \\
\hline \multirow[t]{3}{*}{A7E8} & BB & DC & ';' & & \\
\hline & = 0016 & ' \({ }^{\text {c }}\) & SET & C+1 & \\
\hline & = 0016 & CCR & EQU & c & ; CARRIAGE RETURN \\
\hline \multirow[t]{3}{*}{A7E9} & 9B & DB & CR & & \\
\hline & = Ø017 & ; & SET & C+1 & \\
\hline & \(=\varnothing 017\) & CGTO & EQU & c & \\
\hline A7EA & 474F54CF & DC & 'GOTO ' & & \\
\hline
\end{tabular}
```

Source Code



Source Code

| A825 | $8 \varnothing$ | DB | \$80 |  | ; DOES NOT PRINT |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | ; |  |  |  |  |  |  |
|  | $=003 \mathrm{~A}$ | C | SET | C+1 |  |  |  |  |
|  | $=\varnothing \emptyset 3 \mathrm{~A}$ | CFLPRN | EQU | C | ; | FUNCTION | LEF'T | PAREN |
| A826 | A8 | DC | ' (' |  | , |  |  | PAREN |
|  |  | ; |  |  |  |  |  |  |
|  | $=\varnothing 口 3 \mathrm{~B}$ | C | SET | C+1 |  |  |  |  |
|  | $=\varnothing 03 \mathrm{~B}$ | CDSLPR | EQU | C |  |  |  |  |
| A827 | A8 | DC | ' ${ }^{\prime}$ |  |  |  |  |  |
|  |  | ; |  |  |  |  |  |  |
|  | $=0.3 \mathrm{C}$ | C | SET | $C+1$ |  |  |  |  |
|  | $=0.35$ | CACOM | EQU | C | ; | ARRAY COM | MA |  |
| A828 | AC | DC | ',' |  | , | ARRAY COM |  |  |

Function Name Table


## Memory Manager

A87F

EXPAND


## Source Code

| A8C6 | B5øø |  | LDA |  | $\emptyset, \mathrm{x}$ |  | ADD ECSIZE TO |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A8C8 | 65A4 |  | ADC |  | ECSIZE | ; | all table entries |
| A8CA | 95øø |  | STA |  | $\emptyset, \mathrm{x}$ |  | FROM EXPAND AT ADR |
| A8CC | B501 |  | LDA |  | $1, \mathrm{x}$ |  | TO HIMEM |
| A8CE | 65A5 |  | ADC |  | ECSIZE+1 |  |  |
| A8DØ | 9501 |  | STA |  | $1, \mathrm{x}$ |  |  |
| A8D2 | E8 |  | INX |  |  |  |  |
| A8D3 | E8 |  | INX |  |  |  |  |
| A8D4 | Eø92 |  | CPX |  | \#MEMTOP+2 |  |  |
| A8D6 | 9øEE | ^A8C6 | BCC |  | : EXP 3 |  |  |
| A8D8 | 850F |  | STA |  | APHM +1 | ; | SET NEW APL |
| A8DA | A590 |  | LDA |  | MEMTOP | ; | HI MEM TO |
| A8DC | 850E |  | STA |  | APHM | ; | MEMTOP |
|  |  |  | ; |  |  |  |  |
| A8DE | A6A3 |  | LDX |  | MVLNG+1 | ; | $\mathrm{X}=$ MVLNG[H] |
| A8EØ | E8 |  | INX |  |  | ; | PLUS ONE |
| A8E1 | A4A2 |  | LDY |  | MVLNG | ; | Y = MVLNG[L] |
| A8E3 | DøøB | * A8Fø | BNE |  | : EXP6 | ; | TEST ZERO LENGTH |
| A8E5 | Fø1ø | *A8F7 | BEQ |  | : EXP7 | ; | BR IF LOW = Ø |
| A8E7 | 88 |  | ; EXP4 | DEY |  | ; | DEC MVLNG[L] |
| A8E8 | C69A |  | DEC |  | MVFA +1 |  | DEC MVFA[H] |
| A8EA | C69C |  | DEC |  | MVTA +1 | ; | dEC MVTA[H] |
| A8EC | B199 |  | ; EXP 5 | LDA | [MVFA], Y | ; | MVFA BYTE |
| A8EE | 919B |  | STA |  | [MVTA], Y | ; | TO MVTA |
| A8Fø | 88 |  | : EXP6 | DEY |  | ; | dec Count Low |
| A8Fl | DøF9 | ^A8EC | BNE |  | : EXP5 | ; | BR IF NOT ZERO |
| A8F3 | B199 |  | ; LDA |  | [MVFA], Y | ; | MOVE THE ZERO BYTE |
| A8F5 | 919B |  | STA |  | [MVTA], y |  |  |
|  |  |  |  |  |  |  |  |
| A8F7 |  |  | : EXP7 |  |  |  |  |
| A8F7 | CA |  | DEX |  |  |  | IF MVLNG[H] IS NOT |
| A8F8 | DøED | ^A8E7 | BNE |  | : EXP4 | ; | ZERO THEN MOVE 256 |
|  |  |  | ; |  |  |  | ELSE |
| A8FA | $6 \square$ |  | RTS |  |  |  | DONE |

CONTRACT


## Source Code



## Execute Control

A95F
LOCAL
EXECNL - Execute Next Line

|  | ; START PROGRAM EXECUTOR |  |  |
| :--- | :--- | :--- | :--- |
| A95F | i |  |  |
| A95F 201BB8 | EXECNL |  |  |
| JSR | SETLN1 | ; SET UP LIN \& NXT STMT |  |

## EXECNS - Execute Next Statement

| A962 |  | EXECNS |  |  |
| :---: | :---: | :---: | :---: | :---: |
| A962 | 20F4A9 | JSR | TSTBRK | ; TEST BREAK |
| A965 | D035 ^A99C | BNE | : EXBRK | ; BR IF BREAK |
| A967 | A4A7 | LDY | NXTSTD | ; GET PTR TO NEXT STMT L |
| A969 | C49F | CPY | LLNGTH | ;AT END OF LINE |
| A96B | Bø1C *A989 | BCS | : EXEOL | ; BR IF EOL |

Source Code

| A96D | B18A | LDA |  | [STMCUR], Y | ;GET NEW STMT LENGTH |
| :---: | :---: | :---: | :---: | :---: | :---: |
| A96F | 85A7 | STA |  | NXTSTD | ; SAVE AS FUTURE STMT Length |
| A971 | 98 | TYA |  |  | ; Y=DISPL TO THIS STMT LENGTH |
| A972 | C8 | INY |  |  | ;PLUS 1 IS DISPL TO CODE |
| A973 | B18A | LDA |  | [STMCUR], Y | ;GET CODE |
| A975 | C8 | INY |  |  | ; INC TO STMT MEAT |
| A976 | 84A8 | STY |  | STINDEX | ;SET WORK INDEX |
| A978 | 2Ø7EA9 | JSR |  | : STGO | ;GO EXECUTE |
| A97B | 4C62A9 | JMP |  | EXECNS | ;THEN DO NEXT STMT |
|  |  | ; |  |  |  |
| A97E |  | : STGO | ASLA |  | ;TOKEN*2 |
| A97E | $+\emptyset A$ | ASL |  | A |  |
| A97F | AA | TAX |  |  |  |
| A980 | BDøøAA | LDA |  | Stetab, X | ; GET ADR AND |
| A983 | 48 | PHA |  |  | ; PUSH TO STACK |
| A984 | BDø1AA | LDA |  | STETAB+1, X | ; AND GO TO |
| A987 | 48 | PHA |  |  | ;VIA |
| A988 | 60 | RTS |  |  | ;RTS |
|  |  | ; |  |  |  |
| A989 |  | : EXEOL |  |  |  |
| A989 | Aøø1 | LDY |  | \#1 |  |
| A98B | B18A | LDA |  | [STMCUR], Y |  |
| A98D | $3010{ }^{\text {^A99F }}$ | BMI |  | : EXFD | ; BR IF DIR |
|  |  | IDA |  |  |  |
| A98F | A59F | LDA |  | LLNGTH | ; GET LINE LENGTH |
| A991 | 2øDØA9 | JSR |  | GNXTL | ; INC STMCUR |
| A994 | 2øE2A9 | JSR |  | TENDST | ;TEST END STMT TABLE |
| A997 | 1øC6 ^A95F | BPL |  | EXECNL | ; BR NOT END |
| A999 | $4 \mathrm{C8DB} 7$ | ; EXDONE | JMP | XEND | ; GO BACK TO SYNTAX |
| A99C | $4 \mathrm{C93B7}$ | : EXBRK | JMP | XSTOP | ; BREAK, DO STOP |
| A99F | 4C5DAб | : EXFD | JMP | SNX3 | ; GO to syntax via ready msg |

## GETSTMT - Get Statement in Statement Table



| A9CA | 20DØA9 | JSR |  | GNXTL |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| A9CD | 4CB2A9 | JMP |  | : GS 2 |  |
|  |  | ; |  |  |  |
| A9DØ |  | GNXTL |  |  |  |
| A9Dø | 18 | CLC |  |  |  |
| A9D1 | 658A | ADC |  | STMCUR | ;ADD LENGTH TO STMCUR |
| A9D3 | 858A | STA |  | STMCUR |  |
| A9D 5 | A8 | TAY |  |  |  |
| A9D6 | A58B | LDA |  | STMCUR+1 |  |
| A9D8 | 690ø | ADC |  | \# $\emptyset$ |  |
| A9DA | 858B | STA |  | STMCUR+1 |  |
| A9DC | 60 | RTS |  |  |  |
| A9DD | Aøø2 | GETLL | LDY | \#2 |  |
| A9DF | B18A | LDA |  | [STMCUR], Y |  |
| A9El | 60 | RTS |  |  |  |

## TENDST - Test End of Statement Table

| A9E 2 |  |  | TENDST |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A9E2 | AøØ1 |  | LDY |  | \#1 |  | INDEX TO CNO ['I] |
| A9E4 | B18A |  | LDA |  | [STMCUR], Y | ; | GET CNO [HI] |
| A9E6 | $6 \square$ |  | RTS |  |  |  |  |
| A9E7 |  |  | XREM |  |  |  |  |
| A9E7 |  |  | XDATA |  |  |  |  |
| A9E7 | $6 \emptyset$ |  | TESTRTS | RTS |  |  |  |
| XBYE - Execute BYE |  |  |  |  |  |  |  |
| A9E8 |  |  | XBYE |  |  |  |  |
| A9E8 | 2041 BD |  | JSR |  | CLSALL | ; | CLOSE 1-7 |
| A9EB | 4C71E4 |  | JMP |  | BYELOC | ; | EXIT |
| XDOS - Execute DOS |  |  |  |  |  |  |  |
| A9EE |  |  | XDOS |  |  |  |  |
| A9EE | $2 \emptyset 41 \mathrm{BD}$ |  | JSR |  | CLSALL | ; | CLOSE 1-7 |
| A9Fl | 6СØAøØ |  | JMP |  | [DOSLOC] | ; | GO TO DOS |
| TSTBRK - Test for Break |  |  |  |  |  |  |  |
| A9F4 |  |  | TSTBRK |  |  |  |  |
| A9F4 | Аøøø |  | LDY |  | \#Ø |  |  |
| A9F6 | A511 | A9FE | LDA |  | BRKBY' | ; | LOAD BREAK BYTE |
| A9F8 | DØø4 * |  | BNE |  | : TB2 |  |  |
| A9FA | AØFF |  | LDY |  | \# \$FF |  |  |
| A9FC | 8411 |  | STY |  | BRKBYT |  |  |
| A9FE | 98 |  | : TB2 | TYA |  | ; | SET COND CODE |
| A9FF | $6 \emptyset$ |  | RTS |  |  | ; | DONE |

## Statement Execution Table

; STETAB-STATEMENT EXECUTION TABLE ; -CONTAINS STMT EXECUTION ADR : -MUST BE IN SAME ORDER AS SNTAB

AAøø
AADD
AAØØ
AAØØ
AAØ2
AAØ2
AAØ4
StETAB

| FDB | XREM-1 |
| :--- | :--- |
| DW | REV (XREM-1) |
| FDB | XDATA-1 |
| DW | REV (XDATA-1) |
| TA | EQU |
| FDB | (*-STETAB)/2-1 |
| DW | REV (XINPUT-1 |
| FDB | XCOLOR-1 |
| DW | REV (XCOLOR-1) |
| FDB | XLIST-1 |
| DW | REV (XLIST-1) |
| ST | EQU |
|  | $\quad$ (*-STETAB)/2-1 |

Source Code

| AAøA |  | FDB |  | XENTER-1 |
| :---: | :---: | :---: | :---: | :---: |
| AA®A | +BACA | DW |  | REV (XENTER-1) |
| ААФС |  | FDB |  | XLET-1 |
| ААøС | +AADF | DW |  | REV (XLET-1) |
| ААøE |  | FDB |  | XIF-1 |
| AADE | +B777 | DW |  | REV (XIF-1) |
| AAlø |  | FDB |  | XFOR-1 |
| AAID | $\begin{aligned} & +\mathrm{B} 64 \mathrm{~A} \\ & =ø ø \varnothing 8 \end{aligned}$ | $\mathrm{CFOR}^{\mathrm{DW}}$ | EQU | $\operatorname{REV} \underset{(*-S T E T A B)}{(\mathrm{XFOR}-1)} \mathbf{2 - 1}$ |
| AA12 |  | FDB |  | XNEXT-1 |
| AAl2 | +B6CE | DW |  | REV (XNEXT-1) |
| AA14 |  | FDB |  | XGOTO-1 |
| AAl4 | +B6A2 | DW |  | REV (XGOTO-1) |
| AAl6 |  | FDB |  | XGOTO-1 |
| AA16 | +B6A2 | DW |  | REV (XGOTO-1) |
| AA18 |  | FDB |  | XGOSUB-1 |
| AA18 | $\begin{aligned} & +\mathrm{B} 69 \mathrm{~F} \\ & =\emptyset \emptyset \varnothing \mathrm{C} \end{aligned}$ | $\begin{array}{r} \text { DW } \\ \text { CGOSUB } \end{array}$ | EQU | REV (XGOSUB-1)(*-STETAB) $/ 2-1$ |
| AA1A |  | FDB |  | XTRAP-1 |
| AAlA | +B7EØ | DW |  | REV (XTRAP-1) |
| AAlC |  | FDB |  | XBYE-1 |
| AAIC | +A9E7 | DW |  | REV (XBYE-1) |
| AAIE |  | FDB |  | XCONT-1 |
| AAIE | +B7BD | DW |  | REV ( $\mathrm{XCONT-1)}$ |
| AA2ø |  | FDB |  | XCOM-1 |
| AA2ø | +B1D8 | DW |  | REV (XCOM-1) |
| AA22 |  | FDB |  | XCLOSE-1 |
| AA22 | +BClA | DW |  | REV (XCLOSE-1) |
| AA24 |  | FDB |  | XCLR-1 |
| AA24 | +B765 | DW |  | REV (XCLR-1) |
| AA26 |  | FDB |  | XDEG-1 |
| AA26 | +B260 | DW |  | REV (XDEG-1) |
| AA28 |  | FDB |  | XDIM-1 |
| AA28 | +B1D8 | DW |  | REV (XDIM-1) |
| AA2A |  | FDB |  | XEND-1 |
| AA2A | +B78C | DW |  | REV (XEND-1) |
| AA2C |  | FDB |  | XNEW-1 |
| AA2C | +Aøбв | DW |  | REV (XNEW-1) |
| AA2E |  | FDB |  | XOPEN-1 |
| AA2E | +BBEA | DW |  | REV (XOPEN-1) |
| АА $3 \varnothing$ |  | FDB |  | XLOAD-1 |
| AA30 | +BAFA | DW |  | REV (XLOAD-1) |
| AA32 |  | FDB |  | XSAVE-1 |
| AA32 | +BB5C | DW |  | REV (XSAVE-1) |
| AA34 |  | FDB |  | XSTATUS-1 |
| AA 34 | +BC27 | DW |  | REV (XSTATUS-1) |
| AA36 |  | FDB |  | XNOTE-1 |
| AA36 | +BC35 | DW |  | REV (XNOTE-1) |
| AA38 |  | FDB |  | XPOINT-1 |
| АА 38 | +BC4C | DW |  | REV (XPOINT-1) |
| AA3A |  | FDB |  | XXIO-1 |
| AA3A | +BBE4 | DW |  | REV (XXIO-1) |
| AA3C |  | FDB |  | XON-1 |
| AA3C | +B7EC | DW |  | REV ( $\mathrm{XON-1}$ ) |
|  | $=\varnothing ø 1 \mathrm{E}$ | CON | EQU | (*-STETAB)/2-1 |
| AA3E |  | FDB |  | XPOKE-1 |
| AA3E | +B24B | DW |  | REV (XPOKE-1) |
| AA40 |  | FDB |  | XPRINT-1 |
| AA4ø | +B3B5 | DW |  | REV (XPRINT-1) |
| AA42 |  | FDB |  | XRAD-1 |
| AA42 | +B265 | DW |  | REV (XRAD-1) |
| AA44 |  | FDB |  | XREAD-1 |
| AA44 | +B282 | DW |  | REV (XREAD-1) |
|  | = øø22 | CREAD | EQU | (*-STETAB)/2-1 |
| AA46 |  | FDB |  | XREST-1 |
| AA46 | +B26A | DW |  | REV (XREST-1) |
| AA48 |  | FDB |  | XRTN-1 |
| AA48 | +B718 | DW |  | REV (XRTN-1) |
| AA4A |  | FDB |  | XRUN-1 |
| AA4A | +B74C | DW |  | REV (XRUN-1) |
| AA4C |  | FDB |  | XSTOP-1 |



## Operator Execution Table

OPETAB - OPERATOR EXECUTION TABLE

- CONTAINS OPERATOR EXECUTION ADR
- MUST BE IN SAME ORDER AS OPNTAB

AA7Ø
AA $7 \varnothing$
AA7Ø +ACB4
AA 72
AA 72 +ACBD
AA 74
AA74 +ACD4
AA 76
AA 76 +ACC4
AA 78
AA78
AA7A
AA7A +ACDB
AA7C
AA7C +B164
AA 7 E
AA7E +AC95
AA8 $\emptyset$
АА8 + АС83
AA82
AA82 +AC8C
;
OPETAB
FDB XPLE-1
DW REV (XPLE-1)
FDB XPNE-1
DW REV (XPNE-1)
FDB XPGE-1
DW REV (XPGE-1)
FDB XPLT-1
DW REV (XPLT-1)
FDB XPGT-1
DW REV (XPGT-1)
FDB XPEQ-1
DW REV (XPEQ-1)
FDB XPPOWER-1
DW REV (XPPOWER-1)
FDB XPMUL-1
DW REV (XPMUL-1)
FDB XPPLUS-1
DW REV (XPPLUS-1)
FDB XPMINUS-1
DW REV (XPMINUS-1)
FDB XPDIV-1
DW REV (XPDIV-1)
FDB XPNOT-1
AA84 +AC9E
AA86
AA86 +ACF8
AA88
AA88 +ACED

## Source Code

| AABA |  | FDB | XPAND-1 |
| :---: | :---: | :---: | :---: |
| AA8A | +ACE 2 | DW | REV (XPAND-1) |
| AABC |  | FDB | XPLPRN-1 |
| AA8C | +AB1E | DW | REV (XPLPRN-1) |
| AA8E |  | FDB | XPRPRN-1 |
| AA8E | +AD7A | DW | REV (XPRPRN-1) |
| AA90 |  | FDB | XPAASN-1 |
| AA90 | +AD5E | DW | REV (XPAASN-1) |
| AA92 |  | FDB | XSAASN-1 |
| AA92 | +AEA2 | DW | REV (XSAASN-1) |
| AA94 |  | FDB | XPSLE-1 |
| AA94 | +ACB4 | DW | REV (XPSLE-1) |
| AA96 |  | FDB | XPSNE-1 |
| AA96 | +ACBD | DW | REV (XPSNE-1) |
| AA98 |  | FDB | XPSGE-1 |
| AA98 | +ACD4 | DW | REV (XPSGE-1) |
| AA9A |  | FDB | XPSLT-1 |
| AA9A | +ACC4 | DW | REV (XPSLT-1) |
| AA9C |  | FDB | XPSGT-1 |
| AA9C | +ACCB | DW | REV (XPSGT-1) |
| AA9E |  | FDB | XPEQ-1 |
| AA9E | +ACDB | DW | REV (XPEQ-1) |
| AAA® |  | FDB | XPUPLUS-1 |
| AAAD | +ACB3 | DW | REV (XPUPLUS-1) |
| AAA2 |  | FDB | XPUMINUS-1 |
| AAA2 | +ACA 7 | DW | REV (XPUMINUS-1) |
| AAA 4 |  | FDB | XPSLPRN-1 |
| AAA 4 | +AE25 | DW | REV (XPSLPRN-1) |
| AAA6 |  | FDB | XPALPRN-1 |
| AAA6 | +AD85 | DW | REV (XPALPRN-1) |
| AAAB |  | FDB | XPDLPRN-1 |
| AAAB | +AD81 | DW | REV (XPDLPRN-1) |
| AAAA |  | FDB | XPFLPRN-1 |
| AAAA | +AD7A | DW | REV (XPFLPRN-1) |
| AAAC |  | FDB | XDPSLP-1 |
| AAAC | +AD81 | DW | REV (XDPSLP-1) |
| AAAE |  | FDB | XPACOM-1 |
| AAAE | +AD78 | DW | REV (XPACOM-1) |
| AABØ |  | FDB | XPSTR-1 |
| AABø | +Bø48 | DW | REV (XPSTR-1) |
| AAB2 |  | FDB | XPCHR-1 |
| AAB2 | +Bø66 | DW | REV (XPCHR-1) |
| AAB4 |  | FDB | XPUSR-1 |
| AAB4 | +ВøВ9 | DW | REV (XPUSR-1) |
| AAB6 |  | FDB | XPASC-1 |
| AAB6 | +Bø11 | DW | REV (XPASC-1) |
| AAB8 |  | FDB | XPVAL-1 |
| AAB8 | +AFFF | DW | REV (XPVAL-1) |
| AABA |  | FDB | XPLEN-1 |
| AABA | +AFC9 | DW | REV (XPLEN-1) |
| AABC |  | FDB | XPADR-1 |
| AABC | +Вø1B | DW | REV (XPADR-1) |
| AABE |  | FDB | XPATN-1 |
| AABE | +B12E | DW | REV (XPATN-1) |
| AACD |  | FDB | xpCOS-1 |
| AACD | +B124 | DW | REV (XPCOS-1) |
| AAC2 |  | FDB | XPPEEK-1 |
| AAC2 | +AFEø | DW | REV (XPPEEK-1) |
| AAC4 |  | FDB | XPS IN-1 |
| AAC4 | +B11A | DW | REV (XPSIN-1) |
| AAC6 |  | FDB | XPRND-1 |
| AAC6 | +Bø8A | DW | REV (XPRND-1) |
| AAC8 |  | FDB | XPFRE-1 |
| AAC8 | +AFEA | DW | REV (XPFRE-1) |
| AACA |  | FDB | XPEXP-1 |
| AACA | +B14C | DW | REV (XPEXP-1) |
| AACC |  | FDB | XPLOG-1 |
| AACC | +B138 | DW | REV (XPLOG-1) |
| AACE |  | FDB | XPL1ø-1 |
| AACE | +B142 | DW | REV (XPLIø-1) |


| AADØ |  | FDB | XPSQR-1 |
| :---: | :---: | :---: | :---: |
| AAD® | +B156 | DW | REV (XPSQR-1) |
| AAD2 |  | FDB | XPSGN-1 |
| AAD2 | +AD18 | DW | REV (XPSGN-1) |
| AAD4 |  | FDB | XPABS-1 |
| AAD4 | +BØAD | DW | REV (XPABS-1) |
| AAD6 |  | FDB | XPINT-1 |
| AAD6 | +BøDC | DW | REV (XPINT-1) |
| AAD8 |  | FDB | XPPDL-1 |
| AAD8 | +Bø21 | DW | REV (XPPDL-1) |
| AADA |  | FDB | XPSTICK-1 |
| AADA | +Bø25 | DW | REV (XPSTICK-1) |
| AADC |  | FDB | XPPTRIG-1 |
| AADC | +Bø29 | DW | REV (XPPTRIG-1) |
| AADE |  | FDB | XPSTRIG-1 |
| AADE | +Bø2D | DW | REV (XPSTRIG-1) |

## Execute Expression

## AAEØ <br> LOCAL

## EXEXPR - Execute Expression

| AAEØ |  | XLET |  |  |
| :---: | :---: | :---: | :---: | :---: |
| AAED |  | EXEXPR |  |  |
| AAEØ | 2ø2EAB | JSR | EXPINT | ; GO INIT |
|  |  |  |  |  |
| AAE3 |  | : EXNXT |  |  |
| AAE 3 | $2 ø 3 \mathrm{EAB}$ | JSR | : EGTOKEN | ; GO GET TOKEN |
| AAE6 | Bøø6 ^AAEE | BCS | : EXOT | BR IF OPERATOR |
| AAE8 | 2 ØВAAB | JSR | ARGPUSH | puSh ARg |
| AAEB | 4CE3AA | JMP | : EXNXT | ; GO FOR NEXT TOKEN |
| AAEE | 85AB | : EXOT STA | EXSVOP | SAVE OPERATOR |
| AAFD | AA | TAX |  |  |
| AAFl | BD2FAC | LDA | OPRTAB-16, X | ; GET OP PREC |
| AAF4 |  | LSRA |  | ; SHIFT FOR GOES ON TO PREC |
| AAF4 | +4A | LSR | A |  |
| AAF5 |  | LSRA |  |  |
| AAF5 | +4A | LSR | A |  |
| AAF6 |  | LSRA |  |  |
| AAF6 | +4A | LSR | A |  |
| AAF7 |  | LSRA |  |  |
| AAF7 | +4A | LSR | A |  |
| AAF8 | 85AC | STA | EXSVPR | ; SAVE GOES ON PREC |
| AAFA | A4A9 | ; EXPTST LDY | OPSTKX | ; GET OP STACK INDEX |
| AAFC | B180 | LDA | [ARGSTK], Y | GET TOP OP |
| AAFE | AA | tax |  |  |
| AAFF | BD2FAC | LDA | OPRTAB-16, X | GET TOP OP PREC |
| ABø2 | 290F | AND | \# |  |
| ABø4 | C5AC | CMP | EXSVPR | [TOP OP]: [NEW OP] |
| ABø6 | 9øøD ^AB15 | BCC | : EOPUSH | ; IF T<N, PUSH NEW |
| ABø8 | AA | TAX |  | Lf POLSE POP |
| ABø9 | Fø14 ^AB1F | BEQ | : EXEND | ; THEN DONE |
|  |  |  |  |  |
| ABøB |  | EXOPOP |  |  |
| ABøB | B180 | LDA | [ARGSTK], Y | RE-GET TOS OP |
| ABøD | E6A9 | INC | OPSTKX | ; DEC OP STACK INDEX |
| ABøF | $2 ø 2 Ø$ AB | JSR | : EXOP | ; GET EXECUTE OP |
| AB12 | 4CFAAA | JMP | : EXPTST | ; GO TEST OP WITH NEW TOS |
| AB15 | A5AB | : EOPUSH LDA | EXSVOP | ; GET OP TO PUSH |
| AB17 | 88 | DEY |  | DEC TO NEXT ENTRY |
| AB18 | $918 \emptyset$ | STA | [ARGSTK], Y | ; SET OP IN Stack |
| AB1A | 84A9 | STY | OPSTKX | ; SAVE NEW OP STACK INDEX |
| AB1C | 4CE3AA | JMP | : EXNXT | ; GO GEt Next token |
| AB1F |  | $\dot{\text { XPLPRN }}$ |  |  |

Source Code

| AB1F | 60 | : EXEND |  | ; | DONE EXECUTE EXPR |
| :---: | :---: | :---: | :---: | :---: | :---: |
| AB2ø |  | : EXOP |  |  |  |
| AB2ø | 38 | SEC |  | ; | SUBSTRACT FOR REL Ø |
| AB21 | E91D | SBC | \#CSROP | ; | Value of first real op |
| AB23 |  | ASLA |  | ; | VALUE * 2 |
| AB23 | $+\emptyset A$ | ASL | A |  |  |
| AB24 | AA | TAX |  |  |  |
| AB25 | BD70AA | LDA | OPETAB, X | ; | PUT OP EXECUTION |
| AB28 | 48 | PHA |  | ; | ROUTINE ON STACK |
| AB29 | BD71AA | LDA | OPETAB+1, X | ; | AND GOTO |
| AB2C | 48 | PHA |  | ; | VIA |
| AB2D | $6 \varnothing$ | RTS |  |  | RTS |

Initialize Expression Parameters

| AB2E |  | EXPINT |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| AB2E | AØFF | LDY | \#\$FF |  |  |
| АВ3б | A911 | LDA | \#CSOE | ; | OPERATOR |
| AB32 | 9180 | STA | [ARGSTK], Y | ; | STACK |
| AB34 | 84A9 | STY | OPSTKX |  |  |
| AB36 | C8 | INY |  | ; | AND INITIALIZE |
| AB37 | 84Bø | STY | COMCNT |  |  |
| AB39 | 84AA | STY | ARSTKX | ; | ARG STACK |
| AB3B | 84Bl | STY | ADFLAG | ; | ASSIGN FLAG |
| AB3D | $6 \square$ | RTS |  |  |  |

## GETTOK - Get Next Token and Classify

| AB3E |  |  | GETTOK |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| AB3E |  |  | : EGTOKEN |  |  |
| AB3E | A4A8 |  | LDY | STINDEX | ; GET STMT INDEX |
| AB4ø | E6A8 |  | INC | STINDEX | ; INC TO NEXT |
| AB42 | B18A |  | LDA | [STMCUR], Y | ; GET TOKEN |
| AB44 | 3043 | ${ }^{\wedge} \mathrm{AB} 89$ | BMI | : EGTVAR | ; BR IF VAR |
| AB46 | C90F |  | - CMP | \# ${ }^{\text {® }}$ F | TOKEN: \$øF |
| AB48 | $9 \emptyset \emptyset 3$ | ${ }^{\wedge} \mathrm{AB} 4 \mathrm{D}$ | BCC | : EGNC | ; BR IF \$ØE, NUMERIC CONST |
| AB4A | Fø13 | ${ }^{\wedge} \mathrm{AB} 5 \mathrm{~F}$ | BEQ | : EGSC | ; BR IF \$øF, STR CONST |
| AB4C | $6 \emptyset$ |  | RTS |  | ; RTN IF OPERATOR |
|  |  |  | ; |  |  |
| AB4D |  |  | NCTOFRØ |  |  |
| AB4D | A2øø |  | : EGNC LDX | \#Ø |  |
| AB4F | C8 |  | :EGT1 INY |  | ; INC LINE INDEX |
| AB5 $\square^{\text {a }}$ | B18A |  | LDA | [STMCUR], Y | ; GET VALUE FROM STMT TBL |
| AB52 | 95D4 |  | STA | FRØ, X | ; AND PUT INTO FRØ |
| AB54 | E8 |  | INX |  |  |
| AB55 | Eのø6 |  | CPX | \# 6 |  |
| AB57 | 9ØF6 | ${ }^{\wedge} \mathrm{AB} 4 \mathrm{~F}$ | BCC | : EGT1 |  |
| AB59 | C8 |  | INY |  | ; INY Y BEYOND CONST |
| AB5A | А9øø |  | LDA | \#EVSCALER | ; ACU=SCALER |
| AB5C | AA |  | TAX |  | ; $\mathrm{X}=$ VAL NO $\emptyset$ |
| AB5D | Fø22 | *AB81 | BEQ | : EGST | ; GO SET REM |
| AB5F | C8 |  | : EGSC INY |  | ; INC Y TO LENGTH BYTE |
| AB60 | B18A |  | LDA | [STMCUR], Y | ; GET LENGTH |
| AB62 | A28A |  | LDX | \#STMCUR | ; POINT TO SMCUR |
| AB64 |  |  | RISC |  |  |
| AB64 | 85D6 |  | STA | VTYPE+EVSLEN | ; SET AS LENGTH |
| AB66 | 85D8 |  | STA | VTYPE+EVSDIM | ; AND DIM |
| AB68 | C8 |  | INY |  |  |
| AB69 | 98 |  | TYA |  | ; ACU=DISPL TO STR |
| AB6A | 18 |  | CLC |  |  |
| AB6B | 7500 |  | ADC | $\emptyset, \mathrm{x}$ | ; DISPL PLUS ADR |
| AB6D | 85D4 |  | STA | VTYPE+EVSADR | ; IS STR ADR |
| AB6F | A9øø |  | LDA | \# $\varnothing$ | ; SET $=\emptyset$ |
| AB71 | 85D 7 |  | STA | VTYPE+EVSLEN+1 | ; LENGTH HIGH |
| AB73 | 85D9 |  | STA | VTYPE+EVSDIM+1 | ; DIM HIGH |
| AB75 | 7501 |  | ADC | 1, X | ; FINISH ADR |
| AB77 | 85D5 |  | STA | VTYPE+EVSADR+1 |  |


| AB79 | 98 | TYA |  |  | ACU=DISPL TO STR |
| :---: | :---: | :---: | :---: | :---: | :---: |
| AB7A | 65D6 | ADC |  | VTYPE+EVSLEN | PLUS STR LENGTH |
| AB7C | A8 | TAY |  |  | ; IS NEW INDEX |
| AB7D | A2øø | LDX |  | \#ø | ; VAR NO = Ø |
| AB7F | A983 | LDA |  | \#EVSTR+EVSDTA + E | DIM ; TYPE = STR |
| AB81 | 85D2 | : EGST | STA | vtype | ; SET TYPE |
| AB83 | 86D3 | STX |  | vnum | ; SET NUM |
| AB85 | 84A8 | STY |  | STINDEX | ; SET NEW INDEX |
| AB87 | 18 | CLC |  |  | ; INDICATE VALUE |
| AB88 | 60 | : EGRTS | RTS |  | ; RETURN |
| AB89 |  | GETVAR |  |  |  |
| AB89 |  | : EGTVAR |  |  |  |
| AB89 | 2ø28AC | JSR |  | GVVTADR | ; GET VVT ADR |
| AB8C | B19D | : EGT 2 | LDA | [WVV'PT], Y | ; MOVE VVT ENTRY |
| AB8E | 99D2ø0 | STA |  | VTYPE, Y | ; TO FRØ |
| AB91 | C8 | INY |  |  |  |
| AB92 | Cøø8 | CPY |  | \#8 |  |
| AB94 | 90F6 ^AB8C | BCC |  | : EGT2 |  |
| AB96 | 18 | CLC |  |  | ; INDICATE VALUE |
| AB97 | 60 | RTS |  |  | ; RETURN |

AAPSTR - Pop String Argument and Make Address Absolute

```
AB98 2ØF2AB AAPSTR JSR ARGPOP ; GO POP ARG
```

GSTRAD - Get String [ABS] Address

| AB9B |  | GSTRAD |  |  |
| :---: | :---: | :---: | :---: | :---: |
| AB9B | A902 | LDA | \#EVSDTA | ; LOAD TRANSFORMED BIT |
| AB9D | 24D2 | BIT | VTYPE | ; TEST STRING ADR TRANSFORM |
| AB9F | DØ15 ^ABB6 | BNE | : GSARTS | ; BR IF ALREADY TRANSFORMED |
| ABAI | Ø5D2 | ORA | VTYPE | ; TURN ON TRANS BIT |
| ABA 3 | 85D2 | STA | VTYPE | ; AND SET |
| ABA 5 |  | RORA |  | ; SHIFT DIM BIT TO CARRY |
| ABA5 | +6A | ROR | A |  |
| ABA6 | 900F ^ABB7 | BCC | : GSND |  |
|  |  | ; |  |  |
| ABA8 | 18 | CLC |  |  |
| ABA9 | A5D4 | LDA | VTYPE+EVSADR | $\begin{aligned} & \text {; STRING ADR = STRING DISPL } \\ & \text { + STARP } \end{aligned}$ |
| ABAB | 658C | ADC | STARP |  |
| ABAD | 85D4 | STA | VTYPE+EVSADR |  |
| ABAF | A8 | TAY |  |  |
| ABBø | A5D5 | LDA | VTYPE+EVSADR+1 |  |
| ABB2 | 658D | ADC | STARP+1 |  |
| ABB4 | 85D5 | STA | VTYPE+EVSADR+1 |  |
| ABB6 | 60 | : GSARTS RTS |  |  |
| ABB7 | 202EB9 | : GSND JSR | ERRDIM |  |

ARGPUSH - Push FR0 to Argument Stack


## Source Code



GETPINT - Get Positive Integer from Expression

| ABD5 |  | GETPINT |  |  |
| :---: | :---: | :---: | :---: | :---: |
| ABD5 | 2ØEØAB | JSR | GETINT | ; GO GET INT |
| ABD8 |  | GETPIø |  |  |
| ABD8 | A5D5 | LDA | FRØ+1 | ; GET HIGH BYTE |
| ABDA | $3 \emptyset \emptyset 1$ ^ABDD | BMI | : GPIERR | ; $\mathrm{BR}>32767$ |
| ABDC | 60 | RTS |  | ; DONE |
| ABDD | 4C32B9 | : GPIERR | ERRL |  |

## GETINT - Get Integer from Expression

| ABEØ | 2ØEØAA | GETINT JSR | EXEXPR | ; EVAL EXPR |
| :--- | :--- | ---: | :--- | :--- | :--- |
| ABE3 |  |  |  |  |
| ABE3 | 2ØF2AB | JSR | ARGPOP | ; POP VALUE TO FRØ |
| ABE6 | $4 C 56 A D ~$ | JMP | CVFPI |  |
|  |  |  |  | RETURN |

## GET1INT - Get One-Byte Integer from Expression

| ABE9 |  | GETIINT |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ABE9 | 20D5AB | JSR | GETPINT |  | GET INT < 32768 |  |  |
| ABEC | DØøl *ABEF | BNE | : ERV1 | ; | IF NOT 1 BYTE, | THEN | ERROR |
| ABEE | 60 | RTS |  |  |  |  |  |
| ABEF |  | : ERVI |  |  |  |  |  |
| ABEF | 203AB9 | JSR | ERVAL |  |  |  |  |

## ARGPOP - Pop Argument Stack Entry to FR0 or FR1

| ABF2 |  |  | ARGPOP |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ABF2 | A5AA |  | LDA | ARSLVL |  | GET ARG STACK LEVEL |
| ABF4 | C6AA |  | DEC | ARSLVL |  | DEC AS LEVEL |
| ABF6 |  |  | ASLA |  |  | AS LEVEL * 8 |
| ABF6 | $+\emptyset A$ |  | ASL | A |  |  |
| ABF7 |  |  | ASLA |  |  |  |
| ABF7 | $+\emptyset A$ |  | ASL | A |  |  |
| ABF8 |  |  | ASLA |  |  |  |
| ABF8 | $+\emptyset A$ |  | ASL | A |  |  |
| ABF9 | A8 |  | TAY |  |  | $Y=$ START OF NEXT ENTRY |
| ABFA | 88 |  | DEY |  | ; | MINUS ONE |
| ABFB | A207 |  | LDX | \#7 |  | $X=7$ FOR 8 |
| ABFD | B18ø |  | ; APOPØ LDA | [ARGOPS], Y |  | ; MOVE ARG ENTRY |
| ABFF | 95D2 |  | STA | VTYPE, X |  |  |
| ACø1 | 88 |  | DEY |  |  | BACKWARDS |
| ACø2 | CA |  | DEX |  |  |  |
| ACø3 | 16 F 8 | ${ }^{*} \mathrm{ABFD}$ | BPL | : APOPØ |  |  |
| AC05 | 60 |  | RTS |  |  | DONE |

## ARGP2 - Pop TOS to FR1,TOS-1 to FR0



POP1 - Get a Value in FR0

|  |  | i |  |  |
| :--- | :--- | :--- | :--- | :--- |
|  |  | POP IT INTO FRø |  |  |

## Source Code

RTNVAR - Return Variable to Variable Value Table from FR0


GVVTADR - Get Value's Value Table Entry Address

| AC28 |  | GVVTADR |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| AC28 | Аøøø | LDY | \#Ø | ; | CLEAR ADR HI |
| AC2A | 849 E | STY | WVVTPT+1 |  |  |
| AC2C |  | ASLA |  | ; | MULT VAR NO |
| AC2C | $+\square A$ | ASL | A |  |  |
| AC2D |  | ASLA |  | ; | BY 8 |
| AC2D | $+\emptyset A$ | ASL | A |  |  |
| AC2E | 269 E | ROL | WVVTPT+1 |  |  |
| АСЗб |  | ASLA |  |  |  |
| АСЗø | $+\emptyset A$ | ASL | A |  |  |
| АСЗ1 | 269E | ROL | WVVTPT+1 |  |  |
| AC33 | 18 | CLC |  | ; | THEN |
| AC34 | 6586 | ADC | VVTP | ; | ADD VVTP VALUE |
| AC36 | 859D | STA | WVVTPT | ; | TO FORM ENTRY |
| AC38 | A587 | LDA | VVTP +1 | ; | ADR |
| AC3A | 659 E | ADC | WVVTPT+1 |  |  |
| AC3C | 859 E | STA | WVVTPT+1 |  |  |
| AC3E | 60 | RTS |  |  |  |

## Operator Precedence Table



## Source Code

| AC5B | 4 E | DB | \$4E | ; CRPRN |
| :---: | :---: | :---: | :---: | :---: |
| AC5C | Fl | DB | \$F1 | ; CAASN |
| AC5D | Fl | DB | \$F1 | ; CSASN |
| AC5E | EE | DB | \$EE | ; CSLE |
| AC5F | EE | DB | \$EE | ; CSNE |
| AC60 | EE | DB | \$EE | ; CSGE |
| AC61 | EE | DB | \$EE | ; CSLT |
| AC62 | EE | DB | \$EE | ; CSGT |
| AC63 | EE | DB | \$EE | ; CSEQ |
| AC64 | DD | DB | \$DD | ; cuplus |
| AC65 | DD | DB | \$DD | ; cuminus |
| AC66 | F2 | DB | \$F2 | ; CSLPRN |
| AC67 | F2 | DB | \$F2 | ; CALPRN |
| AC68 | F2 | DB | \$F2 | ; CDLPRN |
| AC69 | F2 | DB | \$F2 | ; CFLPRN |
| AC6A | F2 | DB | \$F2 | ; CDSLPR |
| AC6B | 43 | DB | \$43 | ; CACOM |
| AC6C | F2 | DB | \$F2 | ; FUNCTIONS |
| AC6D | F2 | DB | \$F2 |  |
| AC6E | F2 | DB | \$F2 |  |
| AC6F | F2 | DB | \$F2 |  |
| AC7ø | F2 | DB | \$F2 |  |
| AC71 | F2 | DB | \$F2 |  |
| AC72 | F2 | DB | \$F2 |  |
| AC73 | F2 | DB | \$F2 |  |
| AC74 | F2 | DB | \$F2 |  |
| AC75 | F2 | DB | \$F2 |  |
| AC76 | F2 | DB | \$F2 |  |
| AC77 | F2 | DB | \$F2 |  |
| AC78 | F2 | DB | \$F2 |  |
| AC79 | F2 | DB | \$F2 |  |
| AC7A | F2 | DB | \$F2 |  |
| AC7B | F2 | DB | \$F2 |  |
| AC7C | F2 | DB | \$F2 |  |
| AC7D | F2 | DB | \$F2 |  |
| AC7E | F2 | DB | \$F2 |  |
| AC7F | F2 | DB | \$F2 |  |
| AC8ø | F2 | DB | \$F2 |  |
| AC81 | F2 | DB | \$F2 |  |
| AC82 | F2 | DB | \$F2 |  |
| AC83 | F2 | DB | \$F2 |  |

## Miscellaneous Operators

## Miscellaneous Operators' Executors

| AC84 |  | XPPLUS |  |
| :---: | :---: | :---: | :---: |
| AC84 | 20ø6AC | JSR | ARGP2 |
| AC87 | $2 \emptyset 3 B A D$ | JSR | FRADD |
| AC8A | 4CBAAB | JMP | ARGPUSH |
| AC8D |  | XPMINUS |  |
| AC8D | 2øø๐AC | JSR | ARGP2 |
| AC90 | 2041 AD | JSR | FRSUB |
| AC93 | 4CBAAB | JMP | ARGPUSH |
| AC96 |  | XPMUL |  |
| AC96 | 2øø6AC | JSR | ARGP 2 |
| AC99 | $2 \emptyset 47 \mathrm{AD}$ | JSR | FRMUL |
| AC9C | 4CBAAB | JMP | ARGPUSH |
| AC9F |  | XPDIV |  |
| AC9F | 2øø6AC | JSR | ARGP2 |
| ACA2 | 2ø4DAD | JSR | FRDIV |
| ACA5 | 4CBAAB | JMP | ARGPUSH |
| ACA8 |  | XPUMINUS |  |
| ACA8 | 2øF2AB | JSR | ARGPOP |
| ACAB | A5D4 | LDA | FRø |
| ACAD | 498ø | EOR | \#\$80 |
| ACAF | 85D4 | STA | FRØ |
| ACB1 | 4CBAAB | JMP | ARGPUSH |
| ACB4 |  | XPUPLUS |  |

[^1]| ACB4 | 60 | RTS |  |  |
| :---: | :---: | :---: | :---: | :---: |
| ACB5 | XPLE |  |  |  |
| ACB5 |  | XPSLE |  |  |
| ACB5 | 2ø26AD | JSR | XCMP |  |
| ACB8 | 304B "ADø5 | BMI | XTRUE |  |
| ACBA | F049 ^AD65 | BEQ | Xtrue |  |
| ACBC | 1042 *ADøø | BPL | XFALSE |  |
| ACBE |  | XPNE |  |  |
| ACBE |  | XPSNE |  |  |
| ACBE | 2026AD | JSR | XCMP |  |
| ACCl | Fø3D *ADøø | BEQ | xfalse |  |
| ACC 3 | Dø4の ^ADØ5 | BNE | XTRUE |  |
| ACC5 |  | XPLT XPSLT |  |  |
| ACC5 |  |  |  |  |
| ACC5 | 2ø26AD | JSR | XCMP |  |
| ACC8 | 3ø3B ^ADø5 | BMI | Xtrue |  |
| ACCA | 1 1034 ^ADøø | BPL | xfalse |  |
| ACCC |  | XPGT |  |  |
| ACCC |  | XPSGT |  |  |
| ACCC | 2026AD | JSR | XCMP |  |
| ACCF | 3Ø2F ^ADøø | BMI | XFALSE |  |
| ACD1 | Fø2D ^ADøø | BEQ | xfalse |  |
| ACD 3 | 1ø3ø ^AD65 | BPL | XTRUE |  |
| ACD5 |  | XPGE XPSGE |  |  |
| ACD5 |  |  |  |  |
| ACD5 | 2026AD | JSR | XCMP |  |
| ACD8 | $3 \varnothing 26$ ^ADøб | BMI | Xfalse |  |
| ACDA | $1 \varnothing 29$ ^AD65 | BPL | Xtrue |  |
| ACDC |  | XPEQ XPSEO |  |  |
| ACDC |  |  |  |  |
| ACDC | 2026AD | JSR | XCMP |  |
| ACDF | F624 ^AD65 | BEQ | XtRUE |  |
| ACE1 | Dø1D ^ADøø | BNE | xfalse |  |
|  |  | ; |  |  |
| ACE 3 |  | XPAND |  |  |
| ACE 3 | 20ø6AC | JSR | ARGP2 |  |
| ACE6 | A5D 4 | LDA | FRø |  |
| ACE8 | 25Eø | AND | FR1 |  |
| ACEA | Fø14 ^ADøø | BEQ | xfalse |  |
| ACEC | D017 ^ADø5 | BNE | xtrue |  |
| ACEE |  | XPOR |  |  |
| ACEE | 20ø6AC | JSR | ARGP2 |  |
| ACFl | A5D4 | LDA | FRø |  |
| ACF 3 | Ø5Eø | ORA | FR1 |  |
| ACF5 | Føø9 ^ADøø | BEQ | xpalse |  |
| ACF7 | DøøC ^AD65 | BNE | Xtrue |  |
| ACF9 |  | XPNOT |  |  |
| ACF9 | 2øF2AB | JSR | ARGPOP |  |
| ACFC | A5D4 | LDA | FRø |  |
| ACFE | Føø5 ^ADØ5 | BEQ | xtrue |  |
|  |  | ; | FALL THROUGH TO | XFALSE |
|  |  | ; |  |  |
|  |  | X XFALSE |  |  |
| ADøø |  |  |  |  |
| ADøø | A9øø | LDA | \#ø |  |
| ADø2 | A8 | tay |  |  |
| AD®3 | Føø4 ^ADø9 | BEQ | XTF |  |
|  |  | $\dot{;} \dot{\text { XTRUE }}$ |  |  |
| ADø5 |  |  |  |  |
| AD65 | A94Ø | LDA | \#\$4ø |  |
| AD®7 |  | XTI |  |  |
| AD®7 | Aøø1 | LDY | \#1 |  |
|  |  | $\dot{X}_{\mathrm{X} T \mathrm{~F}}$ |  |  |
| AD69 |  |  |  |  |
| AD69 | 85D4 | STA | FRø |  |
| ADøB | 84D5 | STY | FRø+1 |  |
| ADøD | A2D6 | LDX | \#FRØ+2 | ; POINT TO PART TO CLEAR |
| ADøF | Aøø4 | LDY | \#FPREC-2 | ; GET \# OF bytes to clear |
| ADll | 2048DA | JSR | ZXLY | ; clear rest of frø |
| AD14 | 85D2 | STA | VTYPE |  |
| ADl6 |  | XPUSH |  |  |
| AD16 | 4CBAAB | JMP | ARgPuSh |  |

## Source Code

## XPSGN - Sign Function



XCMP - Compare Executor

| AD26 |  | XCMP |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| AD26 | A4A9 | LDY |  | OPSTKX | ; | GET | OPERATOR | HAT |
| AD28 | 88 | DEY |  |  | ; | GOT | US HERE |  |
| AD29 | B18ø | LDA |  | [ARGSTK], Y |  |  |  |  |
| AD2B | C92F | CMP |  | \#CSLE | ; | IF | OP WAS ARIT | HMETIC |
| AD2D | 9003 *AD32 | BCC |  | FRCMPP | ; | THEN | DO FP REG | COMP |
| AD2F | 4C81AF | JMP |  | STRCMP | ; | ELSE | DO STRING | COMPARE |
| AD32 | 2Øロ6AC | FRCMPP JSR ARGP2 |  |  |  |  |  |  |

FRCMP - Compare Two Floating Point Numbers


FRADD - Floating Point Add


## FRSUB - Floating Point Subtract



FRMUL - Floating Point Multiply



XPAASN - Arithmetic Assignment Operator


## XPACOM - Array Comma Operator

| AD79 | XPACOM |  |
| :--- | ---: | ---: | ---: |
| AD79 | E6Bø | INC $\quad$ COMCNT INCREMENT COMMA COUNT |

XPRPRN - Right Parenthesis Operator

| ; XPFLPRN - FUNCTION RIGHT PAREN OPERATOR |  |  |
| :--- | :--- | :--- |
| ; |  |  |
| XPRPRN |  |  |
| XPFLPRN |  |  |
| LDY | OPSTKX |  |
| PLA |  | GOT OPERATOR STACK TOP |
| PLA | EXOPOP |  |
| JMP |  |  |
|  |  |  |

## XPDLPRN - DIM Left Parenthesis Operator

| AD82 | XDPSLP |  |  |  |  |  |
| ---: | ---: | ---: | :--- | :---: | :---: | :---: |
| AD82 |  | XPDLPRN |  |  |  |  |
| AD82 | A94ø | LDA | $\# \$ 4 \varnothing$ |  |  |  |
| AD84 | 85B1 | STA | ADFLAG |  |  |  |

Source Code

| XPALPRN - Array Left Parenthesis Operator |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| AD86 |  | XPALPRN |  |  |  |
| AD86 | 24Bl | BIT |  | ADFLAG | ; IF NOT ASSIGN |
| AD88 | 1006 *AD90 | BPL |  | :ALPl | ; THE BRANCH |
|  |  | ; |  |  | ELSE |
| AD8A | A5AA | LDA |  | ARSLVL | ; SAVE STACK LEVEL |
| AD8C | 85AF | STA |  | ATEMP | ; OF THE VALUE ASSIGNMENT |
| AD8E | C6AA | DEC |  | ARSLVL | ; AND PSEUDO POP IT |
| AD90 | A96Ø | : ALP1 | LDA | \#ø | ; INIT FOR I2 = Ø |
| AD92 | A8 | TAY |  |  |  |
| AD93 | C5Bø | CMP |  | COMCNT | ; IF COMMA COUNT $=\varnothing$ THEN |
| AD95 | FøØB ^ADA2 | BEQ |  | : ALP2 | ; BR WITH I2 = Ø |
| AD97 | C6B0 |  |  | COMCNT | ELSE |
| AD99 | 2øE3AB | JSR |  | GTINTO | ; ELSE POP I2 AND MAKE INT |
| AD9C | A5D5 | LDA |  | FRØ+1 |  |
| AD9E | 3623 ^ADC3 | BMI |  | : ALPER | ; ERROR IF > 32,767 |
| ADAø | A4D4 | LDY |  | FRø |  |
| ADA 2 | 8598 | ; ALP2 | STA | INDEX $2+1$ | ; SET 12 VALUE |
| ADA 4 | 8497 | STY |  | INDEX2 |  |
|  |  | , |  |  |  |
| ADA6 | 2ØE3AB | JSR |  | GTINTO | ; POP I2 AND MAKE INT |
| ADA9 | A5D4 | LDA |  | FRø | ; MOVE Il |
| ADAB | 85F5 | STA |  | ZTEMP1 | ; TO ZTEMP1 |
| ADAD | A5D5 | LDA |  | FRØ+1 |  |
| ADAF | $3012{ }^{\wedge}$ ADC3 | BMI |  | : ALPER | ; ERROR IF > 32,767 |
| ADB1 | 85F6 | STA |  | ZTEMP1+1 |  |
| ADB3 |  | JSR |  |  |  |
|  | 20F2AB | JSR |  | ARGPOP | ; POP THE ARRAY ENTRY |
|  |  | ; |  |  |  |
| ADB6 | 24B1 | BIT |  | ADFLAG | ; IF NOT EXECUTING DIM |
| ADB8 | $5 \emptyset \emptyset 5$ ^ADBF | BVC |  | :ALP3 | ; THEN CONTINUE |
| ADBA | A90ø | LDA |  | \# $\emptyset$ | ; TURN OFF DIM BIT |
| ADBC | 85B1 | STA |  | ADFLAG | ; IN ADFLAG |
| ADBE | 60 | RTS |  |  | ; AND RETURN |
|  |  | ; |  |  |  |
| ADBF | 66D2 | ROR |  | VTYPE | ; IF ARRAY HAS BEEN |
| ADCl | BøØ3 ^ADC6 | BCS |  | :ALP4 | ; DIMED THEN CONTINUE |
| ADC3 | 2ø2EB9 | : ALPER | JSR | ERRDIM | ; ELSE DIM ERROR |
| ADC6 |  | :ALP4 |  |  |  |
|  |  |  |
| ADC6 | A5F6 |  |  |  | LDA |  | ZTEMP1+1 | ; TEST INDEX 1 |
| ADC8 | C5D7 | CMP |  | VTYPE+EVADI+1 | ; IN RANGE WITH |
| ADCA | $90 \emptyset 8$ *ADD4 | BCC |  | :ALP5 | ; DIM1 |
| ADCC | DøF5 ^ADC3 | BNE |  | : ALPER |  |
| ADCE | A5F5 | LDA |  | ZTEMP1 |  |
| ADDø | C5D6 | CMP |  | VTYPE+EVAD1 |  |
| ADD 2 | BØEF ^ADC3 | BCS |  | : ALPER |  |
|  | A598 | : ALP5 | LDA | INDEX2+1 | ;TEST INDEX 2 |
| ADD6 | C5D9 | CMP |  | VTYPE+EVAD2+1 | ; IN RANGE WITH |
| ADD8 | $90 \emptyset 8$ ^ADE2 | BCC |  | :ALP6 | ; DIM 2 |
| ADDA | DØE7 ^ADC3 | BNE |  | : ALPERR |  |
| ADDC | A 597 | LDA |  | INDEX2 |  |
| ADDE | C5D8 | CMP |  | VTYPE+EVAD2 |  |
| ADEØ | BøE1 ^ADC3 | BCS |  | : ALPER |  |
|  |  | ; ${ }^{\text {a }}$ |  |  |  |
| ADE 2 | 205DAF | : ALP6 | JSR | AMUL 1 | ; INDEXI = INDEXI |
| ADE 5 | A597 | LDA |  | INDEX2 | ; INDEX1 = INDEX1 + INDEX2 |
| ADE 7 | A498 | LDY |  | INDEX2+1 |  |
| ADE9 | 2652AF | JSR |  | AADD |  |
| ADEC | 2046AF | JSR |  | AMUL2 | ; $\mathrm{ZTEMP1}=\mathrm{ZTEMP1}{ }^{\text {* } 6}$ |
| ADEF | A5D4 | LDA |  | VTYPE+EVAADR | ; ZTEMP1 = ZTEMP1 + DISPL |
| ADFl | A4D5 | LDY |  | VTYPE+EVAADR+1 |  |
| ADF 3 | 2052 AF | JSR |  | AADD |  |
| ADF6 | A58C | LDA |  | STARP | ; $\mathrm{ZTEMP1}=\mathrm{ZTEMP1}+\mathrm{ADR}$ |

## Source Code

| ADF8 | A 48 D | LDY |  | STARP+1 |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| ADFA | 2052AF | JSR |  | AADD |  |
|  |  | ; |  |  | ZTEMP1 NOW POINTS |
|  |  | ; |  |  | TO ELEMENT REQD |
| ADFD | 24B1 | BIT |  | ADFLAG | ; IF NOT ASSIGN |
| ADFF | 1015 ^AE16 | BPL |  | : ALP8 | ; THEN CONTINUE |
|  |  | ; |  |  | ELSE ASSIGN |
| AEØ1 | A 5AF | LDA |  | ATEMP | ; RESTORE ARG LEVEL |
| AEØ3 | 85AA | STA |  | ARSLVL | ; TO VALUE AND |
| AEØ5 | 2ØF2AB | JSR |  | ARGPOP | ; POP VALUE |
|  |  | ; |  |  |  |
| AEø8 | AøØ5 | LDY |  | \# 5 |  |
| AEØA | B9D4øø | : ALP7 | LDA | FRø, Y | ; MOVE VALUE |
| AEØD | 91F5 | STA |  | [ZTEMP1],Y | ; TO ELEMENT SPACE |
| AEØF | 88 | DEY |  |  |  |
| AE10 | 10 F 8 ~AEØA | BPL |  | : ALP7 |  |
| AE12 | C8 | INY |  |  | ; TURN OFF |
| AE13 | 84 Bl | STY |  | ADFLAG | ; ADFLAG |
| AEl 5 | 60 | RTS |  |  | ; DONE |
|  |  | ; |  |  |  |
| AE16 | AØØ5 | : ALP8 | LDY | \#5 |  |
| AE18 | B1F5 | : ALP9 | LDA | [ZTEMP1], Y | ; MOVE ELEMENT TO |
| AE1A | 99D4ØØ | STA |  | FRØ, Y | ; FRØ |
| AElD | 88 | DEY |  |  |  |
| AElE | 10 F 8 ^AE18 | BPL |  | : ALP9 |  |
|  |  | ; |  |  |  |
| AE2Ø | C8 | INY |  |  |  |
| AE21 | 84D2 | STY |  | VTYPE |  |
| AE23 | 4CBAAB | JMP |  | ARGPUSH | ; PUSH FRØ BACK TO STACK |
|  |  | ; |  |  | AND RETURN |

## XPSLPRN — String Left Parenthesis

| AE26 |  | XPSLPRN |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| AE26 | A5Bø | LDA |  | COMCNT | ; IF NO INDEX 2 |
| AE28 | FØø7 *AE31 | BEQ |  | : XSLP2 | ; THEN BR |
|  |  | ; |  |  |  |
| AE2A | 2096AE | JSR |  | : XSPV | ; ELSE POP I2 AND |
| AE2D | 8498 | STY |  | INDEX2+1 | ; SAVE IN INDEX 2 |
| AE2F | 8597 | STA |  | INDEX2 |  |
|  |  | ; |  |  |  |
| AE31 | 2096AE | : XSLP2 | JSR | : XSPV | ; POP INDEX 1 |
| AE34 | 38 | SEC |  |  | ; ADD DECREMENT BY ONE |
| AE35 | E901 | SBC |  | \# 1 | ; AND PUT INTO ZTEMPl |
| AE37 | 85F5 | STA |  | ZTEMP1 |  |
| AE39 | 98 | TYA |  |  |  |
| AE3A | E90ø | SBC |  | \# $\varnothing$ |  |
| AE3C | 85F6 | STA |  | ZTEMP1+1 |  |
|  |  | ; |  |  |  |
| AE3E | 2ØF2AB | JSR |  | ARGPOP | ; POP ARG STRING |
|  |  | ; |  |  |  |
| AE41 | A5B1 | LDA |  | ADFLAG | ; IF NOT A DEST STRING |
| AE43 | 10ØB ^AE50 | BPL |  | : XSLP3 | ; THEN BRANCH |
| AE45 | 65BØ | ORA |  | COMCNT |  |
| AE47 | 85B1 | STA |  | ADFLAG |  |
| AE49 | A4D9 | LDY |  | VTYPE+EVSDIM+1 | ; INDEX 2 LIMIT |
| AE4B | A5D8 | LDA |  | VTYPE+EVSDIM | ; IS DIM |
| AE4D | 4C54AE | JMP |  | : XSLP4 |  |
|  |  | ; |  |  |  |
| AE5Ø | A5D6 | : XSLP3 | LDA | VTYPE+EVSLEN | ; INDEX 2 LIMIT |
| AE52 | A4D7 | LDY |  | VTYPE+EVSLEN+1 | ; IS STRING LENGTH |
| AE 54 | A6Bø | ; ${ }^{\text {; }}$ XSLP4 | LDX | COMCNT | ; IF NO INDEX 2 |
| AE56 | F61Ø ^AE68 | BEQ |  | : XSLP6 | ; THEN BRANCH |
| AE58 | C6Bø | DEC |  | COMCNT | ; ELSE |
| AE5A | C498 | CPY |  | INDEX2+1 |  |
| AE5C | 9035 ^AE93 | BCC |  | : XSLER |  |
| AE5E | Døø4 ^AE64 | BNE |  | : XSLP5 | ; INDEX 2 LIMIT |
| AE60 | C597 | CMP |  | INDEX2 |  |
| AE62 | 902F ^AE93 | BCC |  | : XSLER |  |

## Source Code

| AE64 | A498 | : XSLP5 | LDY | INDEX2+1 | ;USE INDEX 2 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| AE66 | A597 | LDA |  | INDEX2 | ;AS LIMIT |
| AE68 | 38 | : $\mathrm{XSLP6}$ | SEC |  | ; LENGTH IS |
| AE69 | E5F5 | SBC |  | ZTEMP1 |  |
| AE6B | 85D6 | STA |  | VTYPE+EVSLEN | ; LIMIT - INDEX 1 |
| AE6D | AA | TAX |  |  |  |
| AE6E | 98 | TYA |  |  |  |
| AE6F | E5F6 | SBC |  | ZTEMP1+1 |  |
| AE71 | 85D7 | STA |  | VTYPE+EVSLEN+1 |  |
| AE73 | 901 E ^AE93 | BCC |  | : XSLER | ; LENGTH MUST BE |
| AE75 | A8 | TAY |  |  | ; GE ZERO |
| AE76 | Døø3 ^AE7B | BNE |  | : XSLP7 |  |
| AE78 | 8A | TXA |  |  |  |
| AE79 | F018 *AE93 | BEQ |  | : XSLER |  |
| AE7B | 209BAB | ; XSLP7 | JSR | GSTRAD | ; GET ABS ADR |
| AE7E | 18 | ; CLC |  |  |  |
| AE7F | A5D4 | LDA |  | VTYPE+EVSADR |  |
| AE81 | 65F5 | ADC |  | ZTEMPI | ; STRING ADR |
| AE83 | 85D4 | STA |  | VTYPE+EVSADR | ; STRING ADR + INDEX 1 |
| AE85 | A5D5 | LDA |  | VTYPE+EVSADR+1 |  |
| AE87 | 65F6 | ADC |  | ZTEMP1+1 |  |
| AE89 | 85D5 | STA |  | VTYPE+EVSADR+1 |  |
|  |  | ; |  |  |  |
| AE8B | 24B1 | BIT |  | ADFLAG | ; IF NOT ASSIGN |
| AE8D | $1 \varnothing \emptyset 1$ *AE9Ø | BPL |  | : XSLP8 | ; THEN BR |
| AE8F | 60 | RTS |  |  | ; ELSE RETURN TO ASSIGN |
| AE9Ø | 4CBAAB | ; XSLP8 | JMP | ARGPUSH | ; PUSH ARG AND RETURN |
| AE93 | 2ø36B9 | : XSLER | JSR | ERRSSL |  |

XSPV - Pop Index Value as Integer and Insure Not Zero

| AE96 |  | : XSPV |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| AE96 | 2øE3AB | JSR |  | GTINTO | ; | GO GET THE INTEGER |
| AE99 | A5D4 | LDA |  | FRø | ; | GET VALUE LOW |
| AE9B | A4D5 | LDY |  | FRø+1 | ; | GET VALUE HI |
| AE9D | DØ03 ^AEA2 | : XSPV1 | BNE | : XSPVR | ; | RTN IF VH NOT ZERO |
| AE9F | AA | TAX |  |  | ; | TEST VL |
| AEAØ | F0Fl ^AE93 | BEQ |  | : XSLER | ; | BR VL, VH $=\emptyset$ |
| AEA2 | 60 | : XSPVR | RTS |  | ; | DONE |

XSAASN - String Assign Operator



## Source Code

| AF32 | 85B1 | STA |  | ADFLAG | ; | FLAG |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | ; |  |  |  | ELSE FOR EXPLICT LENGTH |
| AF 34 | E4D7 | CPX |  | VTYPE+EVSLEN+1 | ; | IF NEW LENGTH |
| AF36 | 90ø6 ^AF3E | BCC |  | :XSA6 | , | GREATER THAN |
| AF38 | Døø5 ^AF3F | BNE |  | :XSA5 | ; | OLD LENGTH THEN |
| AF3A | C4D6 | CPY |  | VTYPE+EVSLEN | ; | SET NEW LENGTH |
| AF3C | Bøø1 *AF3F | BCS |  | : XSA5 | ; | ELSE DO NOTHING |
| AF3E | 60 | : XSA6 | RTS |  |  |  |
|  |  | ; |  |  |  |  |
| AF3F | 84D6 | : XSA5 | STY | VTYPE+EVSLEN |  |  |
| AF41 | 86D7 | STX |  | VTYPE+EVSLEN+1 |  |  |
| AF43 | 4Cl6AC | JMP |  | RTNVAR |  |  |

## AMUL2 - Integer Multiplication of ZTEMP1 by 6



AADD - Integer Addition of [A,Y] to ZTEMP1

| AF52 |  | AADD |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| AF52 | 18 |  | CLC |  |  |  |  |  |
| AF53 | 65 F 5 | ADC |  | ADC | ZTEMP1 | ; | ADD LOW | ORDER |
| AF55 | 85F5 |  | STA |  | ZTEMP1 |  |  |  |
| AF57 | 98 |  | TYA |  |  |  |  |  |
| AF58 | 65F6 |  | ADC |  | ZTEMP1+1 | ; | ADD HIGH | ORDER |
| AF 5A | 85F6 |  | STA |  | ZTEMP1+1 |  |  |  |
| AF5C | 60 |  | RTS |  |  | ; | DONE |  |

## AMUL - Integer Multiplication of ZTEMP1 by DIM2



## STRCMP - String Compare

| AF81 | STRCMP |  |  |  |  |  |
| :--- | ---: | ---: | :--- | :--- | :--- | :--- |
| AF81 | $2098 A B$ | JSR | AAPSTR | POP STRING WITH ABS ADR |  |  |
| AF84 | 2øB6DD | JSR | MVØTO1 | MOVE B TO FRI |  |  |
| AF87 | $2 \emptyset 98 A B$ | JSR | AAPSTR | ; POP STRING WITH ABS ADR |  |  |


| AF8A | A2D6 | $: \text { SCl }$ | LDX | \#FRØ-2+EVSLEN | ; GO DEC STR A LEN |
| :---: | :---: | :---: | :---: | :---: | :---: |
| AF8C | 2øBCAF | JSR |  | zPADEC |  |
| AF8F | ø8 | PHP |  |  | SAVE RTN CODE |
| AF9ø | A2E2 | LDX |  | \#FRI-2+EVSLEN | GO DEC STR B LEN |
| AF92 | $2 ø \mathrm{BCAF}$ | JSR |  | zPADEC |  |
| AF95 | Fø13 ^AFAA | BEQ |  | : SC2 | BR STR B LEN $=\varnothing$ |
| AF97 | 28 | PLP |  | ; | get Str a cond code |
| AF98 | FøøD ^AFA7 | BEQ |  | : SCLT | BR STR A LEN $=\varnothing$ |
| AF9A | Аøøø | LDY |  | \# | COMPARE A BYTE |
| AF9C | BlD4 | LDA |  | [FRø-2+EVSADR], Y | ; OF String A |
| AF9E | DIEø | CMP |  | [FRI-2+EVSADR], Y | ; TO STRING B |
| AFAØ | FøøC ^AFAE | BEQ |  | :SC3 ; | BR IF SAME |
| AFA2 | $9 \varnothing \varnothing 3$ ^AFA7 | BCC |  | : SCLT | BR IF A<B |
| AFA4 | A901 | : SCGT | LDA | \#1 ; | A $>B$ |
| AFA6 | 60 | RTS |  |  |  |
| AFA 7 | A98ø | :SCLT | LDA | \#\$8ø ; | A<B |
| AFA9 | 60 | RTS |  |  |  |
| AFAA | 28 | ; Sc 2 | PLP |  | IF STR A LEN NOT |
| AFAB | D®F7 A AFA4 | BNE |  | : SCGT | ZERO THEN A>B |
| AFAD | 60 | : SCEQ | RTS | ; | ELSE $\mathrm{A}=\mathrm{B}$ |
| AFAE | E6D4 | : Sc3 | INC | FRø-2+EVSADR | ; INC STR A ADR |
| AFBø | Døø2 *AFB4 | BNE |  | : SC4 |  |
| AFB2 | E6D5 | INC |  | FRø-1+EVSADR |  |
| AFB4 | E6Eø | : SC4 | INC | FRI-2+EVSADR | ; INC STR B ADR |
| AFB6 | DøD2 ^AF8A | BNE |  | : SCl |  |
| AFB8 | E6E1 | INC |  | FRI-1+EVSADR |  |
| AFBA | DøCE *AF8A | BNE |  | :SCl |  |

## ZPADEC - Decrement a Zero-Page Double Word



## Functions

## XPLEN - Length Function

| AFCA |  | XPLEN |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| AFCA | 2098AB | JSR |  | AAPSTR | ; | POP STRING WITH ABS ADR |
| AFCD | A5D6 | LDA |  | VTYPE+EVSLEN | ; | MOVE LENGTH |
| AFCF | A4D7 | LDY |  | VTYPE+EVSLEN+1 |  |  |
| AFD1 |  | XPIFP |  |  |  |  |
| AFD1 | 85D4 | STA |  | FRø | ; | TO TOP OF FRø |
| AFD3 | 84D5 | STY |  | FRø+1 |  |  |
| AFD5 | 2ØAAD9 | XPIFPI | JSR | CVIFP | ; | AND CONVERT TO FP |
| AFD8 |  | XPIFP2 |  |  |  |  |
| AFD8 | A900 | ; LDA |  | \# $\varnothing$ | ; | CLEAR |
| AFDA | 85D2 | STA |  | VTYPE | ; | TYPE AND |
| AFDC | 85D3 | STA |  | VNUM | ; | NUMBER |
| AFDE | 4CBAAB | JMP |  | ARGPUSH | ; | PUSH TO STACK AND RETURN |

## XPPEEK - PEEK Function

| AFE1 | XPPEEK |  |  |  |
| :--- | :--- | ---: | :--- | :--- |
| AFE1 | 2ØE3AB | JSR | GTINTO | ; GET INT ARG |
| AFE4 | AØø | LDY | \# |  |
| AFE6 | BlD4 | LDA | [FRØ],Y | ; GET MEM BYTE |
| AFE8 | 4CDIAF | JMP | XPIFP | ; GO PUSH AS FP |

## Source Code

## XPFRE - FRE Function

| AFEB |  | XPFRE |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| AFEB | 20F2AB | JSR | ARGPOP | ; | POP DUMMY ARG |
| AFEE | 38 | SEC |  |  |  |
| AFEF | ADE502 | LDA | HIMEM | ; | NO FREE BYTES |
| AFF2 | E590 | SBC | MEMTOP | ; | $=$ HIMEM-MEMTOP |
| AFF4 | 85D4 | STA | FRØ |  |  |
| AFF6 | ADE602 | LDA | HIMEM +1 |  |  |
| AFF9 | E591 | SBC | MEMTOP +1 |  |  |
| AFFB | 85D5 | STA | FRØ+1 |  |  |
| AFFD | 4CD5AF | JMP | XPIFP1 | ; | GO PUSH AS FP |

XPVAL - VAL Function


Restore Character


XPSTICK - Function Joystick


## Source Code

| Вø3ø |  | : GRF |  |  |
| :---: | :---: | :---: | :---: | :---: |
| Вø30 | 48 | PHA |  |  |
| Bø31 | 20E3AB | JSR | GTinto | GET INTEGER FROM STACK |
| Bø34 | A5D5 | LDA | FR6+1 | ; HIGH ORDER BYTE |
| Bø36 | DøøE ^BØ46 | BNE | : ERGRF | ; SHOULD BE = |
| Вø38 | A5D 4 | LDA | FRø | ; GET \# |
| Bø3A | 68 | PLA |  | ; GET DISPL FROM BASE |
| Вø3B | 18 | CLC |  |  |
| вø3C | 65D4 | ADC | FRø | ; ADD MORE DISPL |
| Bø3E | AA | TAX |  |  |
|  |  | ; |  |  |
| Bø3F | BD7902 | LDA | GRFBAS, X | ; GET VALUE |
| Bø42 | Аøøø | LDY | \# $\emptyset$ |  |
| Bø44 | Fø8B ^AFD1 | BEQ | XPIFP | ; GO CONVERT \& PUSH ON STACK |
|  |  | ; ; |  |  |
|  |  | ; |  |  |
| Bø46 |  | : ERGRF |  |  |
| Bø46 | 2ø3AB9 | JSR | ERVAL |  |

XPSTR - STR Function

| Bø49 |  | XPSTR |
| :--- | :--- | :--- |
| Bø49 | $2 \emptyset F 2 A B$ | JSR |
| Bø4C | $2 \emptyset E 6 D 8$ | $;$ |


| ARGPOP | ; GET VALUE IN FRØ |
| :--- | :--- |
| CVFASC | ; CONVERT TO ASCII |

Build String Element

| Bø4F | A5F3 |  | LDA | INBUFF | ; | SET ADDR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B051 | 85D4 |  | STA | FRØ-2+EVSADR | ; |  |
| B053 | A5F4 |  | LDA | INBUFF+1 |  |  |
| B655 | 85D5 |  | STA | FRØ-1+EVSADR |  |  |
| Get Length |  |  |  |  |  |  |
| Bø57 | AøFF |  | LDY | \# SFF | ; | INIT FOR LENGTH COUNTER |
| Bø59 |  |  | : XSTR1 |  |  |  |
| B059 | C8 |  | INY |  | ; | BUMP COUNT |
| Bø5A | B1F3 |  | LDA | [INBUFF], Y | ; | GET CHAR |
| B05C | 1øFB | ^B059 | BPL | : XSTR1 | ; | IS MSB NOT ON, REPEAT |
| Bø5E | 297F |  | AND | \# \$7F | ; | TURN OFF MSB |
| Bø6Ø | 91F3 |  | STA | [INBUFF], Y | ; | RETURNS CHAR TO BUFFER |
| Bø62 | C8 |  | INY |  | ; | INC TO GET LENGTH |
| Bø63 | 84D6 |  | ; STY | FRØ-2+EVSLEN | ; | SET LENGTH LOW |
| BØ65 | Dø17 | ^BØ7E | ; BNE | : CHR |  | JOIN CHR FUNCTION |

XPCHR - CHR Function

| Bø67 |  | XPCHR |
| :---: | :---: | :---: |
| Bø67 | 20F2AB | JSR |
|  |  | ; |
| B66A | 2056AD | JSR |
| B66D | A5D4 | LDA |
| Bø6F | 8DCøø5 | STA |

## Build String Element

| B072 | A905 | LDA | \# (LBUFF+\$4ø)/256 | ; SET ADDR |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| B074 | 85D5 | STA | FRØ-1+EVSADR | X |  |
| Bø76 | A9Cø | LDA | \# (LBUFF+\$4Ø) \& 255 | ; X |  |
| B978 | 85D4 | STA | FRØ-2+EVSADR | X |  |
|  |  | ; |  |  |  |
| Bø7A | A901 | LDA | \#1 | SET LENGTH | LOW |
| B97C | 85D6 | STA | FRØ-2+EVSLEN | X |  |
| Bø7E |  | : CHR |  |  |  |
| B97E | A90ø | LDA | \# | SET LENGTH | HIGH |
| Bø8ø | 85D7 | STA | FRØ-1+EVSLEN ; | X |  |

## Source Code



| Вø8B |  | XPRND |  |  |
| :---: | :---: | :---: | :---: | :---: |
| вøвв | A2A8 | LDX | \#RNDDIV\&255 | ; POINT TO 65535 |
| Bø8D | АøВø | LDY | \#RNDDIV/256 | ; X |
| Bø8F | 2098DD | JSR | FLDlR | ; MOVE IT TO FRI |
| B692 | 2øF2AB | JSR | ARGPOP | ; Clear dummy arg |
|  |  | LDY | RNDLOC |  |
| Bø95 | ACøAD2 |  |  | ; GET 2 BYTE RANDOM |
| Bø98 | 84D4 | STY | FRø | ; X |
| Bø9A | ACøAD2 | LDY | RNDLOC | ; X |
| Bø9D | 84D5 | STY | FRØ+1 | ; X |
| BØ9F | 20AAD9 | JSR | CVIFP | ; CONVERT TO INTEGER |
| BøA2 | 204DAD | JSR | FRDIV | ;DO DIVIDE |
| BØA5 | 4 CBAAB | JMP | ARGPUSH | ; PUT ON STACK |
|  |  | ; |  |  |
|  |  |  |  |  |
|  |  | R |  |  |
| ВØA8 | $\begin{aligned} & 42 ø 65536 ø \varnothing \\ & \emptyset \varnothing \end{aligned}$ | RNDDIV DB | \$42,\$06 | , \$36, ø, ø |

XPABS - Absolute Value Function

| BøAE | XPABS |  |  |  |
| :--- | :--- | :---: | :--- | :--- |
| BØAE | 2ØF2AB | JSR | ARGPOP | ;GET ARGUMENT |
| BØB1 | A5D4 | LDA | FRØ | ;GET BYTE WITH SIGN |
| BØB3 | $297 F$ | AND | \#\$7F | ;AND OUT SIGN |
| BØB5 | 85D4 | STA | FRØ | ;SAVE |
| BØB7 | 4CBAAB | JMP | ARGPUSH | ;PUSH ON STACK |

XPUSR - USR Function

| BØBA |  | XPUSR |  |  |
| :---: | :---: | :---: | :---: | :---: |
| BØBA | 2øC3Bø | JSR | : USR | ; PUT RETURN ADDR IN CPU STACK |
| BØBD | 2ØAAD9 | JSR | CVIFP | ; CONVERT FRØ TO FP |
| BøCø | 4CBAAB | JMP | ARGPUSH | ; PUSH ON STACK |
|  |  | ; |  |  |
|  |  | ; |  |  |
|  |  | ; |  |  |
| ВØС3 |  | : USR |  |  |
| ВøС3 | A5Bø | LDA | COMCNT | ; GET COMMA COUNT |
| BøC5 | 85C6 | STA | ZTEMP2 | ; SET AS \# OF ARG FOR LOOP CONTROL |
| ВøC7 |  | : USR1 |  |  |
| ВØС7 | 20E3AB | JSR | GTINTO | ; GET AN INTEGER FROM OP STACK |
| BØCA | C6C6 | DEC | ZTEMP2 | ;DECR \# OF ARGUMENTS |
| ВøСС | $3 Ø \emptyset 9$ ^BØD7 | BMI | : USR2 | ; IF DONE THEM ALL, BRANCH |
|  |  | ; |  |  |
| BØCE | A5D4 | LDA | FRø | ; GET ARGUMENT LOW |
| BøDØ | 48 | PHA |  | ; PUSH ON STACK |
| BØD1 | A5D5 | LDA | FR®+1 | ; GET ARGUMENT HIGH |
| BØD3 | 48 | PHA |  | ; PUSH ON STACK |
| BøD4 | 4СС7Bø | JMP | : USRI | ;GET NEXT ARGUMENT |
| BØD7 |  | : USR2 |  |  |
| BøD7 | A5Bø | LDA | COMCNT | ; GET \# OF ARGUMENTS |
| BøD9 | 48 | PHA |  | ; PUSH ON CPU STACK |
| BøDA | 6CD4øø | JMP | [FRø] | ; GO TO USER ROUTINE |
| XPINT - Integer Function |  |  |  |  |
| BØDD |  | XPINT |  |  |
| BØDD | 2øF2AB | JSR | ARGPOP | ; GET NUMBER |
| BØEØ | 20E6BØ | JSR | XINT | ; GET INTEGER |
| BØE3 | 4CBAAB | JMP | ARGPUSH | ; PUSH ON ARGUMENT STACK |

## Source Code



## Transcendental Functions

XPSIN - Sine Function

| B11B |  |  |  |  |
| :--- | :--- | ---: | :--- | :--- |
| B11B | 2øF2AB | JPSIN | ARGPOP | ;GET ARGUMENT |
| B11E | 2øA7BD | JSR | SIN |  |
| B121 | Bø3F ${ }^{\wedge}$ B162 | BCS | :TBAD |  |
| B123 | $9 \emptyset 3 A{ }^{\wedge}$ B15F | BCC | :TGOOD |  |

## XPCOS - Cosine Function

| B125 | XPCOS |  |  |  |
| :--- | :--- | ---: | :--- | :--- |
| B125 | 2øF2AB | JSR | ARGPOP | ;GET ARGUMENT |
| B128 | 2øB1BD | JSR | COS |  |
| B12B | Bø35 AB162 | BCS | :TBAD |  |
| B12D | $9 \emptyset 3 \emptyset ~$ | AB15F | BCC | :TGOOD |

## XPATN - Arc Tangent Function

| B12F |  | XPATN |  |
| :--- | :--- | :---: | :--- |
| B12F | 2øF2AB | JSR | ARGPOP |
| B132 | $2 \emptyset 77 B E$ | JSR | ATAN |
| B135 | Bø2B | AB162 | BCS |
| B137 | $9 \varnothing 26$ | ^B15F | BCC |

## Source Code

## XPLOG - LOG Function

| B139 | XPLOG |  |  |
| :---: | :---: | :---: | :---: |
| B139 | 2øF2AB | JSR | ARGPOP |
| Bl3C | 2øCDDE | JSR | LOG |
| B13F | Bø21 *B162 | BCS | : TBAD |
| B141 | 901C *B15F | BCC | :TG |

XPL10 - LOG Base 10

| B143 | XPL1ø |  |  |
| :---: | :---: | :---: | :---: |
| B143 | 20F2AB | JSR | ARGPOP |
| B146 | 2øD1DE | JSR | LOG1ø |
| B149 | Bø17 ^B162 | BCS | :TBAD |
| B14B | 9012 ^B15F | BCC | : TGOOD |

## XPEXP - EXP Function

| B14D |  | XPEXP |  |
| :---: | :---: | :---: | :---: |
| B14D | 20F2AB | JSR | ARGPOP |
| B15ø | 2øCøDD | JSR | EXP |
| B153 | BøのD *B162 | BCS | : TBAD |
| B155 | 90ø8 *B15F | BCC | : TGOOD |

XPSQR - Square Root Function


## XPPOWER - Exponential Operator [A**B]

| B165 |  | XPPOWER |  |  |
| :---: | :---: | :---: | :---: | :---: |
| B165 | 20ø6AC | JSR | ARGP2 | ; GET ARGUMENT IN FRø,FRI |
| B168 | A5D4 | LDA | FRø | ; IS BASE $=\emptyset$ ? |
| B16A | DøøB ^B177 | BNE | : NØ | ; IF BASE NOT Ø, BRANCH |
| B16C | A5EØ | LDA | FR1 | ;TEST EXPONENT |
| B16E | Føб4 *B174 | BEQ | : Pø | ; IF $=\emptyset$; BRANCH |
| B170 | 10ED ^B15F | BPL | : TGOOD | ; $\mathrm{IF}>\emptyset$, ANSWER $=\varnothing$ |
| B172 | 30EE ^B162 | BMI | : TBAD | ; IF < $¢$, VALUE ERROR |
| B174 |  | : Pø |  |  |
| B174 | 4C05AD | JMP | XTRUE | ; $I F=\emptyset, ~ A N S W E R=1$ |
| B177 |  | : NØ |  |  |
| B177 | 1030 * B1A9 | ; BPL | : SPEVEN | ; IF BASE + THEN NO SPECIAL PROCESS |
| B179 | 297F | AND | \# ${ }^{\text {7 }}$ F | ; AND OUT SIGN BIT |
| B17B | 85D4 | STA | FRØ | ; SET AS BASE EXPONENT |
| B17D | A5EØ | LDA | FR1 | ; GET EXPONENT OF POWER |
| Bl7F | 297F | AND | \#\$7F | ; AND OUT SIGN BIT |
| B181 | 38 | SEC |  |  |
| B182 | E940 | SBC | \# \$4ø | ; IS POWER <1? |
| B184 | 30DC *B162 | BMI | : TBAD | ; IF YES, ERROR |
| B186 | A. 206 | LDX | \#6 | ; GET INDEX TO LAST DIGIT |
| B188 | C905 | CMP | \#5 | ; IF \# CAN HAVE DECIMAL |
| B18A | $9004{ }^{\text {* B190 }}$ | BCC | : SP4 | ; PORTION, THEN BR |
| B18C | Aøø1 | LDY | \#1 |  |
| B18E | DØø8 *B198 | BNE | : SP3 |  |
| B190 |  | : SP4 |  |  |
| B19ø | 85 F 5 | - STA | ZTEMP1 | ; SAVE EXP -4ø |

## Source Code

| B192 | 38 |  | SEC |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| B193 | A905 |  | LDA | \# 5 | ; GET \# BYTES POSSIBLE |
| B195 | E5F5 |  | SBC | ZTEMP1 | ; GET \# byTES THAT COULD BE DECIMAL |
| B197 | A8 |  | TAY |  | ; SET COUNTER |
|  |  |  | ; |  |  |
| B198 |  |  | : SP3 |  |  |
| B198 | CA |  | DEX |  |  |
| B199 | 88 |  | DEY |  | ; DEC COUNTER |
| B19A | Føø6 | * B1A2 | BEQ | : SP2 | ; IF DONE GO TEST EVEN/ODD |
| B19C | B5EØ |  | LDA | FR1, X | ; GET BYTE OF EXPONENT |
| B19E | DøC2 | ${ }^{\text {a }}$ B162 | BNE | :TBAD | ; IF NOT $=\emptyset$, THEN VALUE ERROR |
| BlAø | F6F6 | *B198 | BEQ | :SP3 | ; REPEAT |
|  |  |  | ; |  |  |
| B1A2 |  |  | : SP2 |  |  |
| BlA 2 | Aø8ø |  | LDY | \#\$8ø | ; GET ODD FLAG |
| BlA4 | B5Eø |  | LDA | FRI, X | ; GET BYTE OF EXPONENT |
| B1A6 |  |  | LSRA |  | ; IS IT ODD[LAST BIT OFF]? |
| Bla6 | +4A |  | LSR | A |  |
| B1A7 | ВØØ2 | * B1AB | BCS | : POWR | ; IF YES, BR |
|  |  |  | ; |  |  |
| B1A9 |  |  | : SPEVEN |  |  |
| B1A9 | Аøøø |  | LDY | \# $\varnothing$ |  |
| B1AB |  |  | : POWR |  |  |
| Blab | 98 |  | TYA |  |  |
| BlAC | 48 |  | PHA |  |  |

## Save Exponent [from FR1]

| Blad | A 205 | LDX | \#FMPREC | ;GET POINTER INTO FRI |
| :---: | :---: | :---: | :---: | :---: |
| BlaF |  | : POWR1 |  |  |
| BlAF | B5Eø | LDA | FR1, X | ; GET A BYTE |
| BlB1 | 48 | PHA |  | ; PUSH ON CPU STACK |
| B1B2 | CA | DEX |  | ; POINT TO NEXT BYTE |
| B1B3 | 1ØFA *BlAF | BPL | : POWRI | ; BR IF MORE TO DO |
| B1B5 | 2ØDlDE | ; JSR | LOG1ø | ;TAKE LOG OF BASE |
| B1B8 | BØA8 *B162 | BCS | : TBAD |  |

Pull Exponent into FR1 [from CPU Stack]


## Source Code

Statements
XDIM \& XCOM - Execute DIM and COMMON Statements

| $\begin{aligned} & \text { BlD9 } \\ & \text { BlD9 } \end{aligned}$ |  | XDIM |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | XCOM |  |  |  |
|  |  | ; |  |  |  |
| Bld9 | A4A8 | : DCl | LDY | STINDEX | ; IF NOT AT |
| BlDB | C4A7 | CPY |  | NXTSTD | ; STATEMENT END |
| BlDD | $9 \emptyset \emptyset 1$ * ${ }^{\text {BlEø }}$ | BCC |  | : DC2 | ; THEN CONTINUE |
| BlDF | 60 | RTS |  |  | ; RETURN |
| BlEø | 2øEØAA | : DC2 | JSR | EXEXPR | ; GO SET UP VIA EXECUTE EXPR |
| Ble3 | A5D2 | LDA |  | VTYPE | ; GET VAR TYPE |
| Ble5 |  | RORA |  |  | ; SHIFT DIM BIT TO CARRY |
| Ble5 | +6A | ROR |  | A |  |
| Ble6 | 9003 ^BlEB | BCC |  | : DC3 | ; CONTINUE IF NOT YET DIMMED |
| BlE8 | 2ø2EB9 | : DCERR | JSR | ERRDIM | ; ELSE ERROR |
|  |  | - DC3 |  |  |  |
| BlEB | 38 | : DC3 | SEC |  | ; TURN ON |
| BlEC |  | ROLA |  |  | ; DIM FLAG |
| BlEC | $+2 \mathrm{~A}$ | ROL |  | A |  |
| BlED | 85D2 | STA |  | VTYPE | ; AND RESET |
| B1EF | $3 \emptyset 2 \mathrm{~F}$ * B 220 | BMI |  | : DCSTR | ; AND BR IF STRING |
|  |  | ; |  |  |  |
| B1F1 | A 4 F 5 | LDY |  | ZTEMP1 | ; INC Il BY 1 |
| BlF3 | A6F6 | LDX |  | ZTEMP1+1 | ; AND SET AS DIM1 |
| BlF5 | C8 | INY |  |  |  |
| BlF6 | Døø3 * B1FB | BNE |  | : DC4 |  |
| B1F8 | E8 | INX |  |  |  |
| B1F9 | 30ED *BlE8 | BMI |  | : DCERR | ; BR IF OUT OF BOUNDS |
| BlFB | 84D6 | : DC4 | STY | VTYPE+EVAD1 |  |
| B1FD | 86D7 | STX |  | VTYPE+EVADI+1 |  |
| B1FF | 84F5 | STY |  | ZTEMP1 | ; ALSO PUT BACK ONTO |
| B201 | 86F6 | STX |  | ZTEMP1+1 | ; INDEX 1 FOR MULT |
|  |  | ; |  |  |  |
| B203 | A497 | LDY |  | INDEX2 | ; INC INDEX 2 BY 1 |
| B205 | A698 | LDX |  | INDEX2+1 | ; AND SET AS DIM 2 |
| B207 | C8 | INY |  |  |  |
| B208 | Døø3 * B 2 ¢D | BNE |  | : DC5 |  |
| B20A | E8 | INX |  |  |  |
| B20B | 30 DB * B1E8 | BMI |  | : DCERR | ; BR IF OUT OF BOUNDS |
| B20D | 84D8 | : DC5 | STY | VTYPE+EVAD2 |  |
| B2øF | 86D9 | STX |  | VTYPE+EVAD2+1 |  |
|  |  | ; |  |  |  |
| B211 | 205DAF | JSR |  | AMULI | ; ZTEMP1 = ZTEMP1*D2 |
| B214 | $2 \emptyset 46 \mathrm{AF}$ | JSR |  | AMUL2 | ; ZTEMP1 = ZTEMP1*6 |
|  |  | ; |  |  | RESULT IS AN ARRAY |
|  |  | ; |  |  | SPACE REQD |
| B217 | A4F5 | LDY |  | ZTEMP1 | ; $\mathrm{A}, \mathrm{Y}=\mathrm{LENGTH}$ |
| B219 | A5F6 | LDA |  | ZTEMP1+1 |  |
| B21B | 3øCB ^BlE8 | BMI |  | : DCERR |  |
| B21D | 4C34B2 | JMP |  | : DCEXP | ; GO EXPAND |
|  |  | ; |  |  |  |
| B22ø |  | : DCSTR |  |  |  |
| B22ø | A9øø | LDA |  | \# $\varnothing$ | ; SET CURRENT LENGTH $=\emptyset$ |
| B222 | 85D6 | STA |  | EVSLEN+VTYPE |  |
| B224 | 85D7 | STA |  | EVSLEN+1+VTYPE |  |
|  |  | ; |  |  |  |
| B226 | A 4 F 5 | LDY |  | ZTEMP1 | ; MOVE INDEX |
| B228 | 84D8 | STY |  | VTYPE+EVSDIM | ; TO STR DIM |
| B22A | A5F6 | LDA |  | ZTEMP1+1 | ; [ALSO LOAD A, Y] |
| B22C | 85D9 | STA |  | VTYPE+EVSDIM+1 | ; FOR EXPAND |
| B22E | Dø04 * ${ }^{\text {B234 }}$ | BNE |  | : DCEXP | ; INSURE DIM |
| B23ø | сøøø | CPY |  | \# $\varnothing$ | ; NOT ZERO |
| B232 | FØB4 *B1E8 | BEQ |  | : DCERR | ; FOR STRING |
| B234 B234 |  | : DCEXP |  |  |  |
| B234 | A.28E | LDX |  | \#ENDSTAR | ; POINT TO END ST \& ARRAY SPACE |
| B236 | 2081A8 | JSR |  | EXPAND | ; GO EXPAND |



XPOKE - Execute POKE

| B24C |  | XPOKE |
| :---: | :---: | :---: |
| B24C | 2ØEØAB | JSR |
| B24F | A5D4 | LDA |
| B251 | 8595 | STA |
| B253 | A5D5 | LDA |
| B255 | 8596 | STA |
|  |  | ; JSR |
| B257 | 20E9AB | JSR |
|  |  | ; |
| B25A | A5D4 | LDA |
| B25C | Aøøø | LDY |
| B25E | 9195 | STA |
| B260 | 60 | RTS |



XDEG - Execute DEG

| B261 | XDEG |  |  |  |  |
| :--- | :--- | ---: | :--- | :--- | :--- |
| B261 | A9ø6 | LDA | \#DEGON | GET DEGREES FLAG |  |
| B263 | $85 F B$ | STA | RADFLG | SET FOR TRANSCENDENTALS |  |
| B265 | $6 \emptyset$ | RTS |  |  |  |

XRAD - Execute RAD


XREST - Execute RESTORE Statement

| B26B |  | XREST |  |  | ; | ZERO DATA DISPL |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B26B | A900 | LDA |  | \# $\varnothing$ |  |  |
| B26D | 85B6 | STA |  | DATAD |  |  |
|  |  | ; |  |  |  |  |
| B26F | 2010B9 | - JSR |  | TSTEND | ; | TEST END OF STMT |
| B272 | 9øø3 * 2777 | BCC |  | : XR1 | ; | BR IF NOT END |
| B274 | A8 | TAY |  |  | ; | RESTORE TO LN=Ø |
| B275 | Føø7 *B27E | BEQ |  | : XR2 |  |  |
| B277 | 20D5AB | $\begin{aligned} & \text { : XRI } \end{aligned}$ | JSR | GETPINT | ; | GET LINE NO. |
|  |  | ; |  |  |  |  |
| B27A | A5D 5 | LDA |  | FRØ+1 | ; | LOAD LINE NO. |
| B27C | A4D4 | LDY |  | FRø |  |  |
| B27E | 85B8 | : XR2 | STA | DATALN+1 | ; | SET LINE |
| B28ø | 84B7 | STY |  | DATALN |  |  |
| B282 | $6 \emptyset$ | RTS |  |  | ; | DONE |

## XREAD - Execute READ Statement

| B283 |  | XREAD |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B283 | A5A8 | LDA | STINDEX | ; | SAVE | STINDEX |  |
| B285 | 48 | PHA |  |  |  |  |  |
| B286 | 20C7B6 | JSR | XGS | ; | SAVE | READ STMT | VIA GOSUB |
| B289 | A5B7 | ; LDA | DATALN | - | MOVE | N TO | NUM |
| B28B | 85Aø | STA | TSLNUM |  |  |  |  |
| B28D | A5B8 | LDA | DATALN+1 |  |  |  |  |
| B28F | 85Al | STA | TSLNUM+1 |  |  |  |  |

Source Code

| B291 | 20A2A9 | JSR |  | GETSTMT | ; | GO FIND TSLNUM |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B294 | A58A | LDA |  | STMCUR | ; | MOVE STMCUR TO INBUFF |
| B296 | 85F3 | STA |  | INBUFF |  |  |
| B298 | A58B | LDA |  | STMCUR+1 |  |  |
| B29A | 85F4 | STA |  | INBUFF+1 |  |  |
| B29C | 2ø19B7 | ; JSR |  | XRTN | ; | RESTORE READ STMT VIA RETURN |
| B29F | 68 | PLA |  |  | ' | RESTORE READ STMT VIA RETURN |
| B2Aø | 85A8 | STA |  | STINDEX | ; | SET IT |
| B2A2 |  | : XRD1 |  |  |  |  |
| B2A2 | Аøøø | LDY |  | \# $\varnothing$ | ; | SET CIX=ø |
| B2A4 | 84 F 2 | STY |  | CIX | ; | SET CIX |
| B2A6 | 2007B3 | JSR |  | : XRNT1 | ; | GET LINE NO. LOW |
| B2A9 | 85B7 | STA |  | DATALN | ; | SET LINE NO. LOW |
| B2AB | 2øø5B3 | JSR |  | : XRNT |  |  |
| B2AE | 85B8 | STA |  | DATALN+1 | ; | SET LINE NO. HIGH |
| B2Bø | 200583 | JSR |  | : XRNT |  |  |
| B2B3 | 85 F 5 | STA |  | ZTEMP1 | ; | SET LINE LENGTH |
| B2B5 |  | : XRD2 |  |  |  |  |
| B2B5 | 20ø5B3 | JSR |  | : XRNT |  |  |
| B2B8 | 85 F 6 | STA |  | ZTEMP1+1 | ; | SET STMT LENGTH |
|  |  | ; |  |  |  |  |
| B2BA | 2005B3 | JSR |  | : XRNT | ; | GET STMT LINE TOKEN |
| B2BD | C901 | CMP |  | \#CDATA | ; | IS IT DATA |
| B2BF | F626 *B2E7 | BEQ |  | : XRD4 | ; | BR IF DATA |
|  |  | ; |  |  |  |  |
| B2Cl | A 4 F 6 | LDY |  | ZTEMP1+1 | ; | GET DISPL TO NEXT STMT |
| B2C3 | C4F5 | CPY |  | ZTEMP1 | ; | IS IT EOL |
| B2C5 | Bøø5 * B2CC | BCS |  | : XRD2A | ; | BR IF EOL |
| B2C7 | 88 | DEY |  |  |  |  |
| B2C8 | 84F2 | STY |  | CIX | ; | SET NEW DISPL |
| B2CA | 9ØE9 * B2B5 | BCC |  | : XRD2 | ; | AND CONTINUE THIS STMT |
| B2CC | 84 F 2 | ; XRD2A | STY | CIX |  |  |
| B2CE | C6F2 | DEC |  | CIX |  |  |
| B2Dø |  | ; ${ }^{\text {: XRD3 }}$ | LDY | \#1 | ; |  |
| B2D2 | B1F3 | LDA |  | [INBUFF], Y | , | DIRECT ONE |
| B2D4 | 303D * 313 | BMI |  | : XROOD | ; | BR IF IT WAS [OUT OF DATA] |
|  |  | ; |  |  |  |  |
| B2D6 | 38 | SEC |  |  |  |  |
| B2D7 | A5F2 | LDA |  | CIX | ; | INBUFF + CIX + 1 |
| B2D9 | 65F3 | ADC |  | INBUFF | ; | $=$ ADR NEXT PGM LINE |
| B2DB | 85 F 3 | STA |  | INBUFF |  |  |
| B2DD | A9øø | LDA |  | \#ø |  |  |
| B2DF | 85B6 | STA |  | DATAD |  |  |
| B2El | 65F4 | ADC |  | INBUFF+1 |  |  |
| B2E3 | 85F4 | STA |  | INBUFF+1 |  |  |
| B2E5 | 9ØBB * B2A 2 | BCC |  | : XRD1 | ; | GO SCAN THIS NEXT LINE |
|  |  | ; |  |  |  |  |
| B2E7 |  | : XRD4 |  |  |  |  |
| B2E7 | A9øø | LDA |  | \#Ø | ; | CLEAR ELEMENT COUNT |
| B2E9 | 85F5 | STA |  | ZTEMP1 |  |  |
|  |  | ; |  |  |  |  |
| B2EB |  | : XRD5 |  |  |  |  |
| B2EB | A5F5 | LDA |  | ZTEMP1 | ; | GET ELEMENT COUNT |
| B2ED | C5B6 | CMP |  | DATAD | ; | AT PROPER ELEMENT |
| B2EF | BøØB * B 2 FC | BCS |  | : XRD 7 | ; | BR IF AT |
|  |  | ; |  |  |  | ELSE SCAN FOR NEXT |
| B2F1 | 2005B3 | : XRD6 | JSR | : XRNT | ; | GET CHAR |
| B2F4 | DØFB * B 2 Fl | BNE |  | : XRD6 | ; | BR IF NOT CR OR COMMA |
| B2F6 | BØD8 ^B2DØ | BCS |  | : XRD3 | ; | BR IF CR |
| B2F8 | E6F5 | INC |  | ZTEMP1 | ; | INC ELEMENT COUNT |
| B2FA | DØEF *B2EB | BNE |  | : XRD5 | ; | AND GO NEXT |
|  |  | ; |  |  |  |  |
| B2FC | A940 | : XRD7 | LDA | \#\$40 | ; | SET READ BIT |
| B2FE | 85A6 | STA |  | DIRFLG |  |  |
| В3øø | E6F2 | INC |  | CIX | ; | INC OVER DATA TOKEN |

## Source Code

| B3ø2 | 4C35B3 | JMP |  | : XINA | ; GO DO IT |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | ; |  |  |  |  |  |
|  |  | ; |  |  |  |  |  |
| B305 |  | : XRNT |  |  |  |  |  |
| B305 | E6F2 | INC |  | CIX | ; | INC INDEX |  |
| В307 | A 4 F 2 | : XRNTI | LDY | CIX | ; | GET INDEX |  |
| В309 | B1F3 | LDA |  | [INBUFF], Y | ; | GET CHAR COUNT |  |
| B30В | C92C | CMP |  | \# \$2C | ; | IS IT A COMMA |  |
| B30D | 18 | CLC |  |  | ; | CARRY CLEAR FOR | COMMA |
| B30E | Føø2 *B312 | BEQ |  | : XRNT2 | ; | BR IF COMMA |  |
| B31ø | C99B | CMP |  | \#CR | ; | IS IT CR |  |
| B312 | 60 | : XRNT2 | RTS |  |  |  |  |
|  |  | , XROOD |  |  |  |  |  |
| B313 | 2ø34B9 | : XROOD | JSR | ERROOD |  |  |  |

XINPUT - Execute INPUT


Source Code

|  |  | ; |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| B378 | A4F5 | : XIS 2 | LDY | ZTEMP1 | ; GET SAVED INDEX |
| B37A | A5A8 | LDA |  | STINDEX | ; SAVE INDEX |
| B37C | 48 | PHA |  |  |  |
| B37D | 8A | TXA |  |  | ; ACU $=$ CHAR COUNT |
| B37E | A2F3 | LDX |  | \#INBUFF | ; POINT TO INBUFF |
| B380 | 2664 AB | JSR |  | RISC | ; GO MAKE STR VAR |
| B383 | 68 | PLA |  |  |  |
| B384 | 85A8 | STA |  | STINDEX | ; RESTORE INDEX |
| B386 | 20A6AE | JSR |  | RISASN | ; THEN DO STA ASSIGN |
|  |  | ; |  |  |  |
| B389 | 24A6 | : XINX | BIT | DIRFLG | ; IS THIS READ |
| B38B | 50øF *B39C | BVC |  | : XIN | ; BR IF NOT |
|  |  | ; |  |  |  |
| B38D | E6B6 | INC |  | DATAD | ; INC DATA DISPL |
| B38F | 2010В9 | JSR |  | TSTEND | ; TEST END READ STMT |
| B392 | BøøD *B3A1 | BCS |  | : XIRTS | ; BR IF READ END |
| B394 | 2øø7B3 | ; XIRI | JSR | : XRNTl | ; GET END DATA CHAR |
| B397 | $9 \emptyset 18{ }^{\text {^ }}$ B3B1 | BCC |  | : XINC | ; BR IF COMMA |
| B399 | $4 \mathrm{CDø日2}$ | JMP |  | : XRD3 | ; GO GET NEXT DATA LINE |
|  |  | ; |  |  |  |
| B39C |  | : XIN |  |  |  |
| B39C | 2ø10B9 | JSR |  | TSTEND |  |
| B39F | 9øø8 ^B3A9 | BCC |  | : XIN1 |  |
|  |  | ; |  |  |  |
| B3A1 | 2051 DA | : XIRTS | JSR | INTLBF | ; RESTORE LBUFF |
| B3A4 | A9øø | LDA |  | \# $\varnothing$ | ; RESTORE ENTER |
| B3A6 | 85B4 | STA |  | ENTDTD | ; DEVICE TO ZERO |
| B3A8 | 60 | RTS |  |  | ; DONE |
|  |  | , |  |  |  |
| B3A9 | 2øø7B3 | : XIN 1 | JSR | : XRNT1 | ; IF NOT END OF DATA |
| B3AC | $90 \emptyset 3$ ~B3B1 | BCC |  | : XINC | ; THEN BRANCH |
| B3AE | 4C26B3 | JMP |  | :XINØ | ; AND CONTINUE |
|  |  | ; |  |  |  |
| B3B1 | E6F2 | : XINC | INC | CIX | ; INC INDEX |
| В3B3 | 4C35B3 | JMP |  | : XINA | ; AND CONTINUE |

XPRINT - Execute PRINT Statement



Source Code
B461 4C9FBA JMP PRCHAR OUTPUT CHAR

## XLPRINT - Print to Printer



## XLIST - Execute LIST Command

| B483 |  | XLIST |  |  |
| :---: | :---: | :---: | :---: | :---: |
| B483 | Аøøø | LDY | \# $\varnothing$ | ; SET TABLE SEARCH LINE NO |
| B485 | 84A】 | STY | TSLNUM | ; TO ZERO |
| B487 | 84A1 | STY | TSLNUM+1 |  |
| B489 | 88 | DEY |  |  |
| B48A | 84AD | STY | LELNUM | ; SET LIST END LINE NO |
| B48C | A97F | LDA | \# \$7F | ;TO \$7FFF |
| B48E | 85AE | STA | LELNUM+1 |  |
| B490 | 8DFEØ2 | STA | \$2FE | ; SET NON-DISPLAY MODE |
| B493 | A99B | LDA | \#CR | ; POINT CR |
| B495 | 209FBA | JSR | PRCHAR |  |
| B498 | 20С7B6 | ; JSR | XGS | SAVE CURLINE VIA GOSUB |
| B49B |  | : XLØ |  |  |
| B49B | A4A8 | LDY | STINDEX | ; GET STMT INDEX |
| B49D | C8 | INY |  | ; INC TO NEXT CHAR |
| B49E | C4A7 | CPY | NXTSTD | ; RT NEXT STMT |
| B4Aø | Bø2D * ${ }^{\text {a }}$ 4CF | BCS | : LSTART | ; BR IF AT, NO PARMS |
|  |  | ; |  |  |
| B4A2 | A5A8 | LDA | STINDEX | ; SAVE STINDEX |
| B4A4 | 48 | PHA |  | ; ON STACK |
| B4A5 | 2ø日FAC | JSR | POP1 | ; POP FIRST ARGUMENT |
| B4A8 | 68 | PLA |  | ; RESTORE STINDEX TO |
| B4A9 | 85A8 | STA | STINDEX | ; RE-DO FIRST ARG |
| B4AB | A5D2 | LDA | VTYPE | ; GET VAR TYPE |
| B4AD | $10 \varnothing 6$ ^B4B5 | BPL | : XLI | ; BR IF NOT FILE SPEC STRING |
| B4AF | 2ØD5BA | JSR | FLIST | ; GO OPEN FILE |
| B4B2 | 4C9BB4 | JMP | : XLØ | ; GO BACK TO AS IF FIRST PARM |
|  |  | ; |  |  |
| B4B5 |  | : XLI |  |  |
| B4B5 | 20D5AB | JSR | GETPINT | ; GO GET START LNO |
|  |  | ; STA |  |  |
| B4B8 | 85A1 | STA | TSLNUM+1 |  |
| B4BA | A5D4 | LDA | FRØ | ; MOVE START LNO |
| B4BC | 85AØ | STA | TSLNUM | ; TO TSLNUM |
|  |  | ; |  |  |
| B4BE | A4A8 | LDY | STINDEX | ;GET STMT INDEX |
| B4Cø | C4A7 | CPY | NXTSTD | ; AT NEXT STMT |
| B4C2 | FØØ3 ^B4C7 | BEQ | : LSE | ; BR IF AT, NO PARMS |


| B4C4 | 2øD5AB | JSR |  | GETPINT | ; GO GET LINE NO |
| :---: | :---: | :---: | :---: | :---: | :---: |
| B4C7 | A5D4 | : LSE | LDA | FRø | ; MOVE END LINE NO |
| B4C9 | 85AD | STA |  | LELNUM | ; TO LIST END LINE NO |
| B4CB | A5D 5 | LDA |  | FRø+1 | ; |
| B4CD | 85AE | STA |  | LELNUM+1 |  |
|  |  | ; |  |  |  |
|  |  | ; |  |  |  |
| B4CF |  | : LSTART |  |  |  |
| B4CF | 20A2A9 | JSR |  | GETSTMT | ; GO FIND FIRST LINE |
| B4D2 | 2øE2A9 | : LNXT | JSR | TENDST | ; AT END OF STMTS |
| B4D 5 | $3 \varnothing 24$ "B4FB | BMI |  | : LRTN | ; BR AT END |
| B4D7 | Aøø1 | ':LTERNG | LDY | \#1 | ; COMPARE CURRENT STMT |
| B4D9 | B18A | LDA |  | [STMCUR], Y | ;LINE NO WITH END |
| B4DB | C5AE | CMP |  | LELNUM+1 | ;LINE NO |
| B4DD | 9øøВ *B4EA | BCC |  | :LGO |  |
| B4DF | Dø1A *B4FB | BNE |  | : LRTN |  |
| B4E1 | 88 | DEY |  |  |  |
| B4E2 | B18A | LDA |  | [STMCUR], Y |  |
| B4E4 | C5AD | CMP |  | Lelnum |  |
| B4E6 | $9 \varnothing ø 2$ *B4EA | BCC |  | : LGO |  |
| B4E8 | Dø11 *B4FB | BNE |  | : LRTN |  |
| B4EA | 2ø5CB5 | ; LGO | JSR | :LLINE | ;GO LIST THE LINE |
| B4ED | 20F4A9 | JSR |  | TSTBRK | ; TEST FOR BREAK |
| B4Fø | Døø9 ^B4FB | BNE |  | : LRTN | ; BR IF BREAK |
| B4F2 | 2øDDA9 | JSR |  | GETLL |  |
| B4F5 | 2øDøA9 | JSR |  | GNXTL | ;GO INC TO NEXT LINE |
| B4F8 | 4CD2B4 | JMP |  | : LNXT | ;GO DO THIS LINE |
|  |  | ; |  |  |  |
| B4FB |  | : LRTN |  |  |  |
| B4FB | A5B5 | LDA |  | LISTDTD | ; IF LIST DEVICE |
| B4FD | Føø7 * F 506 | BEQ |  | : LRTN1 | ; IS ZERO, BR |
| B4FF | 20F1BC | JSR |  | CLSYSD | ; ELSE CLOSE DEVICE |
| B502 | A900 | LDA |  | \# 0 | ; AND RESET |
| B504 | 85B5 | STA |  | LISTDTD | ; DEVICE TO zero |
| B5ø6 |  | :LRTN1 |  |  |  |
| B596 | 8DFEø2 | STA |  | \$2FE | ; SET DISPLAY MODE |
| B509 | 4 Cl 19 B 7 | JMP |  | XRTN | ; THEN RESTORE LIST LINE |

## LSCAN - Scan a Table for LIST Token

|  |  | ; | ENTRY PARMS |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | ; | $\mathrm{X}=$ SKIP LENGTH |  |  |
|  |  | ; | A, $\mathrm{Y}=\mathrm{TABLE}$ ADR |  |  |
|  |  | ; | SCANT $=$ TOKEN |  |  |
|  |  | ; |  |  |  |
| B5øC |  | : LSCAN |  |  |  |
| B50C | 86AA | STX |  | SRCSKP | ; SAVE SKIP LENGTH |
| B5øE | 2ø3Ø日5 | JSR |  | : LSST | ; SAVE SRC ADR |
| B511 | A4AA | ; LSCØ | LDY | SRCSKP | ; GET SKIP FACTOR |
|  |  | ; |  |  |  |
| B513 | C6AF | DEC |  | SCANT | ; DECREMENT SRC COUNT |
| B515 | 3øøE ^B525 | BMI |  | : LS INC | ; BR IF DONE |
|  |  | ; |  |  |  |
| B517 | B195 | : LSCl | LDA | [SRCADR], Y | ; GET CHARACTER |
| B519 | $3 ø \emptyset 3$ *B51E | BMI |  | : LSC2 | ; BR IF LAST CHARACTER |
| B51B | C8 | INY |  |  | ; INC TO NEXT |
| B51C | DØF9 *B517 | BNE |  | : LSCl | ; BR ALWAYS |
| B51E | C8 | : LSC2 | INY |  | ; INC TO AFTER LAST CHAR |
| B51F | 2025B5 | JSR |  | : LSINC | ; INC SRC ADR BY Y |
| B522 | 4C11B5 | JMP |  | : LSCの | ; GO TRY NEXT |
|  |  | ; |  |  |  |
| B525 | 18 | : LSINC | CLC |  |  |
| B526 | 98 | TYA |  |  | ; Y PLUS |
| B527 | 6595 | ADC |  | SRCADR | ; SRCADR |
| B529 | 8595 | STA |  | SRCADR | ; IS |

## Source Code

| B52B | A8 | TAY |  |  | ; | NEW |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B52C | A596 | LDA |  | SRCADR+1 | ; | SRCADR |  |
| B52E | 69øø | ADC |  | \#Ø |  |  |  |
| B530 | 8596 | : LSST | STA | SRCADR+1 | ; | Store new | SRCADR |
| B532 | 8495 | STY |  | SRCADR | ; | AND |  |
| B534 | 60 | RTS |  |  |  | RETURN |  |

## LPRTOKEN - Print a Token

| B535 |  | LPRTOKEN |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B535 |  | : LPRTOKEN |  |  |  |  |
| B535 | AØFF | LDY | \# \$FF | ; | INITIALIZE INDEX TO | ZERO |
| B537 | 84AF | STY | SCANT |  |  |  |
| B539 | E6AF | :LPT1 INC | SCANT | ; | INC INDEX |  |
| B53B | A4AF | LDY | SCANT | ; | GET INDEX |  |
| B53D | B195 | LDA | [SRCADR], Y | ; | GET TOKEN CHAR |  |
| B53F | 48 | PHA |  | ; | SAVE CHAR |  |
| B54ø | C99B | CMP | \#CR | ; | IF ATARI CR |  |
| B542 | Føø4 *B548 | BEQ | : LPT1A | ; | THEN DON'T AND |  |
| B544 | 297 F | AND | \#\$7F | ; | TURN OFF MSB |  |
| B546 | Fø03 *B54B | BEQ | : LPT2 | ; | BR IF NON-PRINTING |  |
| B548 |  | : LPTIA |  |  |  |  |
| B548 | 209FBA | JSR | PRCHAR | ; | GO PRINT CHAR |  |
| B54B |  | : LPT2 |  |  |  |  |
| B54B | 68 | PLA |  | ; | GET CHAR |  |
| B54C | 10EB ^B539 | BPL | : LPT1 | ; | BR IF NOT END CHAR |  |
| B54E | 60 | RTS |  | ; | GO BACK TO MY BOSS |  |

## LPTWB - Print Token with Blank Before and After

| B54F |  | : LPTWB |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B54F | A920 | LDA |  | \#\$20 | ; | GET BLANK |  |
| B551 | 209FBA | JSR |  | PRCHAR | ; | GO PRINT IT |  |
| B554 | 2035B5 | : LPTTB | JSR | : LPRTOKEN | ; | GO PRINT TOKEN |  |
| B557 | A920 | : LPBLNK | LDA | \# \$2ø | ; | GET BLANK |  |
| B559 | 4C9FBA | JMP |  | PRCHAR | ; | GO PRINT IT AND | RETURN |
|  |  | ; |  |  |  |  |  |
|  |  | ; |  |  |  |  |  |
|  |  | ; |  |  |  |  |  |

## LLINE - List a Line

| B55C |  | LLINE |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B55C |  | : LLINE |  |  |  |  |
| B55C | Аøøø | LDY | \# $\square$ |  |  |  |
| B55E | B18A | LDA | [STMCUR], Y | ; | MOVE LINE NO |  |
| B560 | 85D4 | STA | FRØ | ; | TO FRø |  |
| B562 | C8 | INY |  |  |  |  |
| B563 | B18A | LDA | [STMCUR], Y |  |  |  |
| B565 | 85D5 | STA | FRø+1 |  |  |  |
| B567 | 2ØAAD9 | JSR | CVIFP | ; | CONVERT TO FP |  |
| B56A | 20E6D8 | JSR | CVFASC | ; | CONVERT TO ASCII |  |
| B56D | A5F3 | LDA | INBUFF | ; | MOVE INBUFF ADR |  |
| B56F | 8595 | STA | SRCADR | ; | TO SRCADR |  |
| B571 | A5F4 | LDA | INBUFF+1 |  |  |  |
| B573 | 8596 | STA | SRCADR+1 |  |  |  |
| B575 | 2054B5 | JSR | : LPTTB | ; | AND PRINT LINE NO |  |
|  |  | ; |  |  |  |  |
| B578 |  | LDLINE |  |  |  |  |
| B578 | Aøø2 | LDY | \#2 |  |  |  |
| B57A | B18A | LDA | [STMCUR], Y | ; | GET LINE LENGTH |  |
| B57C | 859F | STA | LLNGTH | ; | AND SAVE |  |
| B57E | C8 | INY |  |  |  |  |
| B57F | B18A | :LL1 LDA | [STMCUR], Y |  | ; GET STMT LENGTH |  |
| B581 | 85A7 | STA | NXTSTD | ; | AND SAVE AS NEXT ST | DISPL |
| B583 | C8 | INY |  | ; | INC TO STMT TYPE |  |
| B584 | 84A8 | STY | STINDEX | ; | AND SAVE DISPL |  |
| B586 | 2090B5 | JSR | : LSTMT |  | GO LIST STMT |  |


| B589 | A4A7 | LDY | NXTSTD | ; DONE LINE |
| :--- | :--- | :--- | :--- | :--- |
| B58B | C49F | CPY | LLNGTH |  |
| B58D | $9 \emptyset F \emptyset$ | ${ }^{\wedge}$ B57F | BCC | LLL1 |
| B58F | $6 \emptyset$ |  | RTS |  |

## LSTMT - List a Statement

| B590 |  | : LSTMT |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B590 | 2031B6 | JSR |  | : LGCT | ; | GET CURRENT TOKEN |
| B593 | C936 | CMP |  | \#CILET | ; | IF IMP LET |
| B595 | FØ17 *B5AE | BEQ |  | : LADV | ; | BR |
| B597 | 203 DB6 | JSR |  | LSTMC | ; | GO LIST STMT CODE |
|  |  | ; JSR |  |  |  |  |
| B59A | 2031B6 | JSR |  | : LGCT | ; | GO GET CURRENT TOKEN |
| B59D | C937 | CMP |  | \#CERR | ; | BR IF ERROR STMT |
| B59F | Føø4 ^B5A5 | BEQ |  | : LDR |  |  |
| B5A1 | C902 | CMP |  | \#2 | ; | WAS IT DATA OR REM |
| B5A3 | Bøø9 ^B5AE | BCS |  | : LADV | ; | BR IF NOT |
|  |  | ; |  |  |  |  |
| B5A5 | 2ø2FB6 | : LDR | JSR | : LGNT | , | OUTPUT DATA/REM |
| B5A8 | 269FBA | JSR |  | PRCHAR | ; | THEN PRINT THE CR |
| B5AB | 4CA5B5 | JMP |  | : LDR |  |  |
|  |  | ; |  |  |  |  |
| B5AE | $202 \mathrm{FB6}$ | : LADV | JSR | : LGNT | ; | GET NEXT TOKEN |
| B5B1 | 101A *B5CD | BPL |  | : LNVAR | ; | BR IF NOT VARIABLE |
|  |  | ; |  |  |  |  |
| B5B3 | 297F | AND |  | \# \$ 7F | ; | TURN OFF MSB |
| B5B5 | 85AF | STA |  | SCANT | ; | AND SET AS SCAN COUNT |
| B5B7 | A20ø | LDX |  | \# $\varnothing$ | ; | SCAN VNT FOR |
| B5B9 | A583 | LDA |  | VNTP+1 | ; | VAR NAME |
| B5BB | A482 | LDY |  | VNTP |  |  |
| B5BD | $2 \emptyset \emptyset C B 5$ | JSR |  | : LSCAN | ; |  |
| B5Cø | 2035B5 | :LSI | JSR | : LPRTOKEN | ; | PRINT VAR NAME |
| B5C3 | C9A8 | CMP |  | \# \$A8 | , | NAME END IN LPAREN |
| B5C5 | DØE7 ^B5AE | BNE |  | : LADV | ; | BR IF NOT |
| B5C7 | 2ø2FB6 | JSR |  | : LGNT | ; | DON'T PRINT NEXT TOKEN |
| B5CA | 4CAEB5 | JMP |  | : LADV | ; | IF IT IS A PAREN |
|  |  | ; |  |  |  |  |
| B5CD |  | : LNVAR |  |  |  |  |
| B5CD | C90F | CMP |  | \# ${ }^{\text {¢ }}$ ¢ | ; | TOKEN: \$øF |
| B5CF | F018 *B5E9 | BEQ |  | : LSTC | ; | BR IF ØF, STR CONST |
| B5D1 | B036 *B609 | ; BCS |  | : LOP | , | BR IF TOKEN > $\$ \emptyset \mathrm{~F}$ |
|  |  | ; JSR |  |  |  | ELSE IT'S NUM CONST |
| B5D3 | 2ø4DAB | JSR |  | NCTOFRØ | ; | GO MOVE FRØ |
| B5D6 | C6A8 | DEC |  | STINDEX | ; | BACK INDEX TO LAST CHAR |
| B5D8 | 20E6D8 | JSR |  | CVFASC | ; | CONVERT FRØ TO ASCII |
| B5DB | A5F3 | LDA |  | INBUFF | ; | POINT SCRADR |
| B5DD | 8595 | STA |  | SRCADR | ; | TO INBUFF WHERE |
| B5DF | A5F4 | LDA |  | INBUFF+1 | ; | CHAR IS |
| B5E1 | 8596 | STA |  | SRCADR+1 | ; |  |
| B5E3 | 2ø35B5 | : LSX | JSR | : LPRTOKEN | , | GO PRINT NUMBER |
| B5E6 | 4CAEB5 | JMP |  | : LADV | ; | GO FOR NEXT TOKEN |
| B5E9 | 202FB6 | ; LSTC | JSR | : LGNT | ; | GET NEXT TOKEN |
| B5EC | 85AF | STA |  | SCANT | ; | WHICH IS STR LENGTH |
| B5EE | A922 | LDA |  | \#\$22 | ; | PRINT DOUBLE QUOTE CHAR |
| B5Fø | 209FBA | JSR |  | PRCHAR |  |  |
| B5F3 | A5AF | LDA |  | SCANT |  |  |
| B5F5 | FのØA * B6Ø1 | BEQ |  | : LS3 |  |  |
|  |  | ; |  |  |  |  |
| B5F7 | $202 \mathrm{FB6}$ | : LS 2 | JSR | : LGNT | ; | OUTPUT STR CONST |
| B5FA | 2ø9FBA | JSR |  | PRCHAR | ; | CHAR BY CHAR |
| B5FD | C6AF | DEC |  | SCANT | ; | UNTIL COUNT $=\emptyset$ |
| B5FF | DøF6 ^B5F7 | BNE |  | : LS 2 |  |  |
| B601 |  | ${ }^{\text {i }}$ LS3 |  |  |  |  |
| B601 | A922 | LDA |  | \#\$22 | ; | THEN OUTPUT CLOSING |
| B603 | 209FBA | JSR |  | PRCHAR | ; | DOUBLE QUOTE |
| B606 | 4CAEB5 | JMP |  | : LADV |  |  |

Source Code


XFOR - Execute FOR


## Source Code



## XGOSUB - Execute GOSUB

| B6A | XGOSUB |
| :--- | ---: |
| B6A $\varnothing$ | $2 \emptyset C 7 B 6$ |$\quad$ JSR $\quad$ XGS GO TO XGS ROUTINE

XGOTO - Execute GOTO

| $\begin{aligned} & \text { B6A } 3 \\ & \text { B6A } 3 \end{aligned}$ | 2ØD5AB | XGOTO |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | JSR | GETPINT |  | ; GET | POSTIVE | INTEGER IN | IN FRø |
|  |  | ; | GET LINE ADRS \& POINTERS |  |  |  |  |  |
|  |  | ; |  |  |  |  |  |  |
|  |  | ; |  |  |  |  |  |  |
| B6A6 |  | XGO2 |  |  |  |  |  |  |
| B6A6 | A5D5 | LDA | FRø+1 |  | ; X |  |  |  |
| B6A8 | 85A1 | STA | TSLNUM+1 |  | ; X |  |  |  |
| B6AA | A5D4 | LDA | FRø |  | ; PUT | LINE \# | IN TSLNUM |  |
| B6AC | 85Aø | STA | TSLNUM |  | ; X |  |  |  |

## Source Code

| B6AE |  | ; ${ }^{\text {XGO1 }}$ |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B6AE | 20A2A9 | JSR | GETSTMT | ; | LINE POINTERS AND STMT | ADDRESS |
| B6Bl |  | BCS | : ERLN |  | IF NOT FOUND ERROR |  |
| B6B3 | 68 | PLA |  |  | CLEAN UP STACK |  |
| B6B4 | 68 | PLA |  |  |  |  |
| B6B5 | 4C5FA9 | JMP | EXECNL |  | GO TO EXECUTE CONTROL |  |
|  |  | ; |  |  |  |  |
| B6B8 |  | : ERLN |  |  |  |  |
| B6B8 | $20 \mathrm{BEB6}$ | JSR | RESCUR | ; | RESTORE STMT CURRENT |  |
|  |  | ; |  |  |  |  |
|  |  | ; |  |  |  |  |
|  |  | ; |  |  |  |  |
| B6BB | 2028B9 | JSR | ERNOLN | ; | LINE \# NOT FOUND |  |
| B6BE |  | RESCUR |  |  |  |  |
| B6BE | A5BE | LDA | SAVCUR | ; | RESTORE STMCUR |  |
| B6C0 | 858A | STA | STMCUR | ; | X |  |
| B6C2 | A5BF | LDA | SAVCUR+1 | ; | X |  |
| B6C4 | 858B | STA | STMCUR+1 | ; | X |  |
| B6C6 | 60 | RTS |  |  |  |  |

XGS - Perform GOSUB [GOSUB, LIST, READ]

| B6C7 |  | XGS |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B6C7 | 2ø8AB8 | JSR | :SAVDEX | ; | GET | STMT | INDEX |  |  |
| B6CA |  | XGS1 |  |  |  |  |  |  |  |
| B6CA | A9øø | LDA | \# $\varnothing$ | ; | GET | GOSU | TYPE |  |  |
| B6CC | 4C83B6 | JMP | PSHRSTK | ; | PUT | ELEM | NT ON | RUN | STACK |

XNEXT — Execute NEXT


## Source Code




## XEND - Execute END

| B78D | XEND |  |  |
| :---: | :---: | :---: | :---: |
| B78D | 2øA7B7 | JSR | STOP |
| B790 | 4C5@AD | JMP | SNXI |

## XSTOP - Execute STOP

| $\begin{aligned} & \text { B793 } \\ & \text { B793 } \end{aligned}$ |  | XSTOP |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | 2ØA7B7 | JSR | STOP | ; | GO SET UP STOP LINE \# |
|  |  | ; |  |  |  |
|  |  | ; | PRINT MESSAGE |  |  |
|  |  | ; |  |  |  |
| B796 | $2 \emptyset 6 \mathrm{EBD}$ | JSR | PRCR | ; | PRINT CR |
| B799 | A9B6 | LDA | \# : MSTOP\&255 | ; | SET POINTER FOR MESSAGE |
| B79B | 8595 | STA | SRCADR | ; | X |
| B79D | A9B7 | LDA | \# : MSTOP/256 | ; | X |
| B79F | 8596 | STA | SRCADR+1 | ; | X |
|  |  | ; |  |  |  |
| B7A1 | 2035B5 | JSR | LPRTOKEN | ; | PRINT IT |
| B7A4 | 4C74B9 | ; JMP | : ERRM2 | ; | PRINT REST OF MESSAGE |
|  |  | ; |  |  |  |
|  |  | ; |  |  |  |
|  |  | ; |  |  |  |
| B7A7 |  | STOP |  |  |  |
| B7A7 | 20E2A9 | JSR | TENDST | ; | GET CURRENT LINE \# HIGH |
| B7AA | $3 \varnothing \emptyset 7{ }^{\text {^ }}$ В7B3 | BMI | : STOPEND | ; | IF -, THIS IS DIRECT STMT |
|  |  | ; |  | ; | DON'T STOP |
| B7AC | 85BB | STA | STOPLN+1 | ; | SAVE LINE \# HIGH FOR CON |
| B7AE | 88 | DEY |  | ; | DEC INDEX |
| B7AF | B18A | LDA | [STMCUR], Y | ; | GET LINE \# LOW |
| B7B1 | 85BA | STA | STOPLN | ; | SAVE FOR CON |
| B7B3 |  | : STOPEND |  |  |  |
| B7B3 | 4C72BD | JMP | SETDZ | ; | SET L/D DEVICE $=\varnothing$ |
|  |  | ; |  |  |  |
|  |  | ; |  |  |  |
|  |  | ; |  |  |  |
| B7B6 | $\begin{aligned} & 53544 \mathrm{~F} 5 \emptyset 5 \emptyset \\ & 4544 \mathrm{~A} \end{aligned}$ | :MSTOP | DC 'STOPPED |  |  |

XCONT - Execute Continue

| B7BE |  | XCONT |  |  |
| :---: | :---: | :---: | :---: | :---: |
| B7BE | 2øE2A9 | JSR | TENDST | ; IS IT INDIRECT STMT? |
| B7Cl | 1øFØ ^B7B3 | BPL | : STOPEND | ; IF YES, BR |
| B7C3 | A5BA | LDA | STOPLN | ; SET STOP LINE \# AS LINE \# FOR GET |
| B7C5 | 85Aø | STA | TSLNUM | ; X |
| B7C7 | A5BB | LDA | STOPLN+1 | ; X |
| B7C9 | 85Al | STA | TSLNUM+1 | ; X |
| B7CB | 20A2A9 | ; JSR | GETSTMT | ; GET ADR OF STMT WE STOPPED AT |
| B7CE | 2ØE2A9 | JSR | TENDST | ; AT END OF STMT TAB ? |
| B7D1 | 36A2 *B775 | BMI | : RUNEND |  |
| B7D3 | 2ØDDA9 | JSR | GETLL | ; GET NEXT LINE ADDR IN CURSTM |
| B7D6 | 2ØDØA9 | JSR | GNXTL | ; X |
| B7D9 | 2ØE2A9 | JSR | TENDST | ; SEE IF WE ARE AT END OF STMT TABLE |
| B7DC | 3097 ~B775 | BMI | : RUNEND | ; BR IF MINUS |
| B7DE | 4C1BB8 | JMP | SETLN1 | ; SET UP LINE POINTERS |
| XTRAP - Execute TRAP |  |  |  |  |
| B7E1 |  | XTRAP |  |  |
| B7El | 20EØAB | JSR | GETINT | ; CONVERT LINE \# TO POSITIVE |
| B7E4 | A5D4 | LDA | FRØ | ; SAVE LINE \# LOW AS TRAP LINE |
| B7E6 | 85BC | STA | TRAPLN | ; IN CASE OF LATER ERROR |
| B7E8 | A5D5 | LDA | FRØ+1 | ; X |
| B7EA | 85 BD | STA | TRAPLN+1 | ; X |
| B7EC | $6 \emptyset$ | RTS |  |  |

## Source Code

| XON - Execute ON |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| B7ED |  | XON |  |  |
| B7ED | 208AB8 | JSR | : SAVDEX | ; SAVE INDEX INTO LINE |
| B7F0 | 2øE9AB | JSR | GET1INT | ; GET 1 BYTE INTEGER |
| B7F3 | A5D4 | LDA | FRØ | ; GET VALUE |
| B7F5 | F020 * B817 | BEQ | : ERV | ; IF ZERO, FALL THROUGH TO NEXT STMT |
|  |  | ; |  |  |
| B7F7 | A4A8 | LDY | STINDEX | ; GET STMT INDEX |
| B7F9 | 88 | DEY |  | ; BACK UP TO GOSUB/GOTO |
| B7FA | B18A | LDA | [STMCUR], Y | ; GET CODE |
| B7FC | C917 | CMP | \#CGTO | ; IS IT GOTO? |
| B7FE | Fø03 * $\mathrm{B80} 3$ | BEQ | : GO | ; IF YES, DON'T PUSH ON RUN STACK |
|  |  | ; |  |  |
|  |  | ; |  |  |
|  |  | ; | THIS IS ON - GOSUB: | PUT ELEMENT ON RUN STACK |
|  |  | ; |  |  |
| B8øø | $20 \mathrm{CAB6}$ | JSR | XGS 1 | ; PUT ELEMENT ON RUN STACK <br> ; FOR RETURN |
|  |  | ; |  |  |
| B8ø3 |  | : GO |  |  |
| В8ø3 | A5D4 | LDA | FRø | ; GET INDEX INTO EXPRESSIONS |
| B8ø5 | 85B3 | STA | ONLOOP | ; SAVE FOR LOOP CONTROL |
| B807 |  | : ONI |  |  |
| B897 | 20D5AB | JSR | GETPINT | ; GET + INTEGER |
| B80A | C6B3 | DEC | ONLOOP | ; IS THIS THE LINE \# WE WANT? |
| B8øC | Føø6 ^B814 | BEQ | : ON2 | ; IF YES, GO DO IT |
|  | 2010B9 | ; JSR | TSTEND | ; ARE THERE MORE EXPRESSIONS |
| B811 | 90F4 ^ $\mathrm{B} 8 \emptyset 7$ | BCC | : ON1 | ; IF YES, THEN EVAL NEXT ONE |
| B813 | 60 | RTS |  | ; ELSE FALL THROUGH TO NEXT STMT |
| B814 |  | : ON2 |  |  |
| B814 | 4CA6B6 | JMP | XGO2 | ; JOIN GOTO |
|  |  | ; |  |  |
|  |  | ; |  |  |
| B817 |  | : ERV |  |  |
| B817 | $6 \varnothing$ | RTS |  | ; FALL THROUGH TO NEXT STMT. |

## Execution Control Statement Subroutines

## SETLINE - Set Up Line Pointers



FIXRSTK - Fix Run Stack - Remove Old FORs

| $*$ | ON ENTRY | A - VARIAble \# IN CURRENT FOR |
| :--- | :--- | :--- |
| * | ON EXIT | RUNSTK CLEAR OF ALL FOR'S |



POPRSTK - Pop Element from Run Stack


## Source Code


:RCONT - Contract Run Stack

:REXPAN - Expand Run Stack

:SAVRTOP - Save Top of Run Stack in ZTEMP1

| B881 |  | $:$ SAVRTOP |  |  |
| :--- | :--- | :---: | :--- | :--- | :--- |
| B881 | A69ø | LDX | TOPRSTK | SAVE TOPRSTK |
| B883 | $86 C 4$ | STX | TEMPA | ; X |
| B885 | A691 | LDX | TOPRSTK+1 | ; X |
| B887 | $86 C 5$ | STX | TEMPA+1 |  |

:SAVDEX — Save Line Displacement

| B88A |  | SAVDEX |  |  |
| :--- | :--- | ---: | :--- | :--- | :--- |
| B88A | A4A8 | LDY | STINDEX | ; GET STMT INDEX |
| B88C | $84 B 3$ | STY | SAVDEX | SAVE IT |
| B88E | $6 \varnothing$ | RTS |  |  |

:MV6RS - Move 6-Byte Value to Run Stack


## :PL6RS - Pull 6 Bytes from Run Stack to FR1



RSTPTR - Reset Stack Pointers [STARP and RUNSTK]

| B8AF |  | RSTPTR |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| B8AF | A58C | LDA | STARP | ; | GET BASE OF STR/ARRAY SPACE LOW |
| B8B1 | 858E | STA | RUNSTK | ; | RESET |
| B8B3 | 8590 | STA | MEMTOP |  |  |
| B8B5 | 850 E | STA | APHM | ; | SET APPLICATION HIMEM |
| B8B7 | A58D | LDA | STARP+1 | ; | GET BASE STR/ARRAY SPACE HIGH |
| B8B9 | 858F | STA | RUNSTK+1 | ; | RESET |
| B8BB | 8591 | STA | MEMTOP +1 | ; | X |
| B8BD | 850 F | STA | APHM+1 | ; | SET APPLICATION HIMEM |
| B8BF | 60 | RTS |  |  |  |

ZVAR - Zero Variable

| B8Cø |  |  | ZVAR |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | ; |  |  |  |  |
| B8C0 | A686 |  | LDX |  | VVTP | ; | MOVE VARIABLE TABLE POINTER |
| B8C2 | 86F5 |  | STX |  | ZTEMP1 | ; | X |
| B8C4 | A487 |  | LDY |  | VVTP+1 | ; | X |
| B8C6 | 84F6 |  | STY |  | ZTEMP1+1 | ; | X |
|  |  |  | ; |  |  |  |  |
|  |  |  | ; | ARE | WE AT END OF | TABLE | ? |
|  |  |  | ; |  |  |  |  |
| B8C8 |  |  | : ZVARI |  |  |  |  |
| B8C8 | A6F6 |  | LDX |  | ZTEMP1+1 | ; | GET NEXT VARIABLE ADDR HIGH |
| B8CA | E489 |  | CPX |  | ENDVVT+1 | ; | IS IT < END VALUE HIGH |
| B8CC | $9 \varnothing \square 7$ | ^B8D5 | BCC |  | : ZVAR2 | ; | IF YES, MORE TO DO |
| B8CE | A6F5 |  | LDX |  | ZTEMP1 | ; | GET NEXT VARIABLE ADDR LOW |
| B8DØ | E488 |  | CPX |  | ENDVVT | ; | IS IT < END VALUE LOW |
| B8D2 | 9001 | *B8D5 | BCC |  | : ZVAR2 | ; | IF YES, MORE TO DO |
| B8D4 | 60 |  | RTS |  |  | ; | ELSE, DONE |
|  |  |  | ; |  |  |  |  |
|  |  |  | ; | ZERO | A VARIABLE |  |  |
|  |  |  | ; |  |  |  |  |
| B8D5 |  |  | : ZVAR2 |  |  |  |  |
| B8D5 | АØロロ |  | LDY |  | \# $\varnothing$ | ; | TURN OFF |
| B8D7 | BlF5 |  | LDA |  | [ZTEMP1],Y | ; | DIM FLAG |
| B8D9 | 29FE |  | AND |  | \# \$FE |  |  |
| B8DB | 91F5 |  | STA |  | [ZTEMP1],Y |  |  |
| B8DD | Aøø2 |  | LDY |  | \#2 | ; | INDEX PAST VARIABLE HEADER |
| B8DF | A2ø6 |  | LDX |  | \# 6 | ; | GET \# OF BYTES TO ZERO |
| B8E1 | A9øø |  | LDA |  | \# $\varnothing$ | ; | CLEAR A |
|  |  |  | TVAR |  |  |  |  |
| B8E 3 |  |  | : ZVAR3 |  |  |  |  |
| B8E3 | 91F5 |  | STA |  | [ZTEMP1],Y |  | ZERO BYTE |
| B8E5 | C8 |  | INY |  |  | ; | POINT TO NEXT BYTE |
| B8E6 | CA |  | DEX |  |  | ; | DEC POINTER |
| B8E 7 | DØFA | ^B8E3 | BNE |  | : ZVAR3 | ; | IF NOT $=\varnothing$, ZERO NEXT BYTE |

Source Code


| RUNINIT - Initialize Storage Locations for RUN |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| B8F8 |  | RUNINIT |  |  |
| B8F8 | Аøøø | LDY | \# $\varnothing$ | ; CLEAR A |
| B8FA | 84BA | STY | STOPLN | ; CLEAR LINE \# Stopped at |
| B8FC | 84BB | STY | STOPLN+1 | ; X |
| B8FE | 84B9 | STY | ERRNUM | ; CLEAR ERROR \# |
| B9øø | 84FB | STY | RADFLG | ; CLEAR FLAG TOR TRANSENDENTALS |
| B9ø2 | 84B6 | STY | DATAD | ; CLEAR DATA POINTERS |
| B904 | 84B7 | STY | DATALN | ; X |
| B996 | 84B8 | STY | DATALN+1 | ; X |
| B908 | 88 | DEY |  |  |
| B909 | 84 BD | STY | TRAPLN+1 | ; SET TRAP FLAG TO NO TRAP |
| B90B | 8411 | STY | BRKBYT | ; SET BRK BYTE OFF [\$FF] |
| B90D | 4C41BD | JMP | CLSALL | ; GO CLOSE ALL DEVICES |

## TSTEND - Test for End of Statement



## Error Message Routine

## Error Messages

| B916 | E6B9 | ERRNSF | INC | ERRNUM | FILE NOT SAVE FILE |
| :---: | :---: | :---: | :---: | :---: | :---: |
| B918 | E6B9 | ERRDNO | INC | ERRNUM | ; \#DNØ > 7 |
| B91A | E6B9 | ERRPTL | INC | ERRNUM | ; LOAD PGM TOO BIG |
| B91C | E6B9 | ERSVAL | INC | ERRNUM | ; STRING NOT VALID |
| B91E | E6B9 | XERR | INC | ERRNUM | ; EXECUTION OF GARBAGE |
| B920 | E6B9 | ERBRTN | INC | ERRNUM | ; BAD RETURNS |
| B922 | E6B9 | ERGFDE | INC | ERRNUM | ; GOSUB/FOR LINE DELETED |
| B924 | E6B9 | ERLTL | INC | ERRNUM | ; LINE TO LONG |
| B926 | E6B9 | ERNOFOR | INC | ERRNUM | ; NO MATCHING FOR |
| B928 | E6B9 | ERNOLN | INC | ERRNUM | ; LINE NOT FOUND [GOSUB/GOTO] |
| B92A | E6B9 | EROVFL | INC | ERRNUM | ; FLOATING POINT OVERFLOW |
| B92C | E6B9 | ERRAOS | INC | ERRNUM | ARG STACK OVERFLOW |
| B92E | E6B9 | ERRDIM | INC | ERRNUM | ; ARRAY/STRING DIM ERROR |
| B930 | E6B9 | ERRINP | INC | ERRNUM | ; INPUT STMT ERROR |
| B932 | E6B9 | ERRLN | INC | ERRNUM | ; VALUE NOT <32768 |
| B934 | E6B9 | ERROOD | INC | ERRNUM | ; READ OUT OF DATA |
| B936 | E6B9 | ERRSSL | INC | ERRNUM | ; STRING LENGTH ERROR |
| B938 | E6B9 | ERRVSF | INC | ERRNUM | ; VARIABLE TABLE FULL |
| B93A | E6B9 | ERVAL | INC | ERRNUM | VALUE ERROR |
| B93C | E6B9 | MEMFULL | INC | ERRNUM | ; MEMORY FULL |
| B93E | E6B9 | ERON | INC | ERRNUM | NO LINE \# FOR EXP IN ON |

## Source Code

## Error Routine



Print Error Message Part 1 [**ERR]


Print Message Part 2 [AT LINE]


Print Line Number

| B984 | Aøø1 |  | LDY | \#1 | ; SET DISPL |
| :---: | :---: | :---: | :---: | :---: | :---: |
| B986 | B18A |  | LDA | [STMCUR], Y | ; GET LINE \# HIGH |
| B988 | 85D5 |  | STA | FRØ+1 | ; SET IN FRØ FOR CONVERT |
| B98A | 88 |  | DEY |  | ; GET CURRENT LINE \# LOW |
| B98B | Bl8A |  | LDA | [STMCUR], Y | ; GET UNUSED LINE \# LOW |
| B98D | 85D4 |  | STA | FRø | ; SET IN FRØ LOW FOR CONVERT |
| B98F | 209CB9 | ; | JSR | : PRINUM | ; PRINT LINE \# |
| B98F |  | ; |  |  |  |
|  |  | ; |  |  |  |
|  |  | ; |  |  |  |

## Source Code

| B992 | :ERRDONE |  |  |  |
| :--- | :--- | :---: | :--- | :--- |
| B992 | 206EBD | JSR | PRCR | ; PRINT CR |
| B995 | A9Øø | LDA | \#Ø | ; CLEAR A |
| B997 | 85B9 | STA | ERRNUM | ; CLEAR ERROR \# |
| B999 | 4C6ØAø | JMP | SYNTAX |  |

Print Integer Number in FR0


## Execute Graphics Routines

XSETCOLOR - Execute SET COLOR

| B9B7 |  | XSETCOLOR |  |  |
| :---: | :---: | :---: | :---: | :---: |
| B9B7 | 20E9AB | JSR | GETIINT | ; GET REGISTER \# |
| B9BA | A5D4 | LDA | FRØ | ; GET \# |
| B9BC | C905 | CMP | \# 5 | ; IS IT < 5? |
| B9BE | Bø1A *B9DA | BCS | : ERCOL | ; IF NOT, ERROR |
| B9Cø | 48 | PHA |  | ; SAVE |
|  |  | ; |  |  |
| B9Cl | 2ØEØAB | JSR | GETINT | ; GET VALUE |
| B9C4 | A 5D 4 | ; LDA | FRø | ; GET VALUE*16+6 |
| B9C6 |  | ASLA |  | ; X |
| B9C6 | $+\emptyset A$ | ASL | A |  |
| B9C7 |  | ASLA |  | ; X |
| B9C7 | $+\emptyset A$ | ASL | A |  |
| B9C8 |  | ASLA |  | ; X |
| B9C8 | $+\varnothing A$ | ASL | A |  |
| B9C9 |  | ASLA |  | ; X |
| B9C9 | $+\emptyset A$ | ASL | A |  |
| B9CA | 48 | PHA |  | ; SAVE ON STACKS |
| B9CB | $20 \mathrm{E} \square \mathrm{AB}$ | JSR | GETINT | ; GET VALUE 3 |
| B9CE | 68 | PLA |  | ; GET VALUE 2*16 FROM STACK |
| B9CF | 18 | CLC |  |  |
| B9Dø | 65D4 | ADC | FRø | ; ADD IN VALUE 3 |
| B9D2 | A8 | TAY |  | ; SAVE VALUE 2*16 + VALUE 5 |
| B9D3 | 68 | PLA |  | ; GET INDEX |
| B9D4 | AA | TAX |  | ; PUT IN X |
| B9D5 | 98 | TYA |  | ; GET VALUE |
| B9D6 | 9 DC 402 | ; STA | CREGS, X | ; SET VALUE IN REGS |
| B9D9 | $6 \emptyset$ | RTS |  |  |


|  | ; |  |
| :--- | :--- | :--- |
| B9DA | i |  |
| B9DA | :ERSND |  |
| B9DA 2ø3AB9 | :ERCOL |  |
|  |  | JSR |

## XSOUND - Execute SOUND



| B9E6 |  |  | ASLA |  | ; | GET VALUE *2 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B9E6 | $+\varnothing A$ |  | ASL | A |  |  |
| B9E7 | 48 |  | PHA |  |  |  |
|  |  | ; |  |  |  |  |
| B9E8 | A9øø |  | LDA | \# | ; | SET TO ZERO |
| B9EA | 8Dø8D2 |  | STA | SREG1 | ; | X |
|  |  | ; |  |  |  |  |
| B9ED | A903 |  | LDA | \#3 |  |  |
| B9EF | 8DøFD2 |  | STA | SKCTL |  |  |
|  |  | ; |  |  |  |  |
| B9F2 | 2øEØAB |  | JSR | GETINT | ; | GET EXP2 |
| B9F5 | 68 |  | PLA |  | ; | GET INDEX |
| B9F6 | 48 |  | PHA |  | ; | SAVE AGAIN |
| B9F7 | AA |  | tax |  | ; | PUT IN INDEX REG |
| B9F8 | A5D4 |  | LDA | FRØ | ; | GET VALUE |
| B9FA | 9DøøD2 |  | STA | SREG2, X | ; | SAVE IT |
| B9FD | 2øEøAB | , | JSR | GETINT | ; | GET EXP3 |
| BAøø | A5D4 |  | LDA | FRø | ; | GET 16*EXP3 |
| BAø2 |  |  | ASLA |  | ; | x |
| BAø2 | $+\emptyset A$ |  | ASL | A |  |  |
| BAØ3 |  |  | ASLA |  | ; | X |
| BAØ3 | $+\varnothing A$ |  | ASL | A |  |  |
| BAø4 |  |  | ASLA |  | ; | X |
| BAø4 | $+\varnothing A$ |  | ASL | A |  |  |
| BA05 |  |  | ASLA |  | ; | X |
| BA05 | $+\square A$ |  | ASL | A |  |  |
| BAø6 | 48 |  | PHA |  | ; | SAVE IT |
| BA07 | 20EØAB | ; | JSR | GETINT | ; | GET EXP4 |
| BAøA | 68 |  | PLA |  | ; | GET 16*EXP3 |
| BAØB | A8 |  | tay |  | ; | SAVE IT |
| ВАФС | 68 |  | PLA |  | ; | GET INDEX |
| BAøD | AA |  | TAX |  | ; | PUT IN X |
| BAøE | 98 |  | TYA |  | ; | GET EXP3*16 |
| BAøF | 18 |  | CLC |  |  |  |
| BA1ø | 65D4 |  | ADC | FRø | ; | GET 16*EXP3+EXP4 |
| BA12 | 9Dø1D2 |  | STA | SREG3, X | ; | STORE IT |
| BA15 | 60 |  | RTS |  |  |  |

## XPOS - Execute POSITION

| BA16 |  | XPOS |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| BA16 | 2øEøAB | JSR | GETINT | ; | GET | INTEGER | INTO FRØ |
| BA19 | A5D4 | LDA | FRØ | ; | SET | X VALUE |  |
| BA1B | 8555 | STA | SCRX | ; | X |  |  |
| BA1D | A5D5 | LDA | FRØ+1 | ; | X |  |  |
| BAlF | 8556 | STA | SCRX+1 | ; | X |  |  |
|  |  | ; |  |  |  |  |  |
| BA21 | 20E9AB | JSR | GETIINT | ; | SET | Y VALUE |  |
| BA24 | A5D4 | LDA | FRø | ; | X |  |  |
| BA 26 | 8554 | STA | SCRY | ; | X |  |  |
| BA28 | 60 | RTS |  |  |  |  |  |

## XCOLOR - Execute COLOR

| BA29 |  | XCOLOR |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| BA29 | 20EøAB | JSR | GETINT | ; | GET | INTEGER INTO |
| BA2C | A5D4 | LDA | FRø |  |  |  |
| BA2E | 85C8 | STA | COLOR |  |  |  |
| BA3Ø | 60 | RTS |  |  |  |  |
| XDRAWTO - Execute DRAWTO |  |  |  |  |  |  |
| BA31 |  | XDRAWTO |  |  |  |  |
| BA31 | 2016 BA | JSR | xpos | ; | GET | X,Y POSITION |
| BA 34 | A 5 C 8 | LDA | COLOR | ; | GET | COLOR |
| BA36 | $8 \mathrm{DFB} \mathrm{C}^{2}$ | STA | SVCOLOR | ; | SET | IT |

Source Code

| BA 39 | A911 |  | LDA | \#ICDRAW | ; | GET | сом | MAND |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| BA3B | A2ø6 |  | LDX | \# 6 | ; | SET | DEV | CE |
| BA3D | 2øC4BA |  | JSR | GLPCX | ; | SET | THE |  |
| BA4ø | A9øC | ; | LDA | \#SøC | ; | SET | AUX | 1 |
| BA42 | 9D4Aø3 |  | STA | ICAUXI, x |  |  |  |  |
| BA45 | A9øб |  | LDA | \# $\varnothing$ | ; | SET | AUX | 2 |
| BA47 | 9D4Bø3 |  | STA | ICAUX2, X |  |  |  |  |
| BA4A | 2ø24BD |  | JSR | 107 |  |  |  |  |
| BA4D | 4CB3BC |  | JMP | IOTEST |  |  |  |  |

XGR - Execute GRAPHICS


XPLOT - Execute PLOT

| BA76 | XPLOT |  |  |  |
| :--- | :--- | :--- | :--- | :--- |
| BA76 | $2016 B A$ | JSR | XPOS | ; SET X,Y POSITION |
|  |  |  |  |  |
| BA79 | A5C8 | LDA | COLOR | GET COLOR |
| BA7B | A2ø6 | LDX | \#6 | GET DEVICE \# |
| BA7D | 4CA1BA | JMP | PRCX | GO PRINT IT |

## Input/Output Routines

BA8Ø
LOCAL
GETLINE - Get a Line of Input


| BA96 | 2øC4BA | JSR | GLPCX |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| BA99 | 2øøABD | JSR | IO1 | ; GO DO I/O |
| BA9C | 4CB3BC | JMP | IOTEST | ; GO TEST RESULT |

PUTCHAR - Put One Character to List Device

| BA9F |  | PRCHAR |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| BA9F |  | PUTCHAR |  |  |  |  |  |
| BA9F | A6B5 | LDX |  | LISTDTD |  | ; GET LIST DEVICE |  |
| BAAI |  | PRCX |  |  |  |  |  |
| BAAI | 48 | PHA |  |  |  | ; SAVE IO BYTE |  |
| BAA 2 | 2øC6BA | JSR |  | GLPX |  | ; SET DEVICE |  |
| BAA5 | BD4Aø3 | LDA |  | ICAUX1, X |  | ; SET UP ZERO PAG | IOCB |
| BAA8 | 852A | STA |  | ICAUXI-IOCB+ | ICB | ; X |  |
| BAAA | BD4Bø3 | LDA |  | ICAUX2, X |  | ; X |  |
| BAAD | 852B | STA |  | ICAUX2-IOCB+ | ICB | ; X |  |
|  |  | ; |  |  |  |  |  |
| BAAF | 68 | PLA |  |  |  |  |  |
| BABø | AS | TAY |  |  |  |  |  |
| BAB1 | 20B8BA | JSR |  | : PDUM |  |  |  |
|  |  | ; |  |  |  |  |  |
|  |  | ; | RET | JRN HERE FROM | ROU | TINE |  |
| BAB4 | 98 | TYA |  |  |  | ; TEST STATUS |  |
| BAB5 | $4 \mathrm{CB6BC}$ | JMP |  | IOTES 2 |  |  |  |
|  |  | ; |  |  |  |  |  |
|  |  | ; |  |  |  |  |  |
| BAB8 |  | : PDUM |  |  |  |  |  |
| BAB8 | BD47ø3 | LDA |  | ICPUT $+1, \mathrm{X}$ |  | ; GO TO PUT ROUTINE |  |
| BABB | 48 | PHA |  |  |  | ; X |  |
| BABC | BD4603 | LDA |  | ICPUT, X |  | ; X |  |
| BABF | 48 | PHA |  |  |  | ; X |  |
| BACø | 98 | TYA |  |  |  | ; X |  |
| BACl | Aø92 | LDY |  | \#\$92 |  | ; LOAD VALUE FOR CIO | ROUTINE |
| BAC3 | 60 | RTS |  |  |  |  |  |
|  |  | ; |  |  |  |  |  |
| BAC4 | 85 C 0 | GLPCX | STA | IOCMD |  |  |  |
| BAC6 |  | GLPX |  |  |  |  |  |
| BAC6 | 86 Cl | STX |  | IODVC |  | ; AS I/O DEVICE |  |
| BAC8 | 4CA6BC | JMP |  | LDDVX |  | ; LOAD DEVICE X |  |

XENTER - Execute ENTER

| BACB | XENTER |  |  |  |
| :--- | :--- | ---: | :--- | :--- |
| BACB | A9ø4 | LDA | \#SØ4 | ; OPEN INPUT |
| BACD | $2 \emptyset D D B A ~$ | JSR | ELADVC | ; GO OPEN ALT DEVICE |
| BADØ | 85B4 | STA | ENTDTD | ; SET ENTER DEVICE |
| BAD2 | 4C6ØAØ | JMP | SYNTAX |  |

FLIST - Open LIST File


Source Code

| BAF4 | A9ø7 | LDA | $\# 7$ |
| :--- | :--- | :--- | :--- |
| BAF6 | $6 \emptyset$ | RTS |  |

## RUN from File

| BAF7 | A9FF | FRUN LDA |  |
| :--- | :--- | :---: | :---: | :---: |
| BAF9 | Døø2 | BAFD | BNE |

## XLOAD - Execute LOAD Command

| BAFB |  | XLOAD |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| BAFB | A90ø | LDA |  | \# $\varnothing$ | ; SET LOAD MODE |
| BAFD | 48 | : LDø | PHA |  | ; SAVE R/L TYPE |
| BAFE | A904 | LDA |  | \#Ø4 | ; GO OPEN FOR INPUT |
| BBøø | $2 \emptyset \mathrm{DDBA}$ | JSR |  | ELADVC | ; THE SPECIFIED DEVICE |
| BBø3 | 68 | PLA |  |  | ; GET R/L TYPE |
|  |  | ; |  |  |  |
| BB04 |  | XLOAD1 |  |  |  |
| BBø4 | 48 | PHA |  |  | ; SAVE R/L TYPE |
| BB05 | A907 | LDA |  | \# ICGTC | ; CMD IS GET TEXT CHARS |
| BBØ7 | 85C0 | STA |  | IOCMD |  |
| BB09 | 85CA | STA |  | LOADFLG | ; SET LOAD IN PROGRESS |
| BBØВ | 20A6BC | JSR |  | LDDVX | ; LOAD DEVICE X REG |
| BBøE | AØØE | LDY |  | \#ENDSTAR-OUTBUFF | ; $\mathrm{Y}=$ REC LENGTH |
| BBlø | $2 \emptyset 10 \mathrm{BD}$ | JSR |  | 103 | ; GO GET TABLE BLOCK |
| BB13 | 20В3BC | JSR |  | IOTEST | ; TEST I/O |
| BB16 | AD8005 | LDA |  | MISCRAM+OUTBUFF | ; IF FIRST 2 |
| BB19 | ØD8105 | ORA |  | MISCRAM+OUTBUFF+ | 1 ; BYTES NOT ZERO |
| BBlC | Dø38 ^BB56 | BNE |  | : LDFER | ; THEN NOT SAVE FILE |
| BB1E | A 28 C | ; LDX |  | \#STARP | START AT STARP DISPL |
| BB2Ø | 18 | : LD 1 | CLC |  |  |
| BB21 | A580 | LDA |  | OUTBUFF | ; ADD LOMEM TO |
| BB23 | 7Døøø5 | ADC |  | MISCRAM, X | ; LOAD TABLE DISPL |
| BB26 | A8 | TAY |  |  |  |
| BB27 | A 581 | LDA |  | OUTBUFF+1 |  |
| BB29 | 7 D 0105 | ADC |  | MISCRAM $+1, \mathrm{X}$ |  |
|  |  | ; CMP |  |  |  |
| BB2C | CDE602 | CMP |  | HIMEM +1 | ; IF NEW VALUE NOT |
| BB2F | 9øøA ~BB3B | BCC |  | : LD3 | ; LESS THEN HIMEM |
| BB31 | Døø5 ^BB38 | BNE |  | : LD2 | ; THEN ERROR |
| BB33 | CCE502 | CPY |  | HIMEM |  |
| BB36 | $9 \emptyset 03$ * BB3B | BCC |  | : LD3 |  |
| BB38 | 4ClAB9 | : LD2 | JMP | ERRPTL |  |
| BB3B | 9501 | ; LD3 | STA | 1, X | ; ELSE SET NEW TABLE VALUE |
| BB3D | $940 \emptyset$ | STY |  | $\emptyset, ~$ Х |  |
| BB3F | CA | DEX |  |  | ; DECREMENT TO PREVIOUS TBL ENTRY |
| BB4ø | CA | DEX |  |  |  |
| BB41 | E082 | CPX |  | \#VNTP | ; IF NOT AT LOWER ENTRY |
| BB43 | BøDB ^BB2ø | BCS |  | : LD1 | ; THEN CONTINUE |
|  |  | ; JSR |  |  |  |
| BB45 | 2の88BB | JSR |  | : LSBLK | ; LOAD USER AREA |
| BB48 | 2066B7 | JSR |  | XCLR | ; EXECUTE CLEAR |
| BB4B | A9øø | LDA |  | \# $\varnothing$ | ; RESET LOAD IN PROGRESS |
| BB4D | 85CA | STA |  | LOADF LG | ; X |
| BB4F | 68 | PLA |  |  | ; LOAD R/S STATUS |
| BB50 | Føø1 *BB53 | BEQ |  | : LD4 | ; BR IF LOAD |
| BB52 | 60 | RTS |  |  | ; RETURN TO RUN |
| BB53 |  | : LD4 |  |  |  |
| BB53 | 4C5ØAØ | JMP |  | SNXI | ;GO TO SYNTAX |
|  |  | , |  |  |  |
| BB 56 |  | : LDFER |  |  |  |
| BB56 | A9øø | LDA |  | \#ø | ; RESET LOAD IN PROGRESS |
| BB58 | 85CA | STA |  | LOADFLG | ; X |
| BB5A | 2016B9 | JSR |  | ERRNSF | ; NOT SAVE FILE |

## Source Code

## XSAVE - Execute SAVE Command

| BB5D |  | XSAVE |  |  |
| :---: | :---: | :---: | :---: | :---: |
| BB5D | A908 | LDA | \#Ø8 | ; GO OPEN FOR OUTPUT |
| BB5F | 20DDBA | JSR | ELADVC | ; ṪHE SPECIFIED DEVICE |
|  |  | ; |  |  |
| BB62 |  | XSAVEl |  |  |
| BB62 | A90B | LDA | \#ICPTC | ; I/O CMD IS PUT TEXT CHARS |
| BB64 | 85Cb | STA | IOCMD | ; SET I/O CMD |
| BB66 | A28Ø | ; LDX | \#OUTBUFF | ; MOVE RAM TABLE PTRS |
| BB68 | 38 | :SVI SEC |  | ; [OUTBUFF THRU ENSTAR] |
| BB69 | В5øø | LDA | ø, x | ; TO LBUFF |
| BB6B | E58Ø | SBC | OUTBUFE | ; AS DISPLACEMENT |
| BB6D | 9Døøø5 | STA | MISCRAM, X | ; FROM LOW MEM |
| BB70 | E8 | INX |  |  |
| BB71 | B50ø | LDA | $\emptyset, \mathrm{X}$ |  |
| BB73 | E581 | SBC | OUTBUFF+1 |  |
| BB75 | 9DØøø5 | STA | MISCRAM, X |  |
| BB78 | E8 | INX |  |  |
| BB79 | Eø8E | CPX | \#ENDSTAR |  |
| BB7B | 90EB ^ BB68 | BCC | : SVI |  |
|  |  | ; |  |  |
| BB7D | 2ØA6BC | JSR | LDDVX | ; OUTPUT LBUFF |
| BB8ø | AøøE | LDY | \#ENDSTAR-OUTBUFF | ; FOR PROPER LENGTH |
| BB82 | 2010 BD | JSR | 103 |  |
| BB85 | 2 为3BC | JSR | IOTEST | ; TEST GOOD I/O |

LSBLK - LOAD or SAVE User Area as a Block

| BB88 |  | : LSBLK |  |  |
| :---: | :---: | :---: | :---: | :---: |
| BB88 | 20A6BC | JSR | LDDVX | ; LOAD DEVICE X REG |
| BB8B | A582 | LDA | VNTP | ; SET VAR NAME TBL PTR |
| BB8D | 85F3 | STA | INBUFF | ; AS START OF BLOCK ADR |
| BB8F | A583 | LDA | VNTP+1 |  |
| BB91 | 85F4 | STA | INBUFF+1 |  |
| BB93 | AC8Dø5 | LDY | MISCRAM+STARP+1 | ; $\mathrm{A}, \mathrm{Y}=\mathrm{BLOCK}$ LENGTH |
| BB96 | 88 | DEY |  |  |
| BB97 | 98 | TYA |  |  |
| BB98 | AC8C05 | LDY | MISCRAM+STARP |  |
| B69B | $2 \emptyset 12 \mathrm{BD}$ | JSR | 104 | ; GO DO BLOCK I/O |
| B69E | 2øB3BC | JSR | IOTEST |  |
| BBAl | 4CF1BC | JMP | CLSYSl | ;GO CLOSE DEVICE |

## XCSAVE - Execute CSAVE



XCLOAD - Execute CLOAD

| BBAC |  | XCLOAD |  |  |
| :---: | :---: | :---: | :---: | :---: |
| BBAC | A904 | LDA | \#4 | ; GET OPEN FOR OUTPUT |
| BBAE | 20В6BB | JSR | COPEN | ; OPEN CASSETTE |
|  |  | ; |  |  |
| BBBl | А9øø | LDA | \# $\emptyset$ | ; GET LOAD TYPE |
| BBB3 | 4 CO 4 BB | JMP | XLOAD1 | ; DO LOAD |

## COPEN - OPEN Cassette



## Source Code



SOPEN - OPEN System Device


| XXIO - Execute XIO Statement |  |  |  |
| :--- | ---: | :--- | :--- |
| BBE5 | XXIO |  |  |
| BBE5 | 2øø4BD | JSR | GIOCMD |
| BBE8 | 4CEDBB | JMP | XOP1 |

## XOPEN - Execute OPEN Statement

| BBEB |  | XOPEN |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| BBEB | A903 | LDA | \# ICOIO | ; | LOAD OPEN CODE |
| BBED | 85 CD | XOP1 | STA IOCMD |  |  |
| BBEF | 209FBC | JSR | GIODVC | ; | GET DEVICE |
| BBF 2 | 2004 BD | JSR | GIOCMD | ; | GET AUXI |
| BBF5 | 48 | PHA |  |  |  |
| BBF6 | $2 \emptyset 04 \mathrm{BD}$ | JSR | GIOCMD | ; | GET AUX2 |
| BBF9 | A8 | TAY |  | ; | AUX2 IN Y |
| BBFA | 68 | PLA |  | ; | AUXI IN A |
| BBFB |  | XOP 2 |  |  |  |
| BBFB | 48 | PHA |  | ; | SAVE AUXI |
| BBFC | 98 | TYA |  |  |  |
| BBFD | 48 | PHA |  | ; | SAVE AUX2 |
|  |  | ; JSR |  |  |  |
| BBFE | 2ØEØAA | JSR | EXEXPR | ; | GET FS STRING |
| BC01 | 2079BD | JSR | SETSEOL | ; | GIVE STRING AN EOL |
| BC04 | 20A6BC | ; |  |  |  |
| BC07 | 68 | PLA |  |  |  |
| BCø8 | 9D4Bø3 | STA | ICAUX2, X | ; | SET AUX 2 |
| BCØB | 68 | PLA |  | ; | GET AUX 1 |
| BCøC | 9D4AØ3 | STA | ICAUXI, X | ; |  |
| BC ¢ F | 2ØØABD | JSR | IO1 | ; | GO DO I/O |
| BC12 | 2099BD | ; JSR | RSTSEOL | ; | RESTORE STRING EOL |

```
\begin{tabular}{llll} 
BC15 & \(2051 D A\) & JSR & INTLBF \\
BC18 & 4 CB3BC & JMP & IOTEST
\end{tabular}
```


## XCLOSE - Execute CLOSE

| BC1B | XCLOSE |  |
| :--- | ---: | :--- |
| BC1B A9øC | LDA | \#ICCLOSE |

GDVCIO - General Device I/O

| BCID | GDVCIO |  |  |  |
| :--- | :--- | :---: | :--- | :--- |
| BC1D | $85 C \emptyset$ | STA | IOCMD | SET CMD |
| BC1F | $299 F B C$ | JSR | GIODVC | GET DEVICE |
| BC22 | $2024 B D$ | GDIO1 | JSR | IO7 |
| BC25 | 4 CB3BC | JMP | IOTEST | GO DO I/O |
|  |  |  |  | GO TEST STATUS |

## xSTATUS - Execute STATUS

| BC28 |  |  |  |  |
| :--- | :--- | ---: | :--- | :--- |
| BC28 | 2ø9FBC | JSR | GIODVC | ; GET DEVICE |
| BC2B | A9ØD | LDA | \#ICSTAT | ; STATUS CMD |
| BC2D | $2 \emptyset 26 B D$ | JSR | IO8 | GO GET STATUS |
| BC3Ø | 2øFBBC | JSR | LDIOSTA | LOAD STATUS |
| BC33 | 4C2DBD | JMP | ISVARI | GO SET VAR |

## XNOTE - Execute NOTE

| BC36 |  | XNOTE |
| :--- | :--- | ---: |
| BC36 | A926 | LD |
| BC38 | $2 \emptyset 1 D B C$ | JS |
| BC3B | BD4Cø3 | LD |
| BC3E | BC4Dø3 | LD |
| BC41 | $2 \emptyset 2 F B D$ | JS |
| BC44 | $2 \emptyset A 6 B C$ | JS |
| BC47 | BD4EØ3 | LD |
| BC4A | $4 C 2 D B D$ | JM |
|  |  |  |

XPOINT - Execute POINT
BC4D
xPOINT

| BC4D |  |
| :--- | :--- |
| BC4D | $2 \emptyset 9 F B C$ |
| BC5Ø | $2 \emptyset D 5 A B$ |
| BC53 | $2 \emptyset A 6 B C$ |
| BC56 | A5D4 |
| BC58 | 9D4CØ3 |
| BC5B | A5D5 |
| BC5D | 9D4DØ3 |
| BC6Ø | 2ØD5AB |
| BC63 | $2 \emptyset A 6 B C$ |
| BC66 | A5D4 |
| BC68 | 9D4EØ3 |
| BC6B | A925 |
| BC6D | $85 C \emptyset$ |
| BC6F | $4 C 22 B C$ |

JSR
JSR
JSR
LDA
STA
LDA
STA
JSR
JSR
LDA
STA
LDA
STA
JMP

| GIODVC | $;$ GET I/O DEVICE NO. |
| :--- | :--- |
| GETPINT | $;$ GET SECTOR NO. |
| LDDVX | $;$ GET DEVICE X |
| FRØ | SET SECTOR NO. |
| ICAUX3,X |  |
| FRØ+1 |  |
| ICAUX4, X | GET DATA LENGTH |
| GETPINT | $;$ LOAD DEVICE X |
| LDDVX | $;$ SET AL |
| FRØ | $;$ SET DATA LENGTH |
| ICAUX5, X | GOINT CMD |
| \#\$25 | GO DO |
| IOCMD |  |

## XPUT - Execute PUT

| BC72 |  | XPUT |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| BC72 | 209 FBC | JSR | GIODVC | ; | GET DEVICE | \# |
|  |  | ; |  |  |  |  |
| BC75 | 2ØEØAB | JSR | GETINT | ; | GET DATA |  |
| BC78 | A5D4 | LDA | FRø | ; | X |  |
| BC7A | A6Cl | LDX | IODVC | ; | LOAD DEVICE | \# |
| BC7C | 4CAlBA | JMP | PRCX | ; | GO PRINT |  |
| XGET - Execute GET |  |  |  |  |  |  |
| BC7F |  | XGET |  |  |  |  |
| BC7F | 209FBC | JSR | GIODVC | ; | GET DEVICE |  |
|  |  | ; |  |  |  |  |
| BC82 |  | GET1 |  |  |  |  |
| BC82 | A907 | LDA | \# ICGTC | ; | GET COMMAND |  |
| BC84 | 85C0 | STA | IOCMD | ; | SET COMMAND |  |

Source Code

| BC86 | Aøø1 | LDY | \＃ 1 | SET BUFF LENGTH＝1 |
| :---: | :---: | :---: | :---: | :---: |
| BC88 | 2010 BD | JSR | 103 | ；DO IO |
| BC8B | 20B3BC | JSR | IOTEST | ；TEST I／O |
| BC8E | АØøø | LDY | \＃Ø | ；GET CHAR |
| BC9ø | B1F3 | LDA | ［INBUFF］，Y | ；X |
| BC92 | 4C2DBD | JMP | ISVARI | ；ASSIGN VAR |

XLOCATE－Execute LOCATE

| BC95 |  | XLOCATE |
| :--- | :--- | ---: |
| BC95 | $2 \emptyset 16 \mathrm{BA}$ | JSR |
| BC98 | A2Ø6 | LDX |
| BC9A | $2 \emptyset C 6 B A$ | JSR |
| BC9D | DØE3 ${ }^{\wedge}$ BC82 | BNE |


| XPOS | ；GET X，Y POSITION |
| :--- | :--- |
| $\# 6$ | ；GET DEVICE \＃ |
| GLPX | ；X |
| GET1 | ；GO GET |

GIODVC－Get I／O Device Number

| BC9F | GIODVC |  |  |
| :--- | :--- | :--- | :--- |
| BC9F | $2 \emptyset \emptyset 2 B D$ | JSR | GIOPRM |
| BCA2 | 85C1 | STA | IODVC |
| BCA 4 | FøøA |  |  |

LDDVX－Load X Register with I／O Device Offset


IOTEST－Test I／O Status

| BCB3 |  | IOTEST |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| BCB3 | $2 \emptyset \mathrm{FBBC}$ | JSR |  | LDIOSTA | ；LOAD I／O STATUS |
| BCB6 |  | IOTES 2 |  |  |  |
| BCB6 | 3001 ＾BCB9 | BMI |  | SICKIO | ；BR IF BAD |
| BCB8 | 60 | RTS |  |  | ；ELSE RETURN |
| BCB9 |  | SICKIO |  |  |  |
| BCB9 | Aøøø | LDY |  | \＃$\varnothing$ | ；RESET DISPLAY FLAG |
| BCBB | 8CFEØ2 | STY |  | DSPFLG |  |
|  |  | ； |  |  |  |
| BCBE | C980 | CMP |  | \＃ICSBRK | ；IF BREAK |
| BCCø | DøøA＊BCCC | BNE |  | ：SIOl | ；SIMULATE ASYNC |
| BCC2 | 8411 | STY |  | BRKBY＇T | ；BREAK |
| BCC4 | A5CA | LDA |  | LOADFLG | ；IF LOAD FLAG SET |
| BCC6 | FØø3＊ BCCB | BEQ |  | ：SIOS | ； |
| BCC8 | 4СøのАの | JMP |  | COLDSTART | ；DO COLDSTART |
| BCCB |  | ：SIOS |  |  |  |
| BCCB | 60 | RTS |  |  |  |
|  |  | ； |  |  |  |
| BCCC | A4Cl | ：SIO1 | LDY | IODVC | ；PRE－LOAD I／O DEVICE |
| BCCE | C988 | CMP |  | \＃\＄88 | ；WAS ERROR EOF |
| BCD $\varnothing$ | FøØF＊BCEl | BEQ |  | ：SIO4 | ；BR IF EOF |
| BCD2 | 85B9 | ：SIO2 | STA | ERRNUM | ；SET ERROR NUMBER |
| BCD4 | Cø07 | CPY |  | \＃ 7 | ；WAS THIS DEVICE \＃7 |
| BCD6 | Døø3＊BCDB | BNE |  | ：SIO3 | ；BR IF NOT |
| BCD8 | 20F1BC | JSR |  | CLSYSD | ；CLOSE DEVICE 7 |
| BCDB | $2 \emptyset 72 \mathrm{BD}$ | ：SIO3 | JSR | SETDZ | ；SET L／D DEVICE $=$ Ø |
| BCDE | 4C4Ø日9 | JMP |  | ERROR | ；REPORT ERROR |


| BCE 1 | Cøø7 | :SIO4 | CPY | \# 7 | ; | WAS EOF ON D | DEVICE |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| BCE 3 | DøED * BCD2 | BNE |  | :SIO2 | ; | BR IF NOT |  |
| BCE 5 | A25D | LDX |  | \#EPCHAR | ; | WERE WE IN E | ENTER |
| BCE7 | E4C2 | CPX |  | PROMPT | ; |  |  |
| BCE9 | DØE 7 ^BCD2 | BNE |  | : SIO2 | ; | BR NOT ENTER |  |
| BCEB | 20F1BC | JSR |  | CLSYSD | ; | CLOSE DEVICE | E 7 |
| BCEE | 4C53AØ | JMP |  | SNX2 | ; | GO TO SYNTAX |  |

## CLSYSD - Close System Device



| GIOPRM - Get I/O Parameters |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| BDØ2 |  | GIOPRM |  |  |  |  |
| BD02 | E6A8 | INC |  | STINDEX | ; | SKIP OVER \# |
| BDØ4 | 20D5AB | GIOCMD | JSR | GETPINT | ; | GET POSITIVE INT |
| BDø7 | A5D4 | LDA |  | FRø |  | MOVE LOW BYTE TO |
| BD09 | $6 \emptyset$ | RTS |  |  |  |  |

I/O Call Routine

| BDØA | AØFF | IOI |  | LDY | \# 255 | ; BUFL $=255$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| BDøC | Døø2 * BDIØ |  | BNE |  | 103 |  |
| BDØE | Aøøの | IO2 |  | LDY | \# $\emptyset$ | ; BUFL $=\varnothing$ |
| BD1ø | A9øø | 103 |  | LDA | \# $\varnothing$ | ; BUFL < 256 |
| BD12 | 9D4903 | 104 |  | STA | ICBLH, X | ; SET BUFL |
| BD15 | 98 |  | TYA |  |  |  |
| BD16 | 9D4803 |  | STA |  | ICBLL, X |  |
| BD19 | A5F4 | 105 |  | LDA | INBUFF+1 | ; LOAD INBUFF VALUE |
| BD1B | A4F3 |  | LDY |  | INBUFF |  |
| BD1D | 9D4503 | 106 |  | STA | ICBAH, X | ; SE BUF ADR |
| BD2Ø | 98 |  | TYA |  |  |  |
| BD21 | 9D4403 |  | STA |  | ICBAL, $X$ |  |
| BD24 | A5Cø | 107 |  | LDA | IOCMD | ; LOAD COMMAND |
| BD26 | 9D4203 | 108 |  | STA | ICCOM, X | ; SET COMMAND |
| BD29 | 2Ø56E4 |  | JSR |  | CIO | ; GO DO I/O |
| BD2C | 60 |  | RTS |  |  | ; DONE |

ISVAR - I/O Variable Set


## Source Code



## PREADY - Print READY Message

| BD57 |  | PREADY |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| BD57 | A 206 | LDX |  | \#RML-1 | ; | GET READY MSG | LENGTH-1 |
| BD59 | 86F2 | PRDY1 | STX | CIX | ; | SET LEN REM |  |
| BD5B | BD67BD | LDA |  | RMSG, X | ; | GET CHAR |  |
| BD5E | 209FBA | JSR |  | PRCHAR | ; | PRINT IT |  |
| BD61 | A6F2 | LDX |  | CIX | ; | GET LENGTH |  |
| BD63 | CA | DEX |  |  |  |  |  |
| BD64 | 10F3 *BD59 | BPL |  | PRDY1 | ; | BR IF MORE |  |
| BD66 | 60 | RTS |  |  |  |  |  |
| BD67 | $\begin{aligned} & \text { 9B59444145 } \\ & \text { 529B } \end{aligned}$ | RMSG | DB | CR, 'YDAER', CR |  |  |  |
|  | $=\boxed{007}$ | RML | EQU | *-RMSG |  |  |  |

## PRCR — Print Carriage Return



SETDZ - Set Device 0 as LIST/ENTER Device

| BD72 | A90ø | SETDZ | LDA | \# $\varnothing$ |
| :---: | :---: | :---: | :---: | :---: |
| BD74 | 85B4 | STA |  | ENTDTD |
| BD76 | 85B5 | STA |  | LISTDTD |
| BD78 | 60 | RTS |  |  |

SETSEOL - Set an EOL [Temporarily] after a String EOL


## Source Code



## Source Code



| BE77 | A9ø0 | ATAN LDA | \#Ø | ; ARCTAN[X] |
| :---: | :---: | :---: | :---: | :---: |
| BE79 | 85F0 | STA | SGNFLG | ; SIGN FLAG OFF |
| BE7B | 85F1 | STA | XFMFLG | ; \& TRANSFORM FLAG |
| BE7D | A5D4 | LDA | FRø |  |
| BE7F | 297F | AND | \# ${ }^{\text {7 }}$ F |  |
| BE81 | C940 | CMP | \# $\$ 40$ | ; CHECK X VS $1 . \emptyset$ |
| BE83 | 3015 *BE9A | BMI | ATAN1 | ; $\mathrm{X}<1 . \emptyset$ - USE SERIES DIRECTLY |
| BE85 | A5D4 | LDA | FRø | ; $\mathrm{X}>=1 . \emptyset$ - SAVE SIGN \& TRANSFORM |
| BE87 | 2980 | AND | \# $\$ 80$ |  |
| BE89 | 85 Fb | STA | SGNFLG | ; REMEMBER SIGN |
| BE8B | E6Fl | INC | XFMFLG |  |
| BE8D | A 97 F | LDA | \# \$ 7F |  |
| BE8F | 25D4 | AND | FRø |  |
| BE91 | 85D4 | STA | FRØ | ; FORCE PLUS |
| BE93 | A 2EA | LDX | \#FP9S\&\$FF |  |
| BE95 | AØDF | LDY | \#FP9S/\$100 |  |
| BE97 | 2095DE | JSR | XFORM | ; CHANGE ARG TO [ $\mathrm{X}-1] /[\mathrm{X}+1]$ |
| BE9A |  | ATAN1 |  |  |
| BE9A | A2E6 | LDX | \#FPSCR\&\$FF | ; ARCTAN[X], $-1<X<1$ BY SERIES <br> ; OF APPROXIMATIONS |
| BE9C | Aøø5 | LDY | \#FPSCR/\$1øø |  |
| BE9E | 20A7DD | JSR | FSTØR | ; $\mathrm{X} \rightarrow$ - FSCR |
| BEAI | 20B6DD | JSR | FMOVE | ; $\mathrm{X} \rightarrow$ >FRl |
| BEA4 | $2 \emptyset \mathrm{DBDA}$ | JSR | FMUL | ; $\mathrm{X} * \mathrm{X} \rightarrow$ - $\mathrm{FR} \emptyset$ |
| BEA7 | BØ39 ^BEE2 | BCS | ATNOUT | ; Ø'FLOW |
| BEA9 | A90B | LDA | \#NATCF |  |
| BEAB | A 2 AE | LDX | \#ATCOEF\&\$FF |  |
| BEAD | AøDF | LDY | \#ATCOEF/\$16Ø |  |


| BEAF | $2 \emptyset 4 \emptyset D D$ |
| :--- | :--- |
| BEB2 | Bø2E | ^BEE2


| JSR |  | PLYEVL | ; $\mathrm{P}[\mathrm{X} * \mathrm{X}]$ |
| :---: | :---: | :---: | :---: |
| BCS |  | ATNOUT |  |
| LDX |  | \#FPSCR\&\$FF |  |
| LDY |  | \#FPSCR/\$1ØØ |  |
| JSR |  | FLDIR | ; $\mathrm{X} \rightarrow$ > FR1 |
| JSR |  | FMUL | ; $\mathrm{X}^{*} \mathrm{P}$ [ $\left.\mathrm{X} * \mathrm{X}\right]$ |
| BCS |  | ATNOUT | ; Ø'FLOW |
| LDA |  | XFMFLG | ; WAS ARG XFORM'D |
| BEQ |  | ATAN2 | ; NO |
| LDX |  | \#PIOV4\&\$FF | ; YES-ADD ARCTAN [1.ø] = PI/4 |
| LDY |  | \#PIOV4/\$1øø |  |
| JSR |  | FLDIR |  |
| JSR |  | FADD |  |
| LDA |  | SGNFLG | ; GET ORG SIGN |
| ORA |  | FRØ |  |
| STA |  | FRø | ; ATAN[-X] $=-\operatorname{ATAN}[\mathrm{X}]$ |
| ATAN2 | LDA | DEGFLG | ; RADIANS OR DEGREES |
| BEQ |  | ATNOUT | ; RAD - FINI |
| LDX |  | \#PIOV18\&\$FF | ; DEG - DIVIDE BY PI/18Ø |
| LDY |  | \#PIOV18/\$1øØ |  |
| JSR |  | FLDIR |  |
| JSR |  | FDIV |  |
| ATNOUT | RTS |  |  |

## SQR[X] — Square Root

| BEE 3 | 38 |
| :---: | :---: |
| BEE4 | $6 \emptyset$ |
| BEE5 | А9бø |
| BEE7 | 85F1 |
| BEE9 | A5D4 |
| BEEB | 3ØF6 ^BEE3 |
| BEED | C93F |
| BEEF | F017 * BF ¢8 |
| BEFl | 18 |
| BEF2 | 6901 |
| BEF4 | 85Fl |
| BEF6 | 85EØ |
| BEF8 | A901 |
| BEFA | 85El |
| BEFC | A204 |
| BEFE | A900 |
| BFøø | 95E2 |
| BF®2 | CA |
| BF®3 | $1 Ø \mathrm{FB}{ }^{\wedge} \mathrm{BF} \emptyset \emptyset$ |
| BFO5 | 2ø28DB |
| BF98 |  |
| BFø8 | A906 |
| BFØA | 85EF |
| BFøC | A2E6 |
| BFOE | АØø5 |
| BFlØ | 2ØA7DD |
| BF13 | 2ØB6DD |
| BF16 | A 293 |
| BF18 | AØBF |
| BFIA | 2089DD |
| BF1D | 2060DA |
| BF20 | A 2 E 6 |
| BF22 | AøØ5 |
| BF24 | 2ø98DD |
| BF27 | 2øDBDA |
| BF 2 A | A2EC |
| BF2C | Aøø5 |
| BF2E | 20A7DD |
| BF31 | 2ØB6DD |
| BF34 | A 2E6 |
| BF36 | AØロ5 |
| BF38 | 2ø89DD |

```
;
SQRERR SEC ;SET FAIL
    ; SQR
SQR 
    #$3F
    FSQR ; X IN RANGE OF APPROX - GO DO
        #1
        XFMFLG ; NOT IN RANGE - TRANSFORM
        FRI ; MANTISSA = 1
        #1
        FR1+1
        #FPREC-2
        #ø
        STA FRI+2,X
        DEX
        BPL
        SQR1
        FDIV ; x/1暗*N
        ;SQR[X], Ø. 1<=X<1.\emptyset
        #6
        #FSCR&SFF
        #FSCR/$1\emptyset\emptyset ;STASH X IN FSCR
        FMOVE ;X->FRI
        #FTWO&$FF
        #FTWO/$100
        FLDØR ;2.\emptyset->FR\emptyset
        FSUB ;2.\varnothing-X
        #FSCR&$FF
        #FSCR/$1\emptyset\varnothing ; X->FRI
        FMUL ;X*[2.ø-X] :IST APPROX
        #FSCR1&SFF
        #FSCR1/$1\emptyset\emptyset
        FST0R
        FMOVE ;Y->FRI
        #FSCR&$FF
        #FSCR/$1Ø\emptyset
        FLD\emptysetR
```

Source Code


Floating Point

| BF99 | $=$ D8ø |
| :--- | :--- |
| D8øø |  |$\quad$| ORG |
| :--- |
| LOCAL |$\quad$ FPORG

ASCIN - Convert ASCII Input to Internal Form


## Source Code



Source Code

| D850 | FøC6 ~D818 | BEQ | : IN1 | ; GO FOR NEXT CHAR |
| :---: | :---: | :---: | :---: | :---: |
| D852 | C92D | CMP | \#' - ' | ; IS IT MINUS? |
| D854 | Føøø *D856 | BEQ | :MINUS |  |
|  |  | ; |  |  |
|  |  | ; |  |  |
| D856 |  | :MINUS |  |  |
| D856 | 85EE | STA | NSIGN | ; SAVE SIGN FOR LATER |
| D858 | F0BE ^D818 | BEQ | : IN1 | ; UNCONDITIONAL BRANCH FOR NEXT CHAR |
|  |  | ; |  |  |
| D85A |  | : DP |  |  |
| D85A | A6Fl | LDX | DIGRT | ; IS DIGRT STILL = FF? |
| D85C | 1658 ^D8B6 | BPL | : EXIT | ; IF NOT, ALREADY HAVE DP |
| D85E | E8 | INX |  | ; INCR TO zero |
| D85F | 86 Fl | STX | DIGRT | ; SAVE |
| D861 | FØB5 ^D818 | BEQ | : IN1 | ; UNCONDITIONAL BR FOR NEXT CHAR |
|  |  | , |  |  |
| D863 |  | : EXP |  |  |
| D863 | A5F2 | LDA | CIX | ; GET INDEX |
| D865 | 85EC | STA | FRX | ; SAVE |
| D867 | 2094DB | JSR | : GETCHAR | ; GET NEXT CHAR |
| D86A | Bø37 ^D8A3 | BCS | : NON2 | ; BR IF NOT NUMBER |
|  |  | ; |  |  |
|  |  | ; | IT'S A NUMBER IN AN | EXPONENT |
|  |  | ; |  |  |
| D86C |  | : EXP 2 |  |  |
| D86C | AA | TAX |  | ; SAVE IST CHAR OF EXPONENT |
| D86D | A5ED | LDA | EEXP | ; GET \# OF CHAR OVER 9 |
| D86F | 48 | PHA |  | ; SAVE IT |
| D870 | 86ED | STX | EEXP | ; SAVE IST CHAR OF EXPONENT |
| D872 | 2ø94DB | JSR | : GETCHAR | ; GET NEXT CHAR |
|  |  | ; |  |  |
|  |  | ; |  |  |
| D875 | B017 D88E | BCS | : EXP 3 | ; IF NOT \# NO SECOND DIGIT |
| D877 | 48 | PHA |  | ; SAVE SECOND DIGIT |
| D878 | A5ED | ; LDA | EEXP | ; GET lST DIGIT |
| D87A |  | ASLA |  | ; GET DIGIT * $1 \varnothing$ |
| D87A | $+\emptyset A$ | ASL | A |  |
| D87B | 85ED | STA | EEXP | ; X |
| D87D |  | ASLA |  | ; X |
| D87D | $+\emptyset A$ | ASL | A |  |
| D87E |  | ASLA |  | ; X |
| D87E | $+\emptyset A$ | ASL | A |  |
| D87F | 65 ED | ADC | EEXP | ; X |
| D881 | 85ED | STA | EEXP | ; SAVE |
| D883 | 68 | PLA |  | ; GET SECOND DIGIT |
| D884 | 18 | CLC |  |  |
| D885 | 65 ED | ADC | EEXP | ; GET EXPONENT INPUTTED |
| D887 | 85ED | STA | EEXP | ; SAVE |
|  |  | ; |  |  |
| D889 | A 4F2 | LDY | CIX | ; INC TO NEXT CHAR |
| D88B | $2 \emptyset 9 \mathrm{DDB}$ | JSR | : GCHR1 | ; X |
|  |  | ; |  |  |
|  |  | ; |  |  |
| D88E |  | : EXP 3 |  |  |
| D88E | A5EF | LDA | ESIGN | ; GET SIGN OF EXPONENT |
| D890 | F009 *D89B | BEQ | : EXP 1 | ; IF NO SIGN, IT IS + |
| D892 | A5ED | LDA | EEXP | ; GET EXPONENT ENTERED |
| D894 | 49 FF | EOR | \# ${ }^{\text {SFF }}$ | ; COMPLEMENT TO MAKE MINUS |
| D896 | 18 | CLC |  | ; X |
| D897 | 6901 | ADC | \#1 | ; X |
| D899 | 85ED | STA | EEXP | ; SAVE |
| D89B |  | : EXP1 |  |  |
| D89B | 68 | PLA |  | ; GET \# DIGITS MORE THAN 9 |
| D89C | 18 | CLC |  | ; CLEAR CARRY |
| D89D | 65ED | ADC | EEXP | ; ADD IN ENTERED EXPONENT |
| D89F | 85ED | STA | EEXP | ; SAVE EXPONENT |
| D8A1 | D613 *D8B6 | BNE | : EXIT | ; UNCONDITIONAL BR |

## Source Code

|  |  | ; | NON-NUMERIC IN EXPONENT |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | ; |  |  |  |  |  |
| D8A3 |  | :NON2 |  |  |  |  |  |
| D8A3 | C92B | CMP | \#'+' | ; | IS IT PLUS? |  |  |
| D8A5 | Føø6 ~D8AD | BEQ | : EPLUS |  | IF YES BR |  |  |
| D8A7 | C92D | CMP | \#' - ' |  | IS IT A MINUS? |  |  |
| D8A9 | Døø7 ^D8B2 | BNE | : NOTE |  | IF NOT, BR |  |  |
|  |  | ; |  |  |  |  |  |
|  |  | ; |  |  |  |  |  |
| D8AB |  | : EMIN |  |  |  |  |  |
| D8AB | 85EF | STA | ESIGN | ; | SAVE EXPONENT SIGN |  |  |
| D8AD |  | : EPLUS |  |  |  |  |  |
| D8AD | 2094 DB | JSR | : GETCHAR | ; | GET CHARACTER |  |  |
| D8BØ | 99BA * ${ }^{\text {D }} 86 \mathrm{C}$ | BCC | : EXP2 | ; | IF A \#, GO PROCESS | EXPONE | ENT |
|  |  | ; |  |  |  |  |  |
|  |  | ; |  |  |  |  |  |
|  |  | ; |  |  |  |  |  |
|  |  | E | E NOT PART OF OUR |  |  |  |  |
|  |  | ; |  |  |  |  |  |
| D8B2 |  | : NOTE |  |  |  |  |  |
| D8B2 | A 5EC | LDA | FRX | , | POINT TO 1 PAST E |  |  |
| D8B4 | 85F2 | STA | CIX | ; | RESTORE CIX |  |  |
|  |  | ; |  |  |  |  |  |
|  |  | F | FALL THRU TO EXIT |  |  |  |  |
|  |  | ; |  |  |  |  |  |
|  |  | W | WHOLE \# HAS BEEN I | PUT | TTED |  |  |
|  |  | ; |  |  |  |  |  |
| D8B6 |  | : EXIT |  |  |  |  |  |
|  |  | ; |  |  |  |  |  |
|  |  | ; B | BACK UP ONE CHAR |  |  |  |  |
|  |  | - DEC |  |  |  |  |  |
| D8B6 | C6F2 |  | CIX | ; | DECREMENT INDEX |  |  |
|  |  | ; |  |  |  |  |  |
|  |  | ; |  |  |  |  |  |
|  |  | ; |  |  |  |  |  |
|  |  | ; | WHERE EXP = ENTERED EXPONENT [COMPLEMENT OF -] |  |  |  |  |
|  |  |  | + \# DIGITS MORE THAN 9 |  |  |  |  |
|  |  | ; |  |  |  |  |  |
| D8B8 | A5ED | LDA | EEXP | , | GET EXPONENT |  |  |
| D8BA | A6Fl | LDX | DIGRT | , | GET \# DIGITS RIGHT O | OF DEC | CIMAL |
| D8BC | 3 - 5 ~D8C3 | BMI | : EXITI | , | NO DECIMAL POINT |  |  |
| D8BE | Føø3 *D8C3 | BEQ | : EXITI | ; | \# OF DIGITS AFTER D. | . P. $=\varnothing$ |  |
| D8C0 | 38 | SEC |  | , | GET EXP - DIGITS RIG | GHT |  |
| D8Cl | E5Fl | SBC | DIGRT | ; |  |  |  |
|  |  | ; | SHIFT RIGHT ALGEBRAIC TO DIVIDE BY $2=$ POWER OF 1øø |  |  |  |  |
|  | D8C3 |  |  |  |  |  |  |  |
|  |  |  | : EXITI |  |  |  |  |  |
| D8C3 | 48 | PHAROLA |  | ; SET CARRY WITH SIGN OF |  |  |  |
| D8C4 |  |  |  |  |  |  |  |
| D8C4 | +2A | ROL | A |  |  |  |  |
| D8C5 | 68 | PLA |  | ; | GET EXPONENT AGAIN |  |  |
| D8C6 |  | RORA |  | ; | SHIFT RIGHT |  |  |
| D8C6 | +6A | ROR | A |  |  |  |  |
| D8C7 | 85ED | STA | EEXP | ; | SAVE POWER OF $10 \emptyset$ |  |  |
| D8C9 | 9003 ~D8CE | BCC | : EVEN |  | IF NO CARRY \# EVEN |  |  |
|  | 2ØEBDB | ; |  |  |  |  |  |
| D8CE |  | : EVEN |  |  |  |  |  |
| D8CE | A5ED | LDA | EEXP | ; ADD $4 \emptyset$ FOR EXCESS $64+4$ FOR NORM |  |  |  |
|  |  |  |  |  |  |  |  |
| D8Dø | 18 | CLC |  | ; |  |  |  |
| D8D1 | 6944 | ADC | \# \$44 | ; | X |  |  |
| D8D3 | 85D4 | STA | FRØ | ; | SAVE AS EXPONENT |  |  |
|  |  |  |  |  |  |  |  |
| D8D5 | 2øøøDC | JSR | NORM | ; | NORMALIZE NUMBER |  |  |
| D8D8 | BøØ日 ^D8E5 | BCS | : IND2 |  | IF CARRY SET, IT'S A | AN ERR | ROR |
|  |  | ; |  |  |  |  |  |

Source Code

|  |  |  |  | MANTISS |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  |  |  |  |
| D8DA | A6EE |  | LDX | NSIGN | ; | IS SIGN OF \# MI | MINUS? |
| D8DC | Føø6 | *D8E4 | BEQ | : INDON | ; | IF NOT, BR |  |
|  |  |  |  |  |  |  |  |
| D8DE | A5D4 |  | LDA | FRø | ; | GET EXPONENT |  |
| D8EØ | 9980 |  | ORA | \#\$80 | ; | TURN ON MINUS \# | \# BIT |
| D8E2 | 85D4 |  | STA | FRø | ; | SET IN FRØ EXP |  |
| D8E4 |  |  | ON |  |  |  |  |
| D8E4 | 18 |  | CLC |  | ; | CLEAR CARRY |  |
| D8E5 |  |  |  |  |  |  |  |
| D8E5 | 60 |  | RTS |  |  |  |  |

FPASC - Convert Floating Point to ASCII


## Source Code



## Source Code



IFP - Convert Integer to Floating Point


## Source Code

| D9BC | A2б3 | LDX | \#3 | ; | CARRY NOW SET IF THERE WAS A BIT <br> BIGGEST INTEGER IS 3 BYTES |
| :---: | :---: | :---: | :---: | :---: | :---: |
| D9BE |  | : IFP2 |  |  |  |
|  |  |  | DOUBLE \# AND ADD IN | 1 | IF CARRY SET |
|  |  | ; |  |  |  |
| D9BE | B5D4 | LDA | FRø, X | ; | GET BYTE |
| D9Cø | 75D4 | ADC | FRø, X | ; | DOUBLE [ADDING IN CARRY |
|  |  |  |  |  | FROM SHIFT |
| D9C2 | 95D4 | STA | FRø, X | ; | SAVE |
| D9C4 | CA | DEX |  | ; | DECREMENT COUNT OF FRØ BYTES |
| D9C5 | DØF7 ^D9BE | BNE | : IFP2 | ; | IF MORE TO DO, DO IT |
|  |  | ; |  |  |  |
| D9C7 | 88 | DEY |  | ; | DECR COUNT OF INTEGER DIGITS |
| D9C8 | DØEE *D9B8 | BNE | : IFP1 | ; | IF MORE TO DO, DO IT |
| D9CA | D8 | CLD |  | ; | CLEAR DECIMAL MODE |
|  |  | ; |  |  |  |
|  |  | ; | SET EXPONENT |  |  |
|  |  | , |  |  |  |
| D9CB | A942 | LDA | \#\$42 | ; | INDICATE DECIMAL AFTER LAST |
|  |  |  |  |  | DIGIT |
| D9CD | 85D4 | STA | FRØ | ; | STORE EXPONENT |
|  |  | i |  |  |  |
| D9CF | 4CøØDC | JMP | NORM | ; | NORMALIZE |
|  |  | ; |  |  |  |

FPI - Convert Floating Point to Integer


## Source Code

|  |  | ; | MULT INTEGER RESULT | BY 10 |
| :---: | :---: | :---: | :---: | :---: |
|  |  | ; |  |  |
| D9EA | 2ø5ADA | JSR | : ILSHFT | ; GO SHIFT ONCE LEFT |
| D9ED | B653 ^DA42 | BCS | : ERVAL | ; IF CARRY SET THEN \# TOO BIG |
|  |  | ; |  |  |
| D9EF | A5F7 | LDA | ZTEMP4 | ; SAVE INTEGER *2 |
| D9F1 | 85F9 | STA | ZTEMP3 | ; X |
| D9F3 | A5F8 | LDA | ZTEMP4+1 | ; X |
| D9F5 | 85FA | STA | ZTEMP3+1 | ; X |
|  |  | ; |  |  |
| D9F7 | 205ADA | JSR | : ILSHFT | ; MULT BY *2 |
| D9FA | BØ46 *DA42 | BCS | : ERVAL | ; \# TOO BIG |
| D9FC | 205ADA | JSR | : ILSHFT | ; MULT BY *2 [NOW * 8 IN ZTEMP4] |
| D9FF | BØ41 *DA42 | BCS | : ERVAL | ; BR IF \# TOO BIG |
|  |  | ; |  |  |
| DA01 | 18 | CLC |  | ; ADD IN * 2 TO = *1ø |
| DA02 | A5F8 | LDA | ZTEMP4+1 | ; X |
| DAØ4 | 65FA | ADC | ZTEMP3+1 | ; X |
| DAø6 | 85F8 | STA | ZTEMP4+1 | ; X |
| DA08 | A5F7 | LDA | ZTEMP4 | ; X |
| DAØA | 65F9 | ADC | ZTEMP3 | ; X |
| DAØC | 85 F 7 | STA | ZTEMP4 | ; X |
| DAØE | Bø32 ^DA42 | BCS | : ERVAL | ; IF CARRY SET ERROR |
|  |  | ; |  |  |
|  |  | ; |  |  |
|  |  | ; | ADD IN NEXT DIGIT |  |
|  |  | ; |  |  |
| DA1ø | 20B9DC | JSR | : GETDIG | ; GET DIGIT IN A |
| DA13 | 18 | CLC |  |  |
| DA14 | 65F8 | ADC | ZTEMP4+1 | ; ADD IN DIGIT |
| DA16 | 85F8 | STA | ZTEMP4+1 | ; X |
| DA18 | A5F7 | LDA | ZTEMP4 | ; X |
| DA1A | 6900 | ADC | \#Ø | ; X |
| DA1C | Bø24 *DA42 | BCS | : ERVAL | ; BR IF OVERFLOW |
| DAIE | 85F7 | STA | ZTEMP4 | ; X |
|  |  | ; |  |  |
| DA2Ø | C6F5 | DEC | ZTEMP1 | ; DEC COUNTER OF DIGITS TO DO |
| DA22 | DøC6 *D9EA | BNE | :FPII | ; IF MORE TO DO, DO IT |
|  |  | ; |  |  |
|  |  | ; | ROUND IF NEEDED |  |
|  |  | ; |  |  |
| DA24 |  | : ROUND |  |  |
| DA24 | 20B9DC | JSR | : GETDIG | ; GET NEXT DIGIT IN A |
| DA27 | C905 | CMP | \#5 | ; IS DIGIT < 5 ? |
| DA29 | 90øD ^DA38 | BCC | : NR | ; IF YES, DON'T ROUND |
| DA2B | 18 | CLC |  | ; ADD IN 1 TO ROUND |
| DA2C | A5F8 | LDA | ZTEMP4+1 | ; X |
| DA2E | 6901 | ADC | \#1 | ; X |
| DA30 | 85F8 | STA | ZTEMP4+1 | ; X |
| DA32 | A5F7 | LDA | ZTEMP4 | ; X |
| DA34 | 690ø | ADC | \#Ø | ; X |
| DA36 | 85 F 7 | STA | ZTEMP4 | ; X |
|  |  | ; |  |  |
|  |  | ; | MOVE INTEGER TO FRø |  |
|  |  | ; |  |  |
| DA38 |  | : NR |  |  |
| DA38 | A5F8 | LDA | ZTEMP4+1 | ; GET INTEGER LOW |
| DA3A | 85D4 | STA | FRø | ; SAVE |
| DA3C | A5F7 | LDA | ZTEMP4 | ; GET INTEGER HIGH |
| DA3E | 85D5 | STA | FRø+1 | ; SAVE |
|  |  | ; |  |  |
| DA40 | 18 | CLC |  | ; CLEAR CC FOR GOOD RETURN |
| DA41 | 60 | RTS |  |  |
|  |  | ; |  |  |
|  |  | ; |  |  |
| DA42 |  | : ERVAL |  |  |
| DA42 | 38 | SEC |  | ; SET CARRY FOR ERROR RETURN |
| DA43 | 60 | RTS |  |  |
|  |  | * | ZFRø - ZERO FRø |  |
|  |  | * | ZFl - ZERO 6 BYTES A | AT LOC X |



## Floating Point Routines

## FADD - Floating Point Add Routine



FSUB - Floating Point Subtract Routine


## Source Code



## Source Code

| DAB6 |  | $\text { : SUB } 1$ |  |  |
| :---: | :---: | :---: | :---: | :---: |
| DAB6 | B5D5 | LDA | FRgM, X | ; GET FRØ BYTE |
| DAB8 | F5El | SBC | FRIM, X | ; SUB FRI BYTE |
| DABA | 95D5 | STA | FRøM, X | ; STORE |
| DABC | CA | DEX |  | ; DEC POINTER |
| DABD | 10F7 *DAB6 | BPL | : SUB1 | ; SUB NEXT BYTE |
|  |  | ; |  |  |
| DABF | 9004 *DAC5 | BCC | : SUB2 | ; IF THERE IS A BORROW DO |
| DACl | D8 | CLD |  | ; CLEAR DECIMAL MODE |
| DAC 2 | 4CØøDC | JMP | NORM |  |
|  |  | i |  |  |
|  |  | ; | TAKE COMPLEMENT SIGN |  |
|  |  |  |  |  |
| DAC5 |  | : SUB2 |  |  |
| DAC5 | A5D4 | LDA | FRø | ; GET EXPONENT |
| DAC7 | 4980 | EOR | \# \$80 | ; CHANGE SIGN OF MANTISSA |
| DAC9 | 85D4 | STA | FRØ | ; PUT IT BACK |
|  |  | ; |  |  |
|  |  | ; | COMPLEMENT MANTISSA |  |
|  |  | ; |  |  |
| DACB | 38 | SEC |  | ; SET CARRY |
| DACC | A2ø4 | LDX | \#FMPREC-1 | ; GET INDEX COUNTER |
| DACE |  | : SUB3 |  |  |
| DACE | A900 | LDA | \#Ø | ; GET ZERO |
| DAD® | F5D5 | SBC | FRØM, X | ; COMPLEMENT BYTE |
| DAD2 | 95D5 | STA | FRøM, X | ; STORE |
| DAD4 | CA | DEX |  | ; MORE TO DO |
| DAD5 | 10F7 ^DACE | BPL | : SUB3 | ; BR IF YES |
| DAD7 | D8 | ; CLD |  | ; CLEAR DECIMAL MODE |
| DAD8 | 4CøøDC | JMP | NORM | ; GO NORMALIZE |

FMUL - Multiply FRo by FR1


## Source Code



## Source Code

FDIV - Floating Point Divide


## Source Code

| DB63 | DØE9 *DB4E | BNE | :FRDI ; SUB AGAIN |
| :---: | :---: | :---: | :---: |
|  |  | ; |  |
|  |  | ; | SUBTRACT OF FR2 DIDN'T GO |
|  |  | ; |  |
| DB65 |  | :FAIL |  |
| DB65 | 2ØøFDD | JSR | FRA2E ; ADD FR2 BACK TO FRø |
|  |  | ; |  |
|  |  | ; | SHIFT LAST BYTE OF QUOTIENT ONE NIBBLE LEFT |
|  |  | ; |  |
| DB68 | Ø6D9 | ASL | QTEMP ; SHIFT 4 BITS LEFT |
| DB6A | 06D9 | ASL | QTEMP ; X |
| DB6C | 06D9 | ASL | QTEMP ; X |
| DB6E | Ø6D9 | ASL | QTEMP ; X |
| DB7Ø |  | : FRD2 |  |
|  |  | ; |  |
|  |  | ; | SUBTRACT FRI [DIVISOR] FROM FRE [DIVIDEND] |
|  |  | ; |  |
|  |  | ; |  |
| DB7Ø | A005 | LDY | \#FMPREC ; SET LOOP CONTROL |
| DB72 | 38 | SEC | ; SET CARRY |
| DB73 | F8 | SED | ; SET DECIMAL MODE |
| DB74 |  | : FRS 1 |  |
| DB74 | B9DAøØ | LDA | FRE, Y ; GET A BYTE FROM FRE |
| DB77 | F9Eøøø | SBC | FRI, Y ; SUB FRI |
| DB7A | 99DAøø | STA | FRE, Y ; STORE RESULT |
| DB7D | 88 | DEY |  |
| DB7E | 10F4 ^DB74 | BPL | :FRS 1 ; BR IF MORE TO DO |
| DB8Ø | D8 | CLD | ; CLEAR DECIMAL MODE |
|  |  | ; |  |
| DB81 | $9004{ }^{\text {^DB87 }}$ | BCC | :FAIL2 ; IF RESULT < Ø [FRE < FRI] BR |
|  |  | ; |  |
| DB83 | E6D9 | INC | QTEMP ; INCR \# TIMES SUB [QUOTIENT] |
|  |  | ; BNE |  |
| DB85 | DØE9 * DB7Ø | BNE | :FRD2 i SUB AGAIN |
|  |  | ; |  |
|  |  | ; | SUBTRACT OF FRI DIDN'T GO |
|  |  | ; |  |
| DB87 |  | : FAIL2 |  |
| DB87 | 2009DD | JSR | FRAIE ; ADD FRI BACK TO FRø |
|  |  | ; |  |
| DB8A | C6F5 | DEC | ZTEMP1 ; DEC LOOP CONTROL |
| DB8C | DøB5 *DB43 | BNE | :NXTQ ; GET NEXT QUIOTIENT BYTE |
|  | 2062DC | ; JSR | RSHFØE ; SHIFT RIGHT FRØ/FRE TO CLEAR |
|  |  |  | EXP |
| DB91 | 4ClADB | JMP | MDEND ; JOIN MULT END UP CODE |

:GETCHAR - Test Input Character


| DBAI |  |  | Skblank |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| DBAI |  |  | SKPbLANK |  |  |  |
| DBAI | A 4 F 2 |  | LDY |  | CIX | ; GET CIX |
| DBA3 | A920 |  | LDA |  | \# $\$ 2 \varnothing$ | ; Get a blank |
| DBA5 | D1F3 |  | $: \text { SBl }$ | CMP | [INBUFF], Y | ; IS CHAR A BLANK |
| DBA 7 | Døø3 | ^DBAC | BNE |  | : SBRTS | ; BR IF NOT |
| DBA9 | C8 |  | INY |  | :SB1 | ; INC TO NEXT |
| DBAA | DøF9 | ^DBA5 | BNE |  |  | ; GO TEST |
| DBAC | 84F2 |  | :SBRTS STY CIX |  |  | ; SET NON BLANK INDEX |
| DBAE | 60 |  | RTS |  |  | ; RETURN |
|  |  |  | ; TSTNUM-TEST CHAR AT CIX FOR NUM |  |  |  |
|  |  |  | ; TSTNUM - RTNS CARRY SET I |  |  | NuM |
| DBAF |  |  |  |  |  |  |  |  |
| DBAF | A4F 2 |  | LDY |  | CIX | ;GET INDEX |
| DBBI | BlF3 |  | LDA |  | [INBUFF], Y | ;AND GET CHAR |
| DBB3 | 38 |  | SEC |  |  |  |
| DBB4 | E930 |  | SBC |  | \# $\$ 30$ | ; SUBTRACT ASCLT ZERO |
| DBB6 | 9018 | *DBDØ | BCC |  | :TSNFAIL | ; BR CHAR<ASCLT ZERO |
| DBB8 | C90A |  | CMP |  | \#\$øA | ;TEST GT ASCLT 9 |
| DBBA | $6 \varnothing$ |  | RTS |  |  | ; DONE |

## :TSTCHAR - Test to See if This Can Be a Number



[^2]LDX \#FR2+1 ; POINT TO 1ST MANTISSA BYTE

```
}

\section*{Source Code}


\section*{NORM - Normalize Floating Point Number}
\begin{tabular}{|c|c|c|c|c|}
\hline DCøø & & NORM & & \\
\hline DCøø & A 200 & LDX & \#Ø & ; GET ZERO \\
\hline DCØ2 & 86DA & STX & FRØ+FPREC & ; FOR ADD NORM SHIFT IN A ZERO \\
\hline DCØ4 & & NORMI & & \\
\hline DCø4 & A204 & LDX & \#FMPREC-1 & ; GET MAX \# OF BYTES TO SHIFT \\
\hline DCø6 & A5D4 & LDA & FRø & ; GET EXPONENT \\
\hline DCø8 & Fø2E * DC38 & BEQ & : NDONE & ; IF EXP=Ø, \# = \\
\hline DCØA & & : NORM & & \\
\hline DCØA & A5D5 & LDA & FRøM & ; GET IST BYTE OF MANTISSA \\
\hline \multirow[t]{3}{*}{DCøC} & DØ1A *DC28 & BNE & :TSTBIG & ; IF NOT \(=\emptyset\) THEN NO SHIFT \\
\hline & & & SHIFT 1 BYTE LEFT & \\
\hline & & ; & & \\
\hline DCDE & Аøøø & LDY & \#ø & ; GET INDEX FOR 1ST MOVE BYTE \\
\hline DC1Ø & & : NSH & & \\
\hline DC1ø & B9D6øø & LDA & FRØM+1, Y & ; GET MOVE BYTE \\
\hline DC13 & 99D50ø & STA & FRØM, Y & ; STORE IT \\
\hline DC16 & C8 & INY & & \\
\hline DC17 & Cøø5 & CPY & \#FMPREC & ; ARE WE DONE \\
\hline \multirow[t]{5}{*}{DC19} & 9ØF5 ^DC1ø & BCC & : NSH & ; IF NOT SHIFT AGAIN \\
\hline & & ; & & \\
\hline & & ; & & \\
\hline & & ; & DECREMENT EXPONENT & \\
\hline & & ; & & \\
\hline \multirow[t]{2}{*}{DC1B} & C6D4 & DEC & FRø & ; DECREMENT EXPONENT \\
\hline & & ; & & \\
\hline DClD & CA & DEX & & ; DEC COUNTER \\
\hline \multirow[t]{4}{*}{DClE} & DØEA ^DCØA & BNE & : NORM & ; DO AGAIN IF NEEDED \\
\hline & & ; & & \\
\hline & & ; & & \\
\hline & & ; & & \\
\hline DC2Ø & A5D5 & LDA & FRØM & ; IS MANTISSA STILL \(\emptyset\) \\
\hline DC22 & Døø4 ^DC28 & BNE & :TSTBIG & ; IF NOT, SEE IF TOO BIG \\
\hline DC24 & 85D4 & STA & FRø & ; ELSE ZERO EXP \\
\hline DC26 & 18 & CLC & & \\
\hline \multirow[t]{2}{*}{DC27} & 60 & R'TS & & \\
\hline & & ; & & \\
\hline DC28 & & : TSTBIG & & \\
\hline DC28 & A5D4 & LDA & FRø & ; GET EXPONENT \\
\hline DC2A & 297F & AND & \# \$ 7F & ; AND OUT SIGN BIT \\
\hline DC2C & C971 & CMP & \#49+64 & ; IS IT < 49+64? \\
\hline DC2E & \(9001{ }^{\wedge}\) DC31 & BCC & : TSTUND & ; IF YES, TEST UNDERFLOW \\
\hline DC3ø & 60 & RTS & & ; SO RETURN \\
\hline DC31 & & : TSTUND & & \\
\hline DC31 & C9ØF & CMP & \#-49+64 & ; IS IT > \(=-49+64\) ? \\
\hline DC33 & Bøø3 *DC38 & BCS & : NDONE & ; IF YES, WE ARE DONE \\
\hline \multirow[t]{2}{*}{DC35} & 2ø44DA & JSR & ZFRø & ; ELSE \# IS ZERO \\
\hline & & , & & \\
\hline DC38 & & :NDONE & & \\
\hline DC38 & 18 & CLC & & ; CLEAR CARRY FOR GOOD RETURN \\
\hline DC39 & \(6 \varnothing\) & RTS & & \\
\hline
\end{tabular}

\section*{Source Code}

\section*{RSHFT0 - Shift FR0 Right/Increment Exponent}

RSHFT1 - Shift FR1 Right/Increment Exponent


RSHFOE - Shift FR0/FRE 1 Byte Right [They Are Contiguous]
\begin{tabular}{|c|c|c|c|c|c|c|}
\hline DC62 & & & \multicolumn{4}{|l|}{RSHFØE} \\
\hline DC62 & A2øA & & LDX & \#FMPREC*2 & ; & GET LOOP CONTROL \\
\hline & & & ; & & & \\
\hline DC64 & & & \multicolumn{4}{|l|}{: NXTB1} \\
\hline DC64 & B5D4 & & LDA & FRø, X & ; & GET A BYTE \\
\hline DC66 & 95D5 & & STA & FRø+1, X & ; & MOVE IT OVER I \\
\hline & & & \multicolumn{4}{|l|}{\multirow[t]{2}{*}{; DEX ; DEC COUNTER}} \\
\hline DC68 & CA & & & & & \\
\hline DC69 & \(10 \mathrm{F9}\) & *DC64 & BPL & : NXTB1 & ; & MOVE NEXT BYTE \\
\hline DC6B & A9øø & & LDA & \# \(\emptyset\) & ; & GET ZERO \\
\hline DC6D & 85D4 & & STA & FRø & ; & SHIFT IT IN \\
\hline DC6F & 60 & & RTS & & & \\
\hline
\end{tabular}
:CVFR0 - Convert Each Byte in FR0 to 2 Characters in LBUFF


CONVERT A BYTE

\section*{Source Code}


\section*{:STNUM - Put ASCII Number in LBUFF}
\begin{tabular}{lll} 
* ON ENTRY & A - DIGIT TO BE CONVERTED TO ASCII \\
\(\star\) & & AND PUT IN LBUFF \\
\(*\) & \(\mathrm{Y}-\) INDEX IN LBUFF
\end{tabular}

\section*{:STCHAR - Store Character in A in LBUFF}
\begin{tabular}{llrll} 
DC9D & & STNUM & & \\
DC9D & ORA & \(\# \$ 3 \varnothing\) & ; CONVERT TO ASCII \\
DC9F & & STCHAR & & \\
DC9F & \(998 \emptyset \emptyset 5\) & STA & LBUFF, Y & ; PUT IN LBUFF \\
DCA2 & C8 & INY & & INCR LBUFF POINTER \\
DCA3 & \(6 \varnothing\) & RTS & &
\end{tabular}

\section*{:FNZER0 - Find Last Non-zero Character in LBUFF}


:GETDIG - Get Next Digit from FR0

:DECINB — Decrement INBUFF
\begin{tabular}{|c|c|c|c|c|c|c|c|c|}
\hline DCC 1 & & : DECINB & & & & & & \\
\hline DCC1 & 38 & SEC & & ; & SUBTRACT & ONE & FROM & INBUFF \\
\hline DCC2 & A5F3 & LDA & INBUFF & ; & X & & & \\
\hline DCC4 & E901 & SBC & \# 1 & ; & X & & & \\
\hline DCC6 & 85F3 & STA & INBUFF & ; & X & & & \\
\hline DCC8 & A5F4 & LDA & INBUFF+1 & ; & X & & & \\
\hline DCCA & E900 & SBC & \# \(\varnothing\) & ; & X & & & \\
\hline DCCC & 85F4 & STA & INBUFF+1 & ; & X & & & \\
\hline DCCE & 60 & RTS & & & & & & \\
\hline
\end{tabular}

MDESUP - Common Set-up for Multiply and Divide Exponent
\begin{tabular}{|c|c|c|c|c|}
\hline & & * & ON EXIT & FRI - FRI EXP WITH OUT SIGN A - FRø EXP WITHOUT SIGN FRSIGN - SIGN FOR QUOTIENT \\
\hline DCCF & & MDESUP & & \\
\hline DCCF & A5D4 & LDA & FRø & ; GET FRØ EXPONENT \\
\hline DCD1 & 45Eø & EOR & FR1 & ; GET FR1 EXPONENT \\
\hline DCD3 & 2980 & AND & \# \$8Ø & ; AND OUT ALL BUT SIGN BIT \\
\hline DCD5 & 85EE & STA & FRS IGN & ; SAVE SIGN \\
\hline DCD7 & 06E0 & ; ASL & FRl & ; SHIFT OUT SIGN IN FR1 EXP \\
\hline DCD9 & 46Eø & LSR & FR1 & ; RESTORE FRI EXP WITHOUT SIGN \\
\hline DCDB & A5D4 & LDA & FRø & ; GET FRØ EXP \\
\hline DCDD & 297F & AND & \# \$ 7F & ; AND OUT SIGN BIT \\
\hline DCDF & 60 & RTS & & \\
\hline
\end{tabular}

MDSUP - Common Set-up for Multiply and Divide


\section*{Source Code}
\begin{tabular}{|c|c|c|c|c|c|c|}
\hline DCF2 & 290F & & AND & \# ® \(^{\text {F }}\) & ; & AND OUT HIGH ORDER NIBBLE \\
\hline DCF4 & 85E6 & & STA & FR2 & ; & STORE TO FINISH SHIFT \\
\hline & & ; & & & & \\
\hline DCF6 & A905 & & LDA & \#FMPREC & ; & SET LOOP CONTROL \\
\hline DCF8 & 85 F 5 & & STA & ZTEMP1 & ; & X \\
\hline & & ; & & & & \\
\hline DCFA & 2034DD & & JSR & MVFRøE & ; & MOVE FRø TO FRE \\
\hline DCFD & 2044DA & & JSR & ZFRø & ; & CLEAR FRØ \\
\hline DDØØ & 60 & ; & RTS & & & \\
\hline & & ; & & & & \\
\hline
\end{tabular}

FRA


\section*{MVFR12 - Move FR1 to FR2}
\begin{tabular}{|c|c|c|c|c|c|c|}
\hline DD28 & & \multicolumn{5}{|l|}{MVFRI2} \\
\hline DD28 & Aøø5 & LDY & \#FMPREC & ; & SET COUNTER & \\
\hline DD2A & & \multicolumn{5}{|l|}{:MV2} \\
\hline DD2A & B9Eøøø & LDA & FRI, Y & ; & GET A BYTE & \\
\hline DD2D & 99E6ØØ & STA & FR2, Y & ; & STORE IT & \\
\hline & & \multicolumn{5}{|l|}{: DEY . DEC COUNTER} \\
\hline DD30 & 88 & DEY & & ; & DEC COUNTER & \\
\hline DD31 & 1 10F7 *DD2A & BPL & :MV2 & ; & IF MORE TO MOVE, & DO IT \\
\hline DD33 & 60 & RTS & & & & \\
\hline
\end{tabular}

\section*{Source Code}

\section*{MVFR0E - Move FR0 to FRE}
\begin{tabular}{|c|c|c|}
\hline DD34 & & MVFRØE \\
\hline DD34 & AØ05 & LDY \\
\hline DD36 & & : MV1 \\
\hline DD36 & B9D400 & LDA \\
\hline DD39 & 99DAØØ & STA \\
\hline & & ; \\
\hline DD3C & 88 & DEY \\
\hline DD3D & 10F7 *DD36 & BPL \\
\hline DD3F & 60 & RTS \\
\hline
\end{tabular}

\section*{Polynomial Evaluation}


\section*{Floating Load/Store}
\begin{tabular}{|c|c|c|c|c|c|c|}
\hline & & * & LOAD & FRØ FROM [X,Y] & \(\mathrm{X}=\mathrm{LSB}, \mathrm{Y}=\mathrm{MSB}\), & USES FLPTR [PGØ] \\
\hline DD89 & 86FC & FLDGR & STX & FLPTR & ; SET FLPTR \(\Rightarrow\) & [X,Y] \\
\hline DD8B & 84FD & STY & & FLPTR+1 & & \\
\hline DD8D & Aøø 5 & FLDøP & LDY & \#FPREC-1 & ; \# BYTES ENTER & HERE W/FLPTR SET \\
\hline DD8F & B1FC & FLD61 & LDA & [FLPTR], Y & ; MOVE & \\
\hline DD91 & 99D4øø & STA & & FRØ, Y & & \\
\hline DD94 & 88 & DEY & & & & \\
\hline DD95 & 10F8 * DD8F & BPL & & FLDดl & COUNT \& LOOP & \\
\hline DD97 & 60 & RTS & & & & \\
\hline & & * & & & & \\
\hline & & * & LOAD & FRI FROM [X, Y ] & OR [FLPTR] & \\
\hline DD98 & 86FC & FLDIR & STX & FLPTR & ; FLPTR \(\Rightarrow\) [ \(\mathrm{X}, \mathrm{Y}]\) & \\
\hline
\end{tabular}

\section*{Source Code}


\section*{Source Code}


\section*{Source Code}
\(\left.\begin{array}{lllll}\text { DEA4 } & \text { 2ø98DD } & \text { JSR } & \text { FLDlR } & \\ \text { DEA7 } & \text { 2ø66DA } & \text { JSR } & \text { FADD } & ; \mathrm{X}+\mathrm{C} \\ \text { DEAA } & \text { A2E6 } & \text { LDX } & \text { \#FPSCR\&\$FF } & \\ \text { DEAC } & \text { Aøø5 } & \text { 2øA7DD } & \text { LDY } & \text { \#FPSCR/\$1øø }\end{array}\right]\)

\section*{LOG10[X]}



Source Code
\begin{tabular}{|c|c|c|c|c|c|}
\hline DFD8 & BF14283156
\[
\emptyset 4
\] & \multicolumn{2}{|r|}{- BYTE \$} & \multicolumn{2}{|l|}{\$BF, \$14, \$28, \$31, \$56, \$¢4 ; - . \(14283156 \emptyset 4\)} \\
\hline DFDE & 3F19999877 & \multicolumn{2}{|r|}{\multirow[t]{2}{*}{- BYTE \$}} & \multirow[t]{2}{*}{\$3F, \$19,\$99,\$98, \$77, \$44 ;} & \multirow[t]{2}{*}{; 0.1999987744} \\
\hline & 44 & & & & \\
\hline DFE4 & BF33333331 & \multicolumn{2}{|r|}{\multirow[t]{2}{*}{. BYTE \$}} & \multirow[t]{2}{*}{\$BF, \$33, \$33, \$33, \$31, \$13} & -Ø.3333333113 \\
\hline & 13 & & & & \\
\hline \multirow[t]{3}{*}{DFEA} & 3F99999999 & \multirow[t]{2}{*}{FP9} & \multirow[t]{2}{*}{. BYTE} & E \$3F,\$99, \$99, \$99, \$99,\$99 & ; 0.999999999 \\
\hline & 99 & & & & \\
\hline & = ØøøВ & NATCF & E EQU & (*-ATCOEF)/FPREC & \\
\hline \multirow[t]{2}{*}{DFFØ} & 3F78539816 & \multirow[t]{2}{*}{PIOV4} & \multirow[t]{2}{*}{\[
4 \text {. BYTE }
\]} & E \(\quad 3 \mathrm{~F}, \$ 78, \$ 53, \$ 98, \$ 16, \$ 34\) & ; \(\mathrm{PI} / 4=\mathrm{ARCT}\) \\
\hline & 34 & & & & \\
\hline
\end{tabular}

\section*{Atari Cartridge Vectors}
\begin{tabular}{|c|c|c|c|c|}
\hline DFF6 & \(=\mathrm{BFF} 9\) & ORG & \multicolumn{2}{|l|}{CRTGI} \\
\hline BFF9 & & \multicolumn{3}{|l|}{SCVECT} \\
\hline BFF9 & 60 & RTS & & \\
\hline BFFA & ØロАø & DW & COLDSTART & - COLDSTART ADDR \\
\hline BFFC & Øø & DB & \(\emptyset\) & ; CART EXISTS \\
\hline BFFD & 05 & DB & 5 & ; FLAG \\
\hline BFFE & F9BF & DW & SCVECT ; & COLDSTART ENTRY ADDR \\
\hline & & & \multicolumn{2}{|l|}{End of BASIC} \\
\hline Сøøロ & & END & & \\
\hline
\end{tabular}

\section*{Appendix \(\mathbf{A}\)}

\section*{Macros in Source Code}

The following is a listing of the macros used in this source listing. You will be able to tell when a macro was used by a plus ( + ) sign to the left of the hex code produced in column two by the assembler.


\section*{Syntax Table Macro}
```

THIS MACRO IS USED TO SIMULATE THE ACTION OF THE ORIGINAL
ASSEMBLER IN HANDLING SPECIAL SYNTAX TABLE PSEUDO OPS AND
OPERANDS
THE 'SYN' MACRO EXAMINES UP TO 4 ARGUMENTS FOR CERTAIN SPECIAL
CASE NAMES.
IF THE NAME 'JS' IS FOUND, IT GENERATES A SPECIAL 'RELATIVE
SYNTAX JSR' TO THE LABEL FOUND IN THE NEXT PARAMETER

```

\section*{Appendix A}
```

; IF THE NAME 'AD' IS FOUND, IT GENERATES A WORD ADDRESS OF
THE LABEL FOUND IN THE NEXT PARAMETER
; ANY OTHER NAME IS ASSUMED TO BE A SIMPLE BYTE VALUE
SYN: MACRO
:SYAR2 SET '=%2'<>'='
:SYAR3 SET '=83'<>'='
:SYAR4 SET '=%4'<>'='
IF '%1' = 'JS'
%L
ELSE
IF '%1' = 'AD'
%L
DW (%2)
SET \emptyset
ELSE
8L
DB %l
ENDIF
ENDIF
IF :SYAR2
IF (% % % (% ((%3-*)\&\$7F) xOR \$4\emptyset )
:SYAR3 SET Ø
ELSE
IF '%2' = 'AD'
DW (%3)
:SYAR3
SET Ø
ELSE
DB %2
ENDIF
ENDIF
ENDIF
IF :SYAR3
IF '%3' = 'JS'
DB \$8\emptyset+(((%4-*)\&\$7F) XOR \$4Ø )
:SYAR4 SET
ELSE
IF '%3' = 'AD'
DW (84)
:SYAR4
SET
ELSE
DB %3
ENDIF
ENDIF
ENDIF

| IF | : SYAR4 |  |
| :---: | :---: | :---: |
| IF |  |  |
| DB | \$8Ø+( ( $(85-*) \& \$ 7 \mathrm{~F})$ | XOR \$40 |
| ELSE |  |  |
| IF | '\%4' $=$ ' $\mathrm{AD}^{\prime}$ |  |
| DW | (\%5) |  |
| ELSE |  |  |
| DB | \% 4 |  |
| ENDIF |  |  |
| ENDIF |  |  |

        ENDM
    ```

\section*{Appendix B}

\section*{The Bugs in Atari BASIC}

Yes, it's true. There are some bugs in Atari BASIC. Of course, that's not surprising, since Atari released the product as ROM without giving the authors a chance to do second-round bugfixing. But what hurts, a little, is that most of the fixes for the bugs have been known since June of 1979.

As this book is being written, rumor has it that at last Atari is in the final stages of releasing a new version of the BASIC ROMs. Unfortunately, these modified ROMs will appear too late for us to comment upon them in this edition. On the other hand, there are supposed to be fewer than twenty fixes implemented (which isn't a bad record for a product as mature as Atari BASIC), so those of you who are willing to PEEK around a bit can use this listing as at least a road map to the new ROMs.

In any case, though, we thought it would be appropriate to mention a few of the bugs we know about, show you why they exist, and tell how we fixed them back there in the summer of '79.

\section*{The Editing and String Bug}

In the course of editing a BASIC program, sometimes the system loses all or part of the program, or it simply hangs. Often, even SYSTEM RESET will not return control to the user.

Also, string assignments that involve the movement of exact multiples of 256 bytes do not move the bytes properly. For example, \(\mathrm{A} \$=\mathrm{B} \$(257,512)\) would actually move bytes 513 through 768 of \(\mathrm{B} \$\) into bytes 257 through 512 of \(\mathrm{A} \$\), even if neither string were DIMensioned to those values.

Both of these are really the same bug. And both are caused because we strove to be a little too efficient.

There are many ways to move strings of bytes using the 6502 's instruction set. The simplest and most-used methods, though, are excruciatingly slow. So Paul and Kathleen invented a super-fast set of move-memory routines, one for

\section*{Appendix B}
moving up in memory (EXPAND, at \$A881) and one for moving down in memory (CONTRACT, at \$A8FD). Unfortunately, the routines are very complex (which is what makes them fast) and difficult to interface with properly. And so a bug crept into CONTRACT.

Take a look at the code of FMOVER (\$A947). When we get here, we expect MVLNG to contain the complement of the least significant byte of the move length while MVLNG + 1 contains its most significant byte. But look what happens if the original move length was, for example, \(\$ 200\). The complement of the least significant byte \((\$ 00)\) is still zero \((\$ 00)\), so the BEQ to :CONT4 occurs immediately.

But by then, the \(X\) register contains the number of pages to move plus one ( X would contain 3 in this example), so we increment it (it becomes 2 ) and go to label :CONT3, where we bump the high-order byte of both the source and destination addresses. Ah, but therein lies the rub! We haven't yet done anything with the first values in those source and destination addresses, so we have effectively skipped 256 bytes of each!

The solution is to replace the BEQ :CONT4 at \$A94E with the following code:

\section*{DEX \\ BNE :CONT2 RTS}

Do you see the difference? If we enter with MVLNG equal to zero, we immediately move 256 bytes (at :CONT2) before ever attempting to change the source and destination addresses.

And this fix works, honest. We've been using it like this for over two years in BASIC A + .

\section*{Minus Zero}

Taking the unary minus of a number ( \(\mathrm{A}=0:\) PRINT -A) can result in garbage. Usually, this garbage will not affect subsequent calculations, but it does print strangely. And how did this come about?

We simply forgot to take into consideration the fact that zero doesn't really have a sign. Look at the code for the unary minus operator (XPUMINUS, at \$ACA8). Do you see the problem? We simply invert the most significant bit (the sign bit) of the floating point number in FR0.

\section*{Appendix B}

What we should have coded would be something like this:
LDA FR0
BEQ :NOINVERT
EOR \#\$80
STA FRO
:NOINVERT
Luckily, this is not too severe a problem to the BASIC user (one can always use "PRINT 0-A" instead of "PRINT - \(\mathrm{A}^{\prime}\) "), but just think - it only cost two bytes to fix this bug.

\section*{LOCATE and GET}

The GET statement does not reinitialize its buffer pointer, so it can do nasty things to memory if used directly after a statement which has changed the system buffer pointer. For example, GET can change the line number of a DATA statement if it is used after a READ. Also, the same problem exists for the LOCATE statement, since it calls GET.

From BASIC, the easiest solution is to use a function or statement which is known to reset the pointer. Coding " \(\mathrm{XX}=\) STR\$(0)" works just fine, as does PRINTing any number.

Within the source listing, the problem exists at location \(\$ B C 82\), label GET1. If the code had simply read as follows, there would be no bug:

GET1
JSR INTLBF ; reset buffer pointer
LDA \#ICGTC ; continue as before

\section*{INPUT and READ}

Using either an INPUT or READ statement without a following variable does not cause a syntax error (as it should). Then, attempting to execute a statement such as 20 INPUT can cause total system lock-up.

The solution from BASIC? Be careful and don't do it.
And this is one bug that we will not show the fix for, simply because it's too long and involved. We will, however, point to labels :SINPUT and :SREAD (at locations \$A6F4 and \$A6F5) in the Syntax Tables and show why the bug exists.

Note that the :SINPUT does a syntax call (SYN JS,) to the : OPD syntax, which looks for - but does not insist upon - a file number specifier (\# < numeric expression >). Then the

\section*{Appendix B}
syntax joins with :SREAD, which looks for zero or more variables.

Oops! Zero or more? Shouldn't that be one or more? That's where the problem lies.

\section*{Do Not Use NOT}

In all too many cases, the use of the NOT operator is guaranteed to get you in trouble. If you don't believe it, try this: PRINT NOT NOT 1.

The explanation of why the bug occurs is too lengthy to give in detail here; suffice it to say that the precedence of NOT is wrong. Remember the Operator Precedence Table we displayed in Chapter 8 of Part 2? Look at what you got for the go-onto-stack and come-off-stack precedence values for NOT.

Or look at location \$AC57, the NOT entry in OPRTAB. NOT uses a 7 for both its precedence values. But wait a minute. If two operators have the same apparent precedence (as in NOT NOT A or even \(A+B+C\) ), the expression executor will pop the first one off the stack and execute it. But with a unary operator, there is nothing to execute yet.

And the same bug exists for both unary minus and unary plus, so --3 and ++5 don't execute properly. Of course, since unary plus doesn't really do anything, it doesn't matter. In the case of unary minus, though, all but the last minus sign in a string of minus signs is ignored (that is, -3 produces -3 as a result, instead of +3 , as it should). But, by an incredible coincidence, the damage that unary minus causes is invisible to Execute Expression as a whole and only produces the error noted.

The fix? Well, if we want to leave NOT where it is in the order of things, the only way is to restructure the whole precedence table. But if we are willing to accord it a very high precedence, like unary plus and minus, we can fix it - and plus and minus - by changing the bytes at \$AC57, \$AC64, and \$AC65 to \$DC. And, thanks to the differing go-onto-stack and come-off-stack values, we can stack as many NOTs, pluses, or minuses as we want.

Are these all the bugs we know about that can be fixed easily? No. But these are the easiest to understand or the easiest to fix, and we thought they were instructive.

Of course, unless you have an EPROM board and burner handy, you may not be able to take advantage of these fixes.

\section*{Appendix B}

But at least now you may be able to work around them as you program with good old buggy-version Atari BASIC.

And take heart. Remember Richard's Rule: Any nontrivial piece of software has bugs in it. And the corollary: Any piece of software which is bug-free is trivial.


Appendix C

\title{
Labels and Hexadecimal Addresses
}
\begin{tabular}{|c|c|c|}
\hline & AADD & AF52 \\
\hline & AAPSTR & AB98 \\
\hline \multirow[t]{2}{*}{n} & ADC & AF53 \\
\hline & ADFLAG & の日Bl \\
\hline \multirow[t]{11}{*}{n} & AFP & D8øø \\
\hline & AMUL1 & AF5D \\
\hline & AMUL 2 & AF46 \\
\hline & APHM & のøのЕ \\
\hline & ARGOPS & Øø8ø \\
\hline & ARGP2 & ACø6 \\
\hline & ARGPOP & ABF2 \\
\hline & ARGPUS & ABBA \\
\hline & ARGSTK & øø8Ø \\
\hline & ARSLVL & ØøAA \\
\hline & ARSTKX & ØロAA \\
\hline \multirow[t]{12}{*}{n} & ASCIN & D8øの \\
\hline & ASLA & mac \\
\hline & ATAN & BE 77 \\
\hline & ATAN1 & BE9A \\
\hline & ATAN2 & BED4 \\
\hline & ATCOEF & DFAE \\
\hline & ATEMP & øøAF \\
\hline & ATNOUT & BEE2 \\
\hline & BININT & のøD4 \\
\hline & BOTH & BDB3 \\
\hline & BRKBYT & øø11 \\
\hline & BYELOC & E471 \\
\hline \multirow[t]{3}{*}{n} & BYTE & mac \\
\hline & C & 0044 \\
\hline & BYELOC & E471 \\
\hline \multirow[t]{14}{*}{n} & BYTE & mac \\
\hline & C & のø44 \\
\hline & CAASN & のø2D \\
\hline & CACOM & øø3C \\
\hline & CADR & øø43 \\
\hline & CALPRN & のø38 \\
\hline & CAND & Øø2A \\
\hline & CASC & øø4ø \\
\hline & CCHR & Øø3E \\
\hline & CCOM & Øø12 \\
\hline & CCR & ø016 \\
\hline & CDATA & øøø1 \\
\hline & CDIV & Ø027 \\
\hline & CDLPRN & ø039 \\
\hline n & CDOL & øø13 \\
\hline \multirow[t]{13}{*}{n} & CDQ & のø1ø \\
\hline & CDSLPR & のø3B \\
\hline & CEOS & øø14 \\
\hline & CEQ & 0022 \\
\hline & CERR & のø37 \\
\hline & CEXP & のø23 \\
\hline & CFFUN & のø3D \\
\hline & CFLPRN & の日3A \\
\hline & CFOR & øøø8 \\
\hline & CGE & ØO1F \\
\hline & CGOSUB & のøøС \\
\hline & CGS & のø18 \\
\hline & CGT & のø21 \\
\hline
\end{tabular}
\begin{tabular}{|c|c|}
\hline CGTO & Øø17 \\
\hline CILET & Ø036 \\
\hline CIO & E456 \\
\hline CIX & ØロF2 \\
\hline CLALLI & BD4F \\
\hline CLE & Øø1D \\
\hline CLEN & Øø42 \\
\hline CLIST & Øロø4 \\
\hline CLPRN & Øø2B \\
\hline CLSALL & BD \\
\hline CLSYSl & BC \\
\hline CLSYSD & BCF \\
\hline CLT & Øø2Ø \\
\hline CMINU & Øø26 \\
\hline CMUL & Øø24 \\
\hline CNE & 0ø1E \\
\hline CNFNP & Øø44 \\
\hline CNOT & Øロ28 \\
\hline COLDI & Aøø8 \\
\hline COLDST & Аøøø \\
\hline COLOR & øøC8 \\
\hline COMCNT & ØøВø \\
\hline CON & Øø1 \\
\hline CONTLO & A8FB \\
\hline CONTRA & A8F \\
\hline COPEN & BBB6 \\
\hline COR & Øロ29 \\
\hline cos & BDB1 \\
\hline COX & \(0 \square 94\) \\
\hline PC & Øø9D \\
\hline CPLUS & Ø025 \\
\hline CPND & 001 C \\
\hline CR & Ø09B \\
\hline CREAD & øø22 \\
\hline CREGS & Ø2C4 \\
\hline CRPRN & Øø2C \\
\hline CRTGI & BFF9 \\
\hline CSASN & Øø2 \\
\hline CSC & Ø015 \\
\hline CSEQ & Øø34 \\
\hline CSGE & Øロ31 \\
\hline CSGT & Ø03 3 \\
\hline CSLE & Øø2F \\
\hline CSLPRN & ø037 \\
\hline CSLT & 0032 \\
\hline CSNE & Øø3ø \\
\hline CSOE & Øø11 \\
\hline CSROP & Øø1D \\
\hline CSTEP & Øø1A \\
\hline CSTR & のø3D \\
\hline CTHEN & Øø1B \\
\hline CTO & Ø019 \\
\hline CUMINU & Øø 36 \\
\hline CUPLUS & Ø035 \\
\hline CUSR & Øø3F \\
\hline CVAFP & D800 \\
\hline CVAL & 0041 \\
\hline VFAS & \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|}
\hline & CVFPI & AD56 \\
\hline & CVIFP & D9AA \\
\hline & DATAD & ØøВ6 \\
\hline & DATALN & øøВ7 \\
\hline n & DCBORG & Ø30ø \\
\hline & DEGFLG & ØØFB \\
\hline & DEGON & Øøø6 \\
\hline & DIGRT & ØøF1 \\
\hline & DIRFLG & のøA6 \\
\hline & DNERR & \(\mathrm{BCB} \varnothing\) \\
\hline & DOSLOC & のøのA \\
\hline & DSPFLG & Ø2FE \\
\hline & ECSIZE & ØØA 4 \\
\hline & EEXP & ØøED \\
\hline & ELADVC & BADD \\
\hline & ENDSTA & Øø8E \\
\hline & ENDVVT & Ø088 \\
\hline & ENTDTD & ØøВ4 \\
\hline & EPCHAR & の日5D \\
\hline & ERBRTN & B920 \\
\hline & ERGFDE & B922 \\
\hline & ERLTL & B924 \\
\hline & ERNOFO & B926 \\
\hline & ERNOLN & B928 \\
\hline n & ERON & B93E \\
\hline & EROVFL & B92A \\
\hline & ERRAOS & B92C \\
\hline & ERRDIM & B92E \\
\hline & ERRDNO & B918 \\
\hline & ERRINP & B930 \\
\hline & ERRLN & B932 \\
\hline & ERRNSF & B916 \\
\hline & ERRNUM & øøВ9 \\
\hline & ERROOD & B934 \\
\hline & ERROR & B940 \\
\hline & ERRPTL & B91A \\
\hline & ERRSAV & のøく3 \\
\hline & ERRSSL & B936 \\
\hline & ERRVSF & B938 \\
\hline & ERSVAL & B91C \\
\hline & ERVAL & B93A \\
\hline & ESIGN & ØøEF \\
\hline & EVAADR & Øロø2 \\
\hline & EVADI & Øøø 4 \\
\hline & EVAD2 & のøロ6 \\
\hline n & EVARRA & のø4ø \\
\hline & EVDIM & のøø1 \\
\hline n & EVNUM & のøø1 \\
\hline & EVSADR & 0ø02 \\
\hline & EVSCAL & øøøø \\
\hline & EVSDIM & ロøø6 \\
\hline & EVSDTA & øøø2 \\
\hline & EVSLEN & øøø 4 \\
\hline & EVSTR & øø8ø \\
\hline & EVTYPE & øøøø \\
\hline & EVVALU & Øøø2 \\
\hline & EXECNL & A95F \\
\hline & ExECNS & A962 \\
\hline
\end{tabular}

EXEXPR AAEØ EXOPOP ABØB EXP DDCØ EXP1 DEØ3 EXPlø DDCC EXP2 DE2ø EXP3 DE26 EXPAND A881 EXPERR DE4B EXPINT AB2E EXPLOW A87F EXPOUT DE4A EXPSGN DE 39 EXSVOP ØøAB EXSVPR ØØAC FADD DA66
\(n\) FASC D8E6 FBODY ØøøC FCHRFL ØøFの FDB mac FDIV DB28 FHALF DF6C FIXRST B825
FLDØ1 DD8F
n FLDØP DD8D FLDØR DD89 FLDIl DD9E
\(n\) FLDIP DD9C FLDIR DD98 FLIM Øøøø FLIST BAD5
n FLOGIØ DFØ1 FLPTR ØøFC FMOVE DDB6 FMOVE 1 DDB8 FMOVER A947 FMPREC Øøø5 FMUL DADB
\(n\) FNTAB A829
FONE DE8F FP9S DFEA FPI D9D2 FPONE BE71 FPORG D8øø FPREC Øøø6 FPSCR Ø5E6 FPSCR1 Ø5EC FPTR2 ØØFE FRØ ØØD4 FRØM ØØD5 FR1 ØøEの FRIM ØøEl FR2 Ø日E6 FRAlØ DDØI FRAIE DD日9 FRA2Ø DDØ5 FRA2E DDØF FRADD AD3B

Appendix C
\begin{tabular}{|c|c|c|}
\hline & FRCMP & AD35 \\
\hline & FRCMPP & AD32 \\
\hline & FRDIV & AD4D \\
\hline & FRE & のøDA \\
\hline & FRMUL & AD47 \\
\hline & FRSIGN & ØøEE \\
\hline & FRSUB & AD41 \\
\hline & FRUN & BAF7 \\
\hline & FRX & ØøEC \\
\hline & FSCR & の5E6 \\
\hline & FSCR1 & Ø5EC \\
\hline & FSQR & BFø8 \\
\hline & FSTØ1 & DDAD \\
\hline n & FSTØP & DDAB \\
\hline & FSTØR & DDA 7 \\
\hline & FSTEP & Øøø6 \\
\hline & FSUB & DA60 \\
\hline & FTWO & BF93 \\
\hline & GDIOI & BC22 \\
\hline & GDVCIO & BCld \\
\hline & GET 1 & BC82 \\
\hline & GETIIN & ABE9 \\
\hline & GETINT & ABED \\
\hline & GETLL & A9DD \\
\hline n & GETPIØ & ABD8 \\
\hline & GETPIN & ABD5 \\
\hline & GETSTM & A9A2 \\
\hline & GETTOK & AB3E \\
\hline & GETVAR & AB89 \\
\hline & GFDISP & øøøろ \\
\hline & GFHEAD & øøø 4 \\
\hline n & GFLNO & øøø1 \\
\hline & GFTYPE & のøøø \\
\hline & GIOCMD & BDø4 \\
\hline & GIODVC & BC9F \\
\hline & GIOPRM & BDø2 \\
\hline & GLGO & BA92 \\
\hline & GLINE & BA89 \\
\hline & GLPCX & BAC4 \\
\hline & GLPX & BAC6 \\
\hline n & GNLINE & BA8ø \\
\hline & GNXTL & A9D0 \\
\hline & GRFBAS & ø27ø \\
\hline & GSTRAD & AB9B \\
\hline & GTINTO & ABE3 \\
\hline & GVVTAD & AC28 \\
\hline & HIMEM & Ø2E5 \\
\hline & HMADR & Ø2E5 \\
\hline n & IBUFEX & ØØA9 \\
\hline & ICAUXI & Ø34A \\
\hline & ICAUX2 & Ø34B \\
\hline & ICAUX3 & Ø34C \\
\hline & ICAUX4 & Ø34D \\
\hline & ICAUX5 & Ø34E \\
\hline & ICBAH & Ø345 \\
\hline & ICBAL & Ø344 \\
\hline & ICBLH & 0349 \\
\hline & ICBLL & 0348 \\
\hline & ICCLOS & øøロС \\
\hline & ICCOM & 0342 \\
\hline n & ICDDC & のøøE \\
\hline n & ICDNO & \(\emptyset 341\) \\
\hline & ICDRAW & øø11 \\
\hline n & ICFREE & Ø日FF \\
\hline n & ICGBC & Øøø6 \\
\hline n & ICGBR & Øøロ4 \\
\hline & ICGR & øø1C \\
\hline & ICGTC & 0007 \\
\hline & ICGTR & 0005 \\
\hline n & ICHID & Ø340 \\
\hline & ICLEN & Øø1ø \\
\hline & & 80 \\
\hline
\end{tabular}
n ICMAX ØØØE


RNDDIV BØA8 RNDLOC D2ØA ROLA mac ROM Aøøø RORA mac RSHFøE DC62 RSHFTØ DC3A RSHFTI DC3E RSTPTR B8AF RSTSEO BD99 RTNVAR AC16 RUNINI B8F8 RUNSTK Øø8E SAVCUR Øøbe SAVDEX ØøB3 SCANT ØØAF SCOEF BE41 SCRX Øø55 SCRY Øø54 SCVECT BFF9 SEARCH A462 SETDZ BD72 SETLIN B818 SETLNL B81B SETSEO BD79 SGNFLG øøFø SICKIO BCB9 SIN BDA7 SINDON BE40 SINERR BDA5 SINFl BDF6 SINF3 BEØø SINF4 BE11 SINF5 BDE4
n SINF6 BDD5 SINF7 BDCC
\(n\) SINOVF BDCB SIX Ø48Ø SKBLAN DBAl SKCTL D2øF SKPBLA DBAl SNTAB A4AF SNXI A650 SNX2 A653 \(\begin{array}{ll}\text { SNX3 } & \text { A65D } \\ \text { SOPEN } & \text { BBDI }\end{array}\) SOX \(\quad 0481\) \(\begin{array}{ll}\mathrm{SPC} & 9482 \\ \text { SOR }\end{array}\) \(\begin{array}{ll}\text { SQR } & \text { BEE } 5 \\ \text { SQR1 BFøø }\end{array}\) SQR10 DF66 \(\begin{array}{ll}\text { SQR2 } & \text { BF84 } \\ \text { SQR3 BF8A }\end{array}\) SQRCNT ØØEF SQRDON BF64 SQRERR BEE3 SQRLP BF2A SQROUT BF92 SRCADR \(\varnothing \boxed{6} 9\) SRCNXT A49б SRCSKP ØロAA SREGI D2ø8 SREG2 D20ø SREG3 D201 SSTR BA73 STACK \(048 \varnothing\) STARP Øø8C STENUM ØロAF Stetab anøø STINDE ØロA8 STKLVL बøA9 STMCUR øø8A

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[^0]:    register
    X = the zero page address containing the pointer to the location at which space is to be removed

[^1]:    ;GET ARGUMENT INTO FRø
    ;GET BYTE WITH SIGN
    ;FLIP SIGN BIT
    ; RETURN BYTE WITH SIGN CHANGED ;PUSH ON STACKS

[^2]:    NIBSH0 - Shift FR0 One Nibble Left
    DBE7
    DBE7 A2E7

    ```
    * NIBSH2 - SHIFT FR2 ONE NIBBLE LEFT
    ```

    * NIBSH2 - SHIFT FR2 ONE NIBBLE LEFT
    NIBSH2
    NIBSH2
    LDX \#FR2+1 ; POINT TO 1ST MANTISSA BYTE

    ```
    ```

